Fast Computations for Approximation and Compression in Slepian Spaces

Santhosh Karnik¹, Zhihui Zhu², Michael B. Wakin², Justin Romberg¹, Mark A. Davenport¹

> ¹Georgia Institute of Technology ²Colorado School of Mines

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- Spectral Leakage
- Slepian Basis
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Suppose we take N equispaced samples of a bandlimited signal.

$$X_c(t) = \int_{-W/T}^{W/T} X_c(F) e^{j2\pi Ft} dF, \ t \in \mathbb{R}$$

$$x[n] = x_c(nT), n = 0, 1, \dots, N-1$$

If $W = \frac{1}{2}$, we are sampling at the Nyquist rate. We expect N degrees of freedom in x.

If $0 < W < \frac{1}{2}$, we are sampling faster than the Nyquist rate. We expect $\approx 2NW$ degrees of freedom in x.

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What basis yields a sparse representation for x?

Spectral Leakage



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$$\mathbf{x}[n] = \cos\left(\frac{2\pi \cdot 6}{100}n\right), \quad n = 0, 1, \dots, 99$$

Spectral Leakage



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$$\mathbf{x}[n] = \cos\left(\frac{2\pi \cdot 5.5}{100}n\right), \quad n = 0, 1, \dots, 99$$

Motivation

The relationship between a discrete signal $x \in \ell_2(\mathbb{Z})$ and its discrete time Fourier transform (DTFT) $\tilde{x} \in L_2([-\frac{1}{2}, \frac{1}{2}])$ is given by

$$\widetilde{x}(f) = \sum_{n=-\infty}^{\infty} x[n]e^{-j2\pi fn},$$

$$x[n] = \int_{-1/2}^{1/2} \widetilde{x}(f) e^{j2\pi fn} df.$$

For a given positive integer N and bandlimit $0 < W < \frac{1}{2}$, the only signal which is both bandlimited to $f \in [-W, W]$ and timelimited to $n \in \{0, 1, ..., N-1\}$ is the zero signal.

Can we find timelimited signals whose energy is maximally concentrated in the frequency interval [-W, W]?

Definition

For any length N signal, $\mathbf{x} \in \mathbb{C}^N$,

$$\frac{\int_{-W}^{W} |\tilde{\boldsymbol{x}}(f)|^2 df}{\int_{-1/2}^{1/2} |\tilde{\boldsymbol{x}}(f)|^2 df} = \frac{\boldsymbol{x}^* \boldsymbol{B}_{N,W} \boldsymbol{x}}{\boldsymbol{x}^* \boldsymbol{x}},$$

where $\boldsymbol{B}_{N,W}$ is an $N \times N$ matrix with entries

$$\boldsymbol{B}_{N,W}[m,n] = \begin{cases} \frac{\sin(2\pi W(m-n))}{\pi(m-n)} & \text{if } m \neq n \\ 2W & \text{if } m = n \end{cases}$$

The Slepian basis vectors $\boldsymbol{s}_{N,W}^{(0)}, \boldsymbol{s}_{N,W}^{(1)}, \dots, \boldsymbol{s}_{N,W}^{(N-1)}$ are defined as the eigenvectors of $\boldsymbol{B}_{N,W}$, where the respective eigenvalues $\lambda_{N,W}^{(0)}, \lambda_{N,W}^{(1)}, \dots, \lambda_{N,W}^{(N-1)}$ are sorted in decreasing order [1].

Slepian Basis Vector Plots

Signal Length: N = 128, Bandlimit: $W = \frac{1}{8}$

Santhosh Karnik¹, Zhihui Zhu², Michael B. Wakin², Justin Rom Fast Slepian Methods

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Eigenvalue Behavior

Slightly less than 2*NW* eigenvalues are close to 1, and slightly less than N - 2NW eigenvalues are close to 0. For a fixed $\epsilon > 0$, roughly $O(\log N)$ eigenvalues are between ϵ and $1 - \epsilon$.



Slepian Transform





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Slepian Transform





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Goals

- A fast method to compress N dimensional data into $K \approx 2NW$ dimensional data, and still be able to recover the original data if it lies in the span of the first K Slepian basis vectors.
- A fast method to project a vector in \mathbb{C}^N onto the span of the first $K \approx 2NW$ Slepian basis vectors.
- A fast method to apply the rank $K \approx 2NW$ pseudoinverse of $B_{N,W}$ to a vector in \mathbb{C}^N .

