# NESTEROV-BASED ALTERNATING OPTIMIZATION FOR NONNEGATIVE TENSOR COMPLETION: ALGORITHM AND PARALLEL IMPLEMENTATION

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# **Nonnegative Tensor Completion**

- Let  $\mathcal{X} \in \mathbb{R}^{I \times J \times K}_+$  be an incomplete tensor, and  $\Omega \subseteq \{1 \dots I\} \times \{1 \dots J\} \times \{1 \dots K\}$  be the set of indices of its known entries [1]. – Also, let  $\mathcal{M} \in \mathbb{R}^{I \times J \times K}$  with

$$\mathcal{M}(i,j,k) = \begin{cases} 1, \text{ if } (i,j,k) \in \Omega, \\ 0, \text{ otherwise.} \end{cases}$$
(1)

- We consider the Nonnegative Tensor Completion (NTC) problem

$$\min_{\mathbf{A} \ge \mathbf{0}, \mathbf{B} \ge \mathbf{0}, \mathbf{C} \ge \mathbf{0}} f_{\Omega}(\mathbf{A}, \mathbf{B}, \mathbf{C}) + \frac{\lambda}{2} \|\mathbf{A}\|_{F}^{2} + \frac{\lambda}{2} \|\mathbf{B}\|_{F}^{2} + \frac{\lambda}{2} \|\mathbf{C}\|_{F}^{2},$$

where  $\mathbf{A} = [\mathbf{a}_1 \cdots \mathbf{a}_R] \in \mathbb{R}_+^{I \times R}$ ,  $\mathbf{B} = [\mathbf{b}_1 \cdots \mathbf{b}_R] \in \mathbb{R}_+^{J \times R}$ ,  $\mathbf{C} = [\mathbf{c}_1 \cdots \mathbf{c}_R] \in \mathbb{R}_+^{K \times R}$ , and  $f_{\Omega}(\mathbf{A},\mathbf{B},\mathbf{C}) = \frac{1}{2} \| \mathcal{M} \circledast (\mathcal{X} - \llbracket \mathbf{A},\mathbf{B},\mathbf{C} \rrbracket) \|_{F}^{2},$ 

with

$$\llbracket \mathbf{A}, \mathbf{B}, \mathbf{C} \rrbracket = \sum_{r=1}^{R} \mathbf{a}_r \circ \mathbf{b}_r \circ \mathbf{c}_r.$$
(4)

### Alternating optimization framework

- We can derive matrix-based equivalent expressions of  $f_{\Omega}$  as

### **Distributed Memory Implementation**





- We assume p processing elements [2].
- We partition the matricization  $\mathbf{X}_{\mathbf{A}}$  into *p* block rows as  $\mathbf{X}_{\mathbf{A}} = \left[ \left( \mathbf{X}_{\mathbf{A}}^{1} \right)^{T} \cdots \left( \mathbf{X}_{\mathbf{A}}^{p} \right)^{T} \right]^{T}$ , with

 $\mathbf{X}_{\mathbf{A}}^{n} \in \mathbb{R}^{\frac{1}{p} \times JK}$ . We partition similarly  $\mathbf{X}_{\mathbf{B}}$  and  $\mathbf{X}_{\mathbf{C}}$ .

- The *n*-th block row of  $X_A$ ,  $X_B$ ,  $X_C$  have been allocated to the *n*-th processing element, for n = 1, ..., p.
- We partition  $\mathbf{A}_l$  into p block rows as  $\mathbf{A}_l = \left[ \left( \mathbf{A}_l^1 \right)^T \cdots \left( \mathbf{A}_l^p \right)^T \right]^T$ , with  $\mathbf{A}_l^n \in \mathbb{R}^{\frac{l}{p} \times R}$ , for n = 1, ..., p.

- The *n*-th processor knows the whole  $A_i$ , but updates the *n*-th block row of  $A_i$ ,  $A_i^n$ , for n = 1, ..., p.

