Active Learning on Weighted Graphs Using Adaptive and Non-adaptive Approaches

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Challenge in machine learning applications

- Unlabeled data abundant
- · Labels are expensive and scarce

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When can we expect |U| to be less than |V|?

- · Smoothness: strongly connected nodes will have similar signal
- Small cut size: very few edges with oppositely labeled endpoints compared to the total number of edges

Related Work

Global smoothness based sampling

Sample most informative nodes for good signal estimation



- [Guillory and Bilmes 'I I]
- [Ji and Han '12]
- [Anis, G., Ortega '14]

non-adaptive: sample all at once

Boundary refinement sampling

Sample in order to recover the boundary nodes



- [Zhu, Lafferty, Ghahramani '03]
- [Osugi, Kim, Scott '05]
- [Dasarathy, Nowak, Zhu '15]

adaptive: sample one by one

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Which approach is better: depends on error tolerance/sampling budget

A new sampling algorithm: Weighted S^2

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- generalization of S^2 algorithm [Dasarathy, Nowak, Zhu '15]
 - S^2 algorithm assumes an unweighted graph
 - weights capture additional info. about node similarities
 - weighted S^2 exploits the information given by the weights

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- sample complexity of weighted S^2

Hybrid approach: begin with global approach then switch to boundary refinement

- idea is to accelerate the convergence of label prediction using boundary refinement approach
- cutoff maximization [Anis, G., Ortega '14] \rightarrow weighted S^2

Motivation for Weighted S^2

- Weighted S^2 is a generalization of S^2 algorithm [Dasarathy, Nowak, Zhu '15]
- S^2 algorithm works on unweighted graphs
- · Finds cut edges by bisecting paths connecting two oppositely labeled nodes



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- In ML, node $i \Leftrightarrow \mathbf{x}_i \in \mathbb{R}^d$
- Weighted S^2 takes into account $l_{ij} = d(\mathbf{x}_i, \mathbf{x}_j)$ for edge (i, j)
- $d(\mathbf{x}_i, \mathbf{x}_j)$ expected to be larger for cut edges than within class edges
- Bisection based on l_{ij} can find cut edges faster



Weighted S^2 Algorithm



Weighted S^2 Algorithm



Given G = (V, E), Lengths $l \colon E \to \mathbb{R}^+$

I. Random sample until two opposite labeled, connected nodes u, v are found



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- 2. Find the shortest path between u and v
- 3. Bisection search: Find the cut-edge by successively sampling the nodes closest to the midpoint of the path
- 4. Remove the cut-edge and repeat until all the cut-edges are found

Sample Complexity of Weighted S^2

Quantities to parametrize the complexity of cut induced by \boldsymbol{f}



- f partitions G into conn. comp.'s $\{V_i\}$
- Cut C into corresponding cut comp.'s $\{C_{ij}\}$
- eta pprox balanced-ness of $|V_i|$'s
- m = number of cut components
- $l_n = \max$. shortest path length
- $l_{\rm cut} = \min \, {\rm cut} \, {\rm edge} \, {\rm length}$
- + $l_\kappa pprox \max$ dist betn two cut edges in C_{ij}

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Theorem (Sample Complexity)

Weighted S^2 recovers f with prob. $>(1-\epsilon)$ if the sampling budget is at least

$$\underbrace{\frac{\log(1/(\beta\epsilon))}{\log(1/(1-\beta))}}_{A: \text{ random sampling phase}} + \underbrace{m \left[2\log_2\left(\frac{l_n}{l_{\text{cut}}}\right) \right] + (|\partial C| - m) \left[2\log_2\left(\frac{l_\kappa}{l_{\text{cut}}}\right) \right]}_{B: \text{ bisection search phase}}$$

Sample Complexity of Random Sampling Phase

$$A = \frac{\log(1/(\beta\epsilon))}{\log(1/(1-\beta))}$$

f partitions G into similarly labeled connected components $\{V_1, \ldots, V_p\}$

- First sample in each V_i is obtained by random sampling
- A = # samples needed to sample at least one node from each V_i^{I}



- $\beta \coloneqq \min_{1 \le i \le p} |V_i| / |V|$
- measures how balanced V_i 's are
- small $\beta \Rightarrow$ more samples
- less likely to sample from small component

¹[Dasarathy, Nowak, Zhu '15]

Sample Complexity of Bisection Search





Lemma (Bisection search on a path)

Bisection search on path of length l discovers a cut edge of length l_{cut} in no more than $\left[2\log_2\left(\frac{l}{l_{\text{cut}}}\right)\right]$ steps.



- · length of the path of interest is at least halved after two queries
- bisect until discovery of cut edge \sim path of interest has length $l_{
 m cut}$
- number of samples = number of bisections $\left[2 \log_2 \left(\frac{l}{l_{\text{rut}}}\right)\right]$
- more samples if l is large (longer path) and l_{cut} is small (short cut edge)

Sample Complexity of Bisection Search (contd.)

$$B = \underbrace{m \left[2 \log_2 \left(\frac{l_n}{l_{\text{cut}}} \right) \right]}_{B_1} + \underbrace{\left(|\partial C| - m \right) \left[2 \log_2 \left(\frac{l_\kappa}{l_{\text{cut}}} \right) \right]}_{B_2}$$

Question: How many bisection searches and on what path lengths?

- B_1 : To discover the first cut edge (with length $\geq l_{\rm cut}$) in each cut component bisect paths of length $\leq l_n$
- B_2 : To discover the remaining cut edges (with length $\geq l_{\rm cut}$) in each cut component bisect paths of length $\leq l_{\kappa}$

Sample Complexity of Bisection Search (contd.)

