

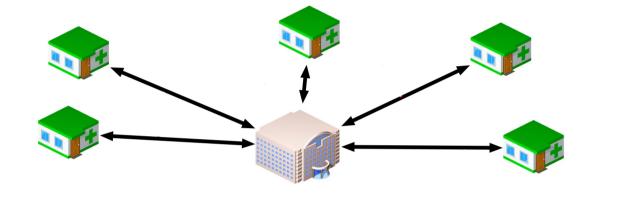


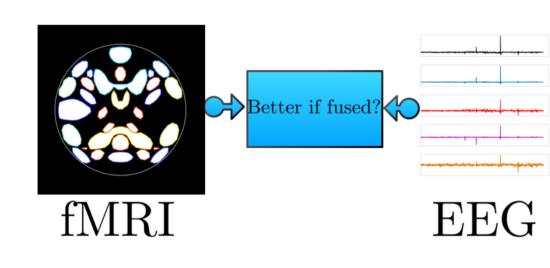
Motivation

Goal: measure linear relationship among variables \rightarrow can use correlation

Challenges: data – privacy-sensitive and distributed

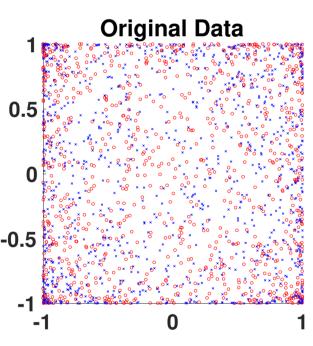
- \rightarrow how to guarantee privacy?
- \rightarrow how to measure the best correlation metric?
- \rightarrow how to do it in distributed setting?

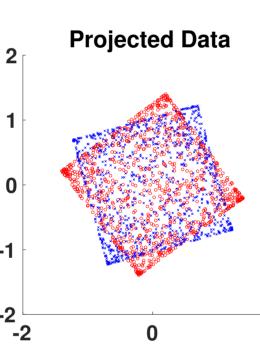




Canonical Correlation Analysis (CCA)

CCA finds subspaces for different "views" of data [1] \rightarrow "views" are maximally correlated after projection





 $= \mathbf{I},$

Can we have a CCA algorithm that preserves privacy, provides good utility and operates in distributed-data setting?

Problem Formulation

 \rightarrow consider a system with S different sites

- \rightarrow site s contains views: $\mathbf{X}_s \in \mathbb{R}^{D_x imes N_s}$, $\mathbf{Y}_s \in \mathbb{R}^{D_y imes N_s}$
- \rightarrow pooled data scenario: $\mathbf{X} = [\mathbf{X}_1 \dots \mathbf{X}_S] \in \mathbb{R}^{D_x \times N}$ and $\mathbf{Y} = [\mathbf{Y}_1 \dots \mathbf{Y}_S] \in \mathbb{R}^{D_y \times N}$
- \rightarrow goal: find subspaces $\mathbf{U} \in \mathbb{R}^{D_x imes K}$, $\mathbf{V} \in \mathbb{R}^{D_y imes K}$ [3]

$$\begin{array}{ll} \mbox{minimize} & \|\mathbf{U}^{\top}\mathbf{X} - \mathbf{V}^{\top}\mathbf{Y}\|_{F}^{2} \\ \mbox{ubject to} & \frac{1}{N}\mathbf{U}^{\top}\mathbf{X}\mathbf{X}^{\top}\mathbf{U} = \mathbf{I}, \frac{1}{N}\mathbf{V}^{\top}\mathbf{Y}\mathbf{Y}^{\top}\mathbf{V} \end{array}$$

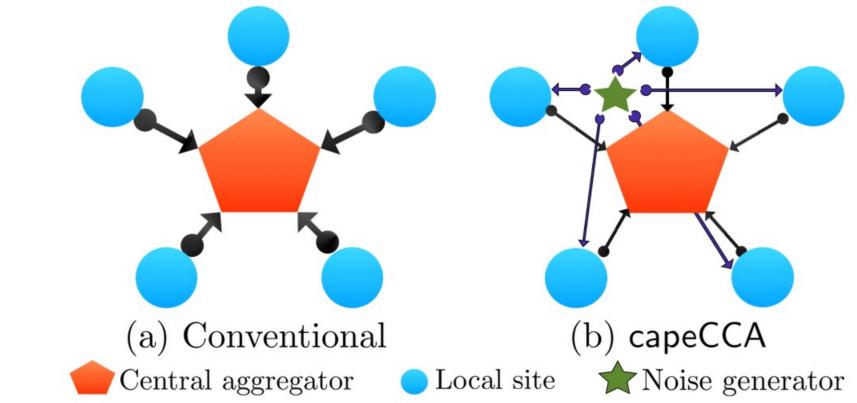
 $\frac{1}{N} \mathbf{U}^{\top} \mathbf{X} \mathbf{Y}^{\top} \mathbf{V} = \mathbf{I}.$