## **Factor Update Implementation**

$$\begin{split} f_{\Omega}(\mathbf{A},\mathbf{B},\mathbf{C}) &= \frac{1}{2} \|\mathbf{M}_{\mathbf{A}} \circledast \left( \mathbf{X}_{\mathbf{A}} - \mathbf{A}(\mathbf{C} \odot \mathbf{B})^{T} \right) \|_{F}^{2} = \frac{1}{2} \|\mathbf{M}_{\mathbf{B}} \circledast \left( \mathbf{X}_{\mathbf{B}} - \mathbf{B}(\mathbf{C} \odot \mathbf{A})^{T} \right) \|_{F}^{2} \\ &= \frac{1}{2} \|\mathbf{M}_{\mathbf{C}} \circledast \left( \mathbf{X}_{\mathbf{C}} - \mathbf{C}(\mathbf{B} \odot \mathbf{A})^{T} \right) \|_{F}^{2}, \end{split}$$

where  $M_A$ ,  $M_B$ ,  $M_C$ , and  $X_A$ ,  $X_B$ ,  $X_C$  are the matrix unfoldings of  $\mathcal{M}$  and  $\mathcal{X}$ , respectively.

- Solving (2) for **A**, **B**, **C** is a non-convex problem.
- Alternating optimization (AO):
  - Initialize  $A_0, B_0, C_0, I = 0$ .

$$1 \quad \mathbf{A}_{l+1} = \operatorname*{argmin}_{\mathbf{A} \ge \mathbf{0}} f_{\Omega}(\mathbf{A}) := \frac{1}{2} \left\| \mathbf{M}_{\mathbf{A}} \circledast \left( \mathbf{X}_{\mathbf{A}} - \mathbf{A} \left( \mathbf{C}_{l} \odot \mathbf{B}_{l} \right)^{T} \right) \right\|_{F}^{2} + \frac{\lambda}{2} \left\| \mathbf{A} \right\|_{F}^{2}.$$

2 
$$\mathbf{B}_{l+1} = \operatorname*{argmin}_{\mathbf{B}>\mathbf{0}} f_{\Omega}(\mathbf{B}) := \frac{1}{2} \|\mathbf{M}_{\mathbf{B}} \otimes (\mathbf{X}_{\mathbf{B}} - \mathbf{B} (\mathbf{C}_{l} \odot \mathbf{A}_{l+1})^{T})\|_{F}^{2} + \frac{\lambda}{2} \|\mathbf{B}\|_{F}^{2}.$$

3 
$$\mathbf{C}_{l+1} = \operatorname*{argmin}_{\mathbf{C} \ge \mathbf{0}} f_{\Omega}(\mathbf{C}) := \frac{1}{2} \left\| \mathbf{M}_{\mathbf{C}} \circledast \left( \mathbf{X}_{\mathbf{C}} - \mathbf{C} \left( \mathbf{B}_{l+1} \odot \mathbf{A}_{l+1} \right)^T \right) \right\|_F^2 + \frac{\lambda}{2} \left\| \mathbf{C} \right\|_F^2$$

- Iterate 1, 2, 3 until convergence.

# **Nonnegative Matrix Completion**

- Let  $\mathbf{X} \in \mathbb{R}^{m \times n}_+$  be an incomplete matrix, and  $\Omega \subseteq \{1 \dots m\} \times \{1 \dots n\}$  be the set of indices of its known entries.
- Also, let  $\mathbf{A} \in \mathbb{R}^{m \times r}_+$ ,  $\mathbf{B} \in \mathbb{R}^{n \times r}_+$ , and  $\mathbf{M} \in \mathbb{R}^{m \times n}$  with

$$\mathbf{M}(i,j) = \left\{ egin{array}{ll} \mathbf{1}, \ \mathsf{if} \ (i,j) \in \Omega, \ \mathbf{0}, \ \mathsf{otherwise}. \end{array} 
ight.$$