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Prolate Matrix vs. Partial Fourier Matrix

Let $F_{N,W}$ be an $N \times 2NW'$ matrix whose columns are the lowest 2NW' frequency DFT vectors of length N.

Theorem 1

For any $N \in \mathbb{N}$, $W \in (0, \frac{1}{2})$, and $\epsilon \in (0, \frac{1}{2})$, there exist $N \times R_L$ matrices L_1, L_2 and an $N \times N$ matrix E_F such that

$$oldsymbol{B}_{N,W}=oldsymbol{F}_{N,W}oldsymbol{F}_{N,W}^*+oldsymbol{L}_1oldsymbol{L}_2^*+oldsymbol{E}_F$$

where

$$R_L \leq \left(rac{4}{\pi^2}\log(8N) + 6
ight)\log\left(rac{15}{\epsilon}
ight) ext{ and } \|m{E}_F\| \leq \epsilon.$$

Due to the length limitations, we have included the proof in Section 2 of "The Fast Slepian Transform". A preprint available at https://arxiv.org/abs/1611.04950.

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Number of Eigenvalues in Transition Region

Slepian [1] showed that for a fixed W and ϵ ,

$$\#\{\ell:\epsilon\leq\lambda_{N,W}^{(\ell)}\leq 1-\epsilon\}\sim \frac{2}{\pi^2}\log N\log\left(\frac{1}{\epsilon}-1\right).$$

Zhu and Wakin [2] showed that for any N, W, and ϵ ,

$$\#\{\ell:\epsilon \leq \lambda_{N,W}^{(\ell)} \leq 1-\epsilon\} \leq \frac{\frac{2}{\pi^2}\log(N-1) + \frac{2}{\pi^2}\frac{2N-1}{N-1}}{\epsilon(1-\epsilon)}$$

Using the previous theorem, and the Courant-Fischer-Weyl min-max principle, we can improve this bound.

Theorem 2

For any
$$N \in \mathbb{N}$$
, $W \in (0, \frac{1}{2})$, and $\epsilon \in (0, \frac{1}{2})$,

$$\#\{\ell:\epsilon<\lambda_{\textit{\textit{N}},\textit{\textit{W}}}^{(\ell)}<1-\epsilon\}\leq \left(\frac{8}{\pi^2}\log(8\textit{\textit{N}})+12\right)\log\left(\frac{15}{\epsilon}\right)$$

Fast Projection

Define
$$\boldsymbol{S}_{[K]} = \begin{bmatrix} \boldsymbol{s}_{N,W}^{(0)} & \boldsymbol{s}_{N,W}^{(1)} & \cdots & \boldsymbol{s}_{N,W}^{(K-1)} \end{bmatrix}$$
.

Theorem 3

For any $N \in \mathbb{N}$, $W \in (0, \frac{1}{2})$, and $\epsilon \in (0, \frac{1}{2})$, pick K such that $\lambda_{N,W}^{(K-1)} > \epsilon$ and $\lambda_{N,W}^{(K)} < 1 - \epsilon$. Then, there exist $N \times R_U$ matrices U_1, U_2 and an $N \times N$ matrix \mathbf{E}_S such that

$$\boldsymbol{S}_{[K]}\boldsymbol{S}_{[K]}^* = \boldsymbol{B}_{N,W} + \boldsymbol{U}_1\boldsymbol{U}_2^* + \boldsymbol{E}_S$$

where

$${\mathcal R}_U \leq \left(rac{8}{\pi^2}\log(8{\mathcal N}) + 12
ight)\log\left(rac{15}{\epsilon}
ight) \; {\it and} \; \|{m E}_{\mathcal S}\| \leq \epsilon.$$

Idea: $\boldsymbol{B}_{N,W}$ and $\boldsymbol{S}_{[K]}\boldsymbol{S}_{[K]}^*$ have the same eigenvectors, and most of the eigenvalues are almost the same.

