

Rumor Source Detection: A Probabilistic Perspective

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Rumor Source Detection Problem

At time 0, start spreading.



At unknown time t, observe G_I .

SI Model Spreading

Susceptible and Infected. Takes time $au_{ij} \sim P_{ij}$ for *i* to infect *j*.



Belief Propagation Algorithm

- $\blacktriangleright m_{i \to j}(t)$: probability that if j is the source, j causes the infection in the subtree rooted at i.
- ▶ $f_{ij}(t)$: PDF of the time it takes for node *i* to infect node *j*.

 $m_{C \to D} = (m_{A \to C} m_{B \to C}) * f_{DC}$

- ▶ Goal: Find *s* from G_I .
- ▶ It is natural to estimate *s* through the MLE:

 $\hat{s}_{\mathsf{MLE}} = \arg\max L(v|G_I), \ L(v|G_I) \triangleq \Pr(\exists t, \mathcal{G}_I^s(t) = G_I|s = v).$ $v \in \mathcal{V}$

Challenge and Previous Works

✓ Challenge:

- MLE's computation time is exponential in # of infected nodes [1].
- Previous works resort to approximations.
- Previous Approximation Schemes:
 - ▶ [1] Rumor Center \hat{s}_R :





 $m_{k \to D}$

* denotes convolution.

 $L(j,t|G_I) = m_{i \to j}$ $i \in Nei(j)$

- By Fast Fourier Transform, a convolution takes time $O(L \log L)$.
- Each edge needs to propagate two times (forward and backward).
- Since a tree has n 1 edges, the time complexity is $O(nL \log L)$.

Extension to General Graph (GGT Heuristic)

Generate a weighted spanning tree. Then apply the BP algo.

- 1. Given G_I and the rates on edges, Assume A is the source. $I = \{A\}$
- 4. Choose the next infected node again. Assume C is chosen. $I = \{A, B, C\}$ Update C's parent and infection rate.

 $R(v, G_I)$:# of ways infection starting from v leading to G_I .

 $\hat{s}_R = \arg \max R(v, G_I).$ $v \in \mathcal{V}_I$

▶ [2] Jordan Center \hat{s}_{J} :

d(u, v): minimum # of hops from node u to node v.

 $\hat{s}_J = \underset{u \in \mathcal{V}_I}{\operatorname{arg\,min\,}} \max_{v \in \mathcal{V}_I} d(u, v).$

▶ [3] Dynamic Message Passing \hat{s}_D :

 $\mathbb{P}_{v}(i,t)$: the estimate of the probability that when node v is the source,

node *i* is infected at time *t*.

 $\hat{s}_D = \underset{v \in \mathcal{V}_I}{\operatorname{arg\,max\,max}} \max_{t \in \mathbb{R}_+} \prod_{i \in \mathcal{V}_I} \tilde{\mathbb{P}}_v(i,t) \prod_{i \notin \mathcal{V}_I} (1 - \tilde{\mathbb{P}}_v(i,t)).$

Key Ideas

 $\mathbb{P}(G|s,t) = \mathbb{P}(\mathscr{T}_1|s,t)\mathbb{P}(\mathscr{T}_2|s,t).$

If G_I is composed of two trees, \mathscr{T}_1 and \mathscr{T}_2 , then the joint likelihood



Experimental Results



is factorized:

Recursively factorizing leads to the Belief Propagation Algorithm.

In Contrast, the likelihood of the node is not factorizable because

 $\mathbb{P}(G|s) = \int_{t} \mathbb{P}(\mathscr{T}_{1}|s,t) \mathbb{P}(\mathscr{