- We consider the Nonnegative Matrix Completion (NMC) problem

$$\min_{\mathbf{A}\geq\mathbf{0}}f_{\Omega}(\mathbf{A}):=\frac{1}{2}\|\mathbf{M}\otimes\left(\mathbf{X}-\mathbf{A}\mathbf{B}^{T}\right)\|_{F}^{2}+\frac{\lambda}{2}\|\mathbf{A}\|_{F}^{2}.$$

- The gradient and the Hessian of  $f_{\Omega}$ , at point **A**, are given by

$$abla f_{\Omega}(\mathbf{A}) = -\left(\mathbf{M} \circledast \mathbf{X} - \mathbf{M} \circledast \mathbf{AB}^{T}\right)\mathbf{B} + \lambda \mathbf{A},$$

- The update of  $A_I$  is achieved via the updates of  $A_I^n$ , for n = 1, ..., p:
- The *n*-th processing element uses its local data  $X^n_A$ , as well as the whole matrices  $B_l$  and  $C_l$ , and computes the *n*-th block row of matrix  $A_{l+1}$ ,  $A_{l+1}^n$ , via the Nesterov Matrix Completion algorithm.
- Each processing element broadcasts its output to all other processing elements; this operation can **be implemented via an** Allgather **operation**.

At this point, all processors know  $A_{l+1}$  and are ready for the update of  $B_l$  (and, then, of  $C_l$ ).

# **Numerical Experiments**

- Results obtained from a Message Passing Interface (MPI) implementation of the AO NTC.
- The program is executed on a DELL PowerEdge R820 system with SandyBridge Intel(R) Xeon(R) CPU E5 – 4650v2 (in total, 16 nodes with 40 cores each at 2.4 Gz) and 512 GB RAM per node.
- The matrix operations are implemented using routines of the C++ library Eigen [3].
- The performance metric we compute is the speedup attained using *p* processors.

# **Real Data**

- The MovieLens 10M dataset [4], which contains time-stamped ratings of movies.
- Binning the time into seven-day-wide bins, results in a tensor of size  $71567 \times 65133 \times 171$ .
- The number of samples is 8M (99.99% sparsity).
- We first perform a random permutation on our data to resolve load imbalance issues.
- We measure the completion accuracy by measuring the mean squared error of 2M known ratings with our predictions. The mean squared error we achieved is 0.0033

(For the *n*-th known rating, with indices  $(i_n, j_n, k_n)$ , we compute our prediction after rounding the quantity  $\sum_{r=1}^{R} \mathbf{A}(i_n, :) \circledast \mathbf{B}(j_n, :) \circledast \mathbf{C}(k_n, :)$  to the closest integer.)

# $abla^2 f_{\Omega}(\mathbf{A}) = \left(\mathbf{B}^T \otimes \mathbf{I}\right) \operatorname{diag}^2\left(\operatorname{vec}\left(\mathbf{M}\right)\right) \left(\mathbf{B} \otimes \mathbf{I}\right) + \lambda \mathbf{I}.$