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Fast Projection

Define

$$\begin{split} \mathcal{I}_1 &= \{\ell < \mathcal{K} : 1 - \epsilon \leq \lambda_{\mathcal{N},\mathcal{W}}^{(\ell)} \leq 1\}, \\ \mathcal{I}_2 &= \{\ell < \mathcal{K} : \epsilon < \lambda_{\mathcal{N},\mathcal{W}}^{(\ell)} < 1 - \epsilon\}, \\ \mathcal{I}_3 &= \{\ell \geq \mathcal{K} : \epsilon < \lambda_{\mathcal{N},\mathcal{W}}^{(\ell)} < 1 - \epsilon\}, \\ \mathcal{I}_4 &= \{\ell \geq \mathcal{K} : 0 \leq \lambda_{\mathcal{N},\mathcal{W}}^{(\ell)} \leq \epsilon\}, \end{split}$$

and accordingly partition

$$\begin{split} \boldsymbol{S}_{N,W} &:= \begin{bmatrix} \boldsymbol{s}_{N,W}^{(0)} & \cdots & \boldsymbol{s}_{N,W}^{(N-1)} \end{bmatrix} = \begin{bmatrix} \boldsymbol{S}_1 & \boldsymbol{S}_2 & \boldsymbol{S}_3 & \boldsymbol{S}_4 \end{bmatrix} \\ \boldsymbol{\Lambda}_{N,W} &:= \operatorname{diag}(\lambda_{N,W}^{(0)}, \dots, \lambda_{N,W}^{(N-1)}) = \operatorname{blockdiag}(\boldsymbol{\Lambda}_1, \boldsymbol{\Lambda}_2, \boldsymbol{\Lambda}_3, \boldsymbol{\Lambda}_4). \end{split}$$

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Fast Projection

So,
$$B_{N,W} = S_1 \Lambda_1 S_1^* + S_2 \Lambda_2 S_2^* + S_3 \Lambda_3 S_3^* + S_4 \Lambda_4 S_4^*$$
,
and $S_{[K]} S_{[K]}^* = S_1 S_1^* + S_2 S_2^*$.

Then, if we define

$$\begin{split} \boldsymbol{U}_1 &= \begin{bmatrix} \boldsymbol{S}_2 (\boldsymbol{I} - \boldsymbol{\Lambda}_2)^{1/2} & \boldsymbol{S}_3 \boldsymbol{\Lambda}_3^{1/2} \end{bmatrix}, \\ \boldsymbol{U}_2 &= \begin{bmatrix} \boldsymbol{S}_2 (\boldsymbol{I} - \boldsymbol{\Lambda}_2)^{1/2} & -\boldsymbol{S}_3 \boldsymbol{\Lambda}_3^{1/2} \end{bmatrix}, \\ \boldsymbol{E}_5 &= \boldsymbol{S}_1 (\boldsymbol{I} - \boldsymbol{\Lambda}_1) \boldsymbol{S}_1^* - \boldsymbol{S}_4 \boldsymbol{\Lambda}_4 \boldsymbol{S}_4^*, \end{split}$$

we have

$$\boldsymbol{S}_{[\mathcal{K}]}\boldsymbol{S}_{[\mathcal{K}]}^* = \underbrace{\boldsymbol{B}_{\mathcal{N},\mathcal{W}}}_{\text{Toeplitz}} + \underbrace{\boldsymbol{U}_1\boldsymbol{U}_2^*}_{\text{Low Rank}} + \underbrace{\boldsymbol{E}_S}_{\text{Small}}.$$

So, $\boldsymbol{B}_{N,W} + \boldsymbol{U}_1 \boldsymbol{U}_2^*$ is an approximation for $\boldsymbol{S}_{[K]} \boldsymbol{S}_{[K]}^*$ which can be applied to a vector in $O(N \log N \log \frac{1}{\epsilon})$ operations.

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Fast Compression

Theorem 4

For any $N \in \mathbb{N}$, $W \in (0, \frac{1}{2})$, and $\epsilon \in (0, \frac{1}{2})$, pick K such that $\lambda_{N,W}^{(K-1)} > \epsilon$ and $\lambda_{N,W}^{(K)} < 1 - \epsilon$. Then, there exist $N \times K'$ matrices T_1, T_2 where

$$\mathcal{K}' \leq \lceil 2\mathcal{N}\mathcal{W}
ceil + \left(rac{12}{\pi^2}\log(8\mathcal{N}) + 18
ight)\log\left(rac{15}{\epsilon}
ight)$$

such that $\|\boldsymbol{T}_1\boldsymbol{T}_2^* - \boldsymbol{S}_{[K]}\boldsymbol{S}_{[K]}^*\| \le 2\epsilon$, and \boldsymbol{T}_1 and \boldsymbol{T}_2^* can be applied to a vector in $O(N \log N \log \frac{1}{\epsilon})$ operations.