Want to estimate U and V in the distributed setting while preserving privacy

DISTRIBUTED DIFFERENTIALLY-PRIVATE CANONICAL CORRELATION ANALYSIS Hafiz Imtiaz and Anand D. Sarwate

Rutgers University **Privacy Analysis** • Analyze Gauss (AG) algorithm: input perturbation on 2nd-moment matrix [2] • DP is post-processing invariant \Rightarrow computation of U and V is (ϵ, δ) -DP • Projection/clustering do not satisfy DP \Rightarrow can be modified at the cost of utility **Simulation Results** (b) capeCCA • U of Wisc. X-ray Microbeam (XRMB) Dataset \rightarrow view 1: speech; view 2: jaw movement • fMRI+EEG Dataset \rightarrow view 1: fMRI; view 2: EEG • Clustering performance on XRMB \rightarrow CHIndex • Estimation of correlation on fMRI+EEG $\rightarrow err_{corr}$ XRMB (p = 30, ϵ = 0.2) Performance Variation on XRMB Dataset ---------1e⁻⁵ 1e⁻⁴ 1e⁻³ 1e⁻² 5e⁻² a) Privacy param (ϵ fMRI+EEG ($\epsilon = 0.5$) IRI+EEG (N = 1000) RI+EEG (N = 2000, ϵ = 0.5 Performance 0.15 -o-non-priv Variation on ^{ິວ} 0.1 ---local 0.1 fMRI + -+-capeCCA **0 - 00 - 00 - 00 - 0** EEG 1e⁻³ 1e⁻² 1e⁻¹ $1e^{-5}$ $1e^{-4}$ $1e^{-3}$ $1e^{-2}$ $5e^{-2}$ Dataset b) Total samples (N) c) Privacy param (δ a) Privacy param (**Conclusion and Future Works** capeCCA achieves the same utility as pooled-data scenario in the honest-but-curious setting **Takeaway: Future directions:** capeCCA has better utility than local and the same privacy level • capeCCA can reach non - priv in some reg • for fixed ϵ : more samples \rightarrow better perform • for fixed N and S: higher $\epsilon \rightarrow$ better per

Differential Privacy (DP)



Definition: Algorithm $\mathcal{A}(\mathbb{D})$ taking values in a set \mathbb{T} provides (ϵ, δ) -differential privacy [2] if

 $P(\mathcal{A}(\mathbb{D}) \in \mathbb{S}) \leq e^{\epsilon} P(\mathcal{A}(\mathbb{D}') \in \mathbb{S}) + \delta$ for all measurable $\mathbb{S} \subseteq \mathbb{T}$ and all *neighboring* data sets \mathbb{D} and \mathbb{D}' differing in a single entry. A conventional scheme:

- Compute $\mathbf{Z}_s = \begin{bmatrix} \mathbf{X}_s^\top & \mathbf{Y}_s^\top \end{bmatrix}^\top$ and $\mathbf{C}_s = \frac{1}{N_s} \mathbf{Z}_s \mathbf{Z}_s^\top$
- Send $\hat{\mathbf{C}}_s = \mathbf{C}_s + \mathbf{E}_s$ to aggregator, where $\{ [\mathbf{E}_s]_{ij} : i \in [D], j \leq i \}$ drawn i.i.d. from $\mathcal{N}(0, \tau_s^2)$
- Aggregator computes $\hat{\mathbf{C}} = \frac{1}{S} \sum_{s=1}^{S} \hat{\mathbf{C}}_{s}$
- Variance of the estimator: $\tau_{ag}^2 \triangleq \frac{\tau_s^2}{S}$
- \rightarrow In pooled-data setting: noise variance $\tau_c^2 = \frac{\tau_s^2}{S^2}$
- How can we achieve the same noise variance in the distributed setting? \rightarrow employ CAPE protocol [4]

Proposed Algorithm: capeCCA

Input: 0-centered samples \mathbf{X}_s and \mathbf{Y}_s as $\mathbf{Z}_s = \begin{bmatrix} \mathbf{X}_s^\top & \mathbf{Y}_s^\top \end{bmatrix}^{\prime}$, $\|\mathbf{z}_{s,n}\|_2 \leq 1$ for $s \in [S]$; privacy parameters ϵ, δ

Stage 1: Generate $\mathbf{E}_s \in \mathbb{R}^{D \times D}$

Stage 1: Generate $D \times D$ symmetric \mathbf{G}_s

Stage 2: Compute and send $\hat{\mathbf{C}}_s \leftarrow \frac{1}{N_s} \mathbf{Z}_s \mathbf{Z}_s^\top + \mathbf{E}_s + \mathbf{F}_s + \mathbf{G}_s$