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Number of samples needed to recover f increases with

- number of boundary nodes $|\partial C|$ and number of cut components m
- graph diameter l_n and distance between cut edges l_κ
- shorter cut edges (i.e., small $l_{\rm cut}$)

Experiment Setup

<u>Graph construction</u>: data $\{\mathbf{x}_1, \ldots, \mathbf{x}_n\} \subset \mathbb{R}^d$ and distances $d(\mathbf{x}_i, \mathbf{x}_j)$

- G: unweighted, symmetric k-nn graph (with k = 4)
- G_d : same topology as G but edge-weights $w_{ij} = d(\mathbf{x}_i, \mathbf{x}_j)$
- G_s : same topology as G with $w_{ij} = sim(\mathbf{x}_i, \mathbf{x}_j) \dots (\uparrow d \Leftrightarrow sim \downarrow)$

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Sampling algorithms

- Weighted S^2 on G_d
- S^2 on G [Dasarathy, Nowak, Zhu '15]
- Cutoff maximization on $G_{\rm S}$ [Anis, G., Ortega '15]

Label prediction from observed samples

- soft labels \hat{f} using bandlimited interpolation [Narang et al. '13]
- threshold $\hat{\mathbf{f}}$ to get the final predictions

Synthetic Data: Advantage of Weighted S^2 over S^2



- + 900 points (red) with f=+1 on inner circle of mean radius 1 and var 0.05
- 100 points (blue) with f = -1 on outer circle of mean radius 1.1 and var 0.45
- + 4-nn graph using Euclidean distance in \mathbb{R}^2

n	<i>C</i>	$ \partial C $	$\frac{\text{mean}(l_{\text{cut}})}{\text{mean}(l_{\text{non-cut}})}$	Unweighted S^2	Weighted S^2	Cutoff
1000	129	160	4.0533	237	179.2	999
	+	+ +	+ + + +	·		



An illustration of advantage of weighted S^2 (2 samples) over unweighted S^2 (3 samples)

Real World Data: Samples for Exact Recovery

USPS: handwritten digits

- $\mathbf{x}_i \in \mathbb{R}^{256}$ 16 imes 16 image
- $sim(i,j) = exp\left(-\frac{\|\mathbf{x}_i \mathbf{x}_j\|^2}{2\sigma^2}\right)$

•
$$d(i,j) = \|\mathbf{x}_i - \mathbf{x}_j\|$$

Newsgroups: documents

- $\mathbf{x}_i \in \mathbb{R}^{3000}$ tf-idf of words

•
$$sim(i,j) = \frac{\mathbf{x}_i^{\top} \mathbf{x}_j}{\|\mathbf{x}_i\| \|\mathbf{x}_j\|}$$

•
$$d(i,j) = \sqrt{1 - \sin^2(i,j)}$$

Data	n	C	$ \partial C $	$\frac{\text{mean}(l_{\text{cut}})}{\text{mean}(l_{\text{non-cut}})}$	UW. <i>S</i> ²	W . <i>S</i> ²	Cutoff	Hybrid	$n_{\rm switch}$
7 v 9	400	154	180	1.1074	312.37	312.07	399	277	47
2 v 4	400	29	39	1.1183	49.13	48.37	394	76	38
ΒvΗ	400	255	235	1.0691	368.07	368.17	399	384	42

- Weights don't help much (since $l_{\mathrm{cut}} pprox l_{\mathrm{non-cut}}$)
- · Global approach (max cutoff) not good at recovering exact boundary
- But good at signal approximation with fewer samples

Real World Data: Error vs. Number of Samples





- Fewer samples \Rightarrow cutoff max.
- More samples \Rightarrow weighted S^2
- Hybrid: start with cutoff max. then switch to weighted S^2
- Switch at sample i

$$1 - \frac{\langle \hat{\mathbf{f}}_i, \hat{\mathbf{f}}_{i-1} \rangle}{\|\hat{\mathbf{f}}_i\|\|\hat{\mathbf{f}}_{i-1}\|} < \delta \ (\Rightarrow \hat{\mathbf{f}}_i = \hat{\mathbf{f}}_{i-1})$$

Conclusion

- Weighted S^2 algorithm: generalization of S^2 to weighted graphs
- Analysis of sample complexity
- Demonstration of advantage of weighted S^2 over unweighted S^2
- Active learning approach given sampling budget / error tolerance:
 - small budget / more error tolerance \Rightarrow global smoothness approach (e.g., cutoff maximization)
 - large budget / less error tolerance \Rightarrow boundary refinement approach (e.g., weighted $S^2)$
- Hybrid approach: best of both methods

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Appendix: Clustered-ness of the Cut

•
$$e_1, e_2 \in C$$
: $\delta(e_1, e_2) = d^{G-C}(x_1, x_2) + d^{G-C}(y_1, y_2) + \max\{l_{e_1}, l_{e_2}\}$



• $H_r(C, \mathcal{E})$: graph with nodes \leftrightarrow cut edges in G

for
$$e_1, e_2 \in C$$
: $\{e_1, e_2\} \in \mathcal{E}$ if and only if $\delta(e_1, e_2) \leq r$

• As r increases, number of connected components in H_r decreases l_{κ} = the smallest r for which H_r has m connected components

Larger $l_\kappa \Rightarrow$ need to bisect a longer path to get the next cut edge