Idea: Combine Theorem 1 and Theorem 3 to get

$$\boldsymbol{S}_{[K]}\boldsymbol{S}_{[K]}^* = \boldsymbol{F}_{N,W}\boldsymbol{F}_{N,W}^* + \boldsymbol{L}_1\boldsymbol{L}_2^* + \boldsymbol{U}_1\boldsymbol{U}_2^* + \boldsymbol{E}_F + \boldsymbol{E}_S.$$

Then, set $\boldsymbol{T}_1 = \begin{bmatrix} \boldsymbol{F}_{N,W} & \boldsymbol{L}_1 & \boldsymbol{U}_1 \end{bmatrix}$ and $\boldsymbol{T}_2 = \begin{bmatrix} \boldsymbol{F}_{N,W} & \boldsymbol{L}_2 & \boldsymbol{U}_2 \end{bmatrix}$.

Let $\boldsymbol{B}_{N,W}^{\dagger}$ be the rank-K truncated pseudoinverse of $\boldsymbol{B}_{N,W}$.

Theorem 5

For any $N \in \mathbb{N}$, $W \in (0, \frac{1}{2})$, and $\epsilon \in (0, \frac{1}{2})$, pick K such that $\lambda_{N,W}^{(K-1)} > \epsilon$ and $\lambda_{N,W}^{(K)} < 1 - \epsilon$. Then, there exist $N \times R_V$ matrices V_1, V_2 and an $N \times N$ matrix \mathbf{E}_B such that

$$\boldsymbol{B}_{N,W}^{\dagger} = \boldsymbol{B}_{N,W} + \boldsymbol{V}_1 \boldsymbol{V}_2^* + \boldsymbol{E}_B$$

where

$${\sf R}_V \leq \left(rac{8}{\pi^2}\log(8{\sf N})+12
ight)\log\left(rac{15}{\epsilon}
ight) \; {\it and} \; \|{m E}_B\| \leq 3\epsilon.$$

Idea: $\boldsymbol{B}_{N,W}^{\dagger}$ and $\boldsymbol{B}_{N,W}$ have the same eigenvectors, and most of the eigenvalues are almost the same.

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Fast Pseudoinverse

So,
$$B_{N,W} = S_1 \Lambda_1 S_1^* + S_2 \Lambda_2 S_2^* + S_3 \Lambda_3 S_3^* + S_4 \Lambda_4 S_4^*$$
,
and $B_{N,W}^{\dagger} = S_1 \Lambda_1^{-1} S_1^* + S_2 \Lambda_2^{-1} S_2^*$.

Then, if we set

$$\begin{split} \boldsymbol{V}_1 &= \begin{bmatrix} \boldsymbol{S}_2 (\boldsymbol{\Lambda}_2^{-1} - \boldsymbol{\Lambda}_2)^{1/2} & \boldsymbol{S}_3 \boldsymbol{\Lambda}_3^{1/2} \end{bmatrix}, \\ \boldsymbol{V}_2 &= \begin{bmatrix} \boldsymbol{S}_2 (\boldsymbol{\Lambda}_2^{-1} - \boldsymbol{\Lambda}_2)^{1/2} & -\boldsymbol{S}_3 \boldsymbol{\Lambda}_3^{1/2} \end{bmatrix}, \\ \boldsymbol{E}_B &= \boldsymbol{S}_1 (\boldsymbol{\Lambda}_1^{-1} - \boldsymbol{\Lambda}_1) \boldsymbol{S}_1^* - \boldsymbol{S}_4 \boldsymbol{\Lambda}_4 \boldsymbol{S}_4^*, \end{split}$$

we have

$$\boldsymbol{B}_{N,W}^{\dagger} = \underbrace{\boldsymbol{B}_{N,W}}_{\text{Toeplitz}} + \underbrace{\boldsymbol{V}_{1}\boldsymbol{V}_{2}^{*}}_{\text{Low Rank}} + \underbrace{\boldsymbol{E}_{B}}_{\text{Small}}.$$

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So, $\boldsymbol{B}_{N,W} + \boldsymbol{V}_1 \boldsymbol{V}_2^*$ is an approximation for $\boldsymbol{B}_{N,W}^{\dagger}$ which can be applied to a vector in $O(N \log N \log \frac{1}{\epsilon})$ operations.

Simulations



Simulations



- The Slepian basis is a good choice of basis for working with samples of a bandlimited signal.
- We now have fast (approximate) algorithms for working with the Slepian basis, whose complexity scales comparably to the FFT.

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