Stage 1: Generate $\mathbf{F}_s \in \mathbb{R}^{D \times D}$

Stage 3: Compute $\hat{\mathbf{C}} \leftarrow \frac{1}{S} \sum_{s=1}^{S} \left(\hat{\mathbf{C}}_s - \mathbf{F}_s \right)$

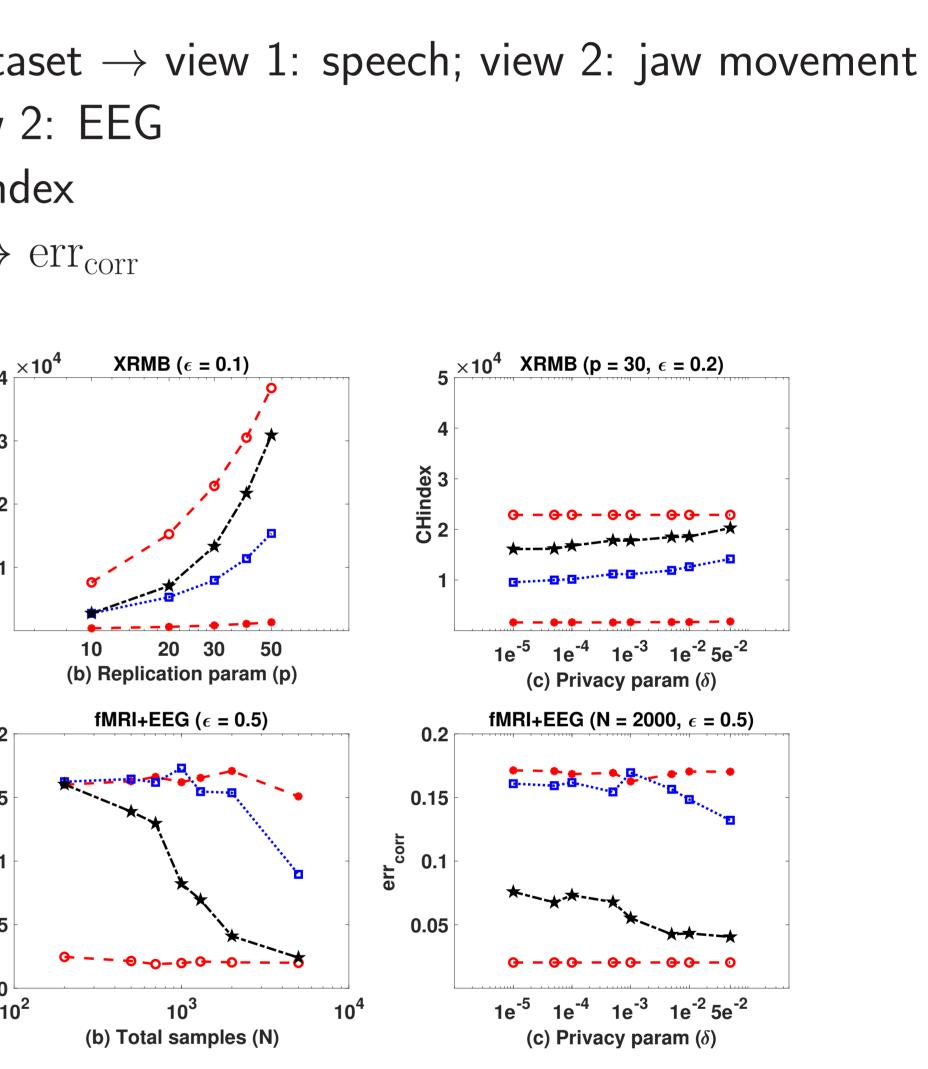
Stage 4: Extract sub-matrices from $\hat{\mathbf{C}}$

Output: Differentially-private approximates: $\hat{\mathbf{U}}^*$ and $\hat{\mathbf{V}}^*$

References

[1] Hotelling, H. (1936). Relations Between Two Sets of Variates. Biometrika, 28(3/4), 321-377. doi:10.2307/2333955 [2] Dwork, C. et al. (2014). Analyze Gauss: Optimal Bounds for Privacy-preserving Principal Component Analysis. doi: 10.1145/2591796.2591883 [3] Hardoon, D. R. et al. (2004). Canonical Correlation Analysis: An Overview with Application to Learning Methods. doi: 10.1162/0899766042321814 [4] Imtiaz, H. et al. (2019). Distributed Differentially Private Computation of Functions with Correlated Noise. arXiv e-print: http://arxiv.org/abs/1904.10059





l conv for	can we scrap the "trusted"
	noise generator? [4]
egimes	can we achieve the same in
rmance	an asymmetric network? [4]
rformance	can we achieve adapt our
	approach to $\delta = 0$?