## **Nesterov-Type Algorithm for NMC**

Algorithm 1:  $A_{opt} = N_NMC(X, M, B, A_*)$ **Input:** X, M  $\in \mathbb{R}^{m \times n}_+$ , B  $\in \mathbb{R}^{n \times r}_+$ , A<sub>\*</sub>  $\in \mathbb{R}^{m \times r}_+$ ,  $\lambda$ ,  $\mu$ , L  $\mathbf{W} = -(\mathbf{M} \circledast \mathbf{X})\mathbf{B}$  $oldsymbol{q} = rac{\mu + \lambda}{L + \lambda}$  $\mathbf{A}_0 = \mathbf{Y}_0 = \mathbf{A}_*$  $\alpha_0 = 1, I = 0$ while (1) do  $\nabla f_{\Omega}(\mathbf{Y}_{I}) = \mathbf{W} + (\mathbf{M} \circledast \mathbf{Y}_{I} \mathbf{B}^{T}) \mathbf{B} + \lambda \mathbf{Y}_{I}$ if (*term\_cond* is *TRUE*) then break else  $\mathbf{A}_{l+1} = \left(\mathbf{Y}_l - \frac{1}{L+\lambda} \nabla f_{\Omega}(\mathbf{Y}_l)\right)_+$  $\alpha_{l+1}^2 = (1 - \alpha_{l+1})\alpha_l^2 + q\alpha_{l+1}$  $\beta_{l+1} = \frac{\alpha_l(1-\alpha_l)}{\alpha_l^2 + \alpha_{l+1}}$  $\mathbf{Y}_{l+1} = \mathbf{A}_{l+1} + \beta_{l+1} \left( \mathbf{A}_{l+1} - \mathbf{A}_{l} \right)$ return A<sub>/</sub>.

### **Nesterov Based AO NTC**

Algorithm 2: Nesterov-based AO NTC **Input:** X,  $\Omega$ ,  $A_0 > 0$ ,  $B_0 > 0$ ,  $C_0 > 0$ . I = 0while (1) do  $\mathbf{A}_{l+1} = \mathrm{N}_{\mathrm{NMC}}(\mathbf{X}_{\mathbf{A}}, \mathbf{M}_{\mathbf{A}}, (\mathbf{C}_{l} \odot \mathbf{B}_{l}), \mathbf{A}_{l})$  $\mathbf{B}_{l+1} = \mathrm{N_NMC}(\mathbf{X}_{\mathbf{B}}, \mathbf{M}_{\mathbf{B}}, (\mathbf{C}_l \odot \mathbf{A}_{l+1}), \mathbf{B}_l)$  $\mathbf{C}_{l+1} = \mathrm{N}_{-}\mathrm{NMC}(\mathbf{X}_{\mathbf{C}}, \mathbf{M}_{\mathbf{C}}, (\mathbf{A}_{l+1} \odot \mathbf{B}_{l+1}), \mathbf{C}_{l})$ if (term\_cond is TRUE) then break; endif



Figure: Speedup achieved for the MovieLens 10M dataset of size  $71567 \times 65133 \times 171$  with p cores, for p = 1, 5, 20, 171.

### Synthetic Data

- Synthetic data of the same size and sparsity level.

- True latent factors with i.i.d elements, uniformly distributed in [0, 1].

### Speed Up



# Computation of $W_A$ and $Z_A$

 $W_{A} = (M_{A} \circledast X_{A})(C \odot B), \qquad Z_{A} = (M_{A} \circledast A(C \odot B)^{T})(C \odot B).$ 

- The *i*-th row of  $W_A$ , for i = 1, ..., I, is computed as

 $W_{A}(i,:) = \left(M_{A}(i,:) \circledast X_{A}(i,:)\right)(C \odot B).$ 

- The computation involves the multiplication of a  $(1 \times JK)$  row vector and a  $(JK \times R)$  matrix.
- In order to reduce the computational complexity, we must exploit the sparsity of  $\mathcal{X}$ .
- Let  $nnz_i$  be the number of known entries in the *i*-th horizontal slice of  $\mathcal{X}$ . Also, let these known entries have indices  $(i, j_q, k_q) \in \Omega$ , for  $q = 1, \ldots, nnz_i$ .
- The computation of the *i*-th row of  $W_A$  reduces to

$$\mathbf{W}_{\mathbf{A}}(i,:) = \sum_{q=1}^{nnz_i} \mathcal{X}\left(i, j_q, k_q\right) \mathbf{C}\left(k_q,:\right) \circledast \mathbf{B}\left(j_q,:\right).$$

Efficient computation of 
$$Z_A$$
 can be achieved following similar arguments.



Figure: Speedup achieved for a 71567  $\times$  65133  $\times$  171 tensor with *p* cores, for *p* = 1, 5, 20, 171.

#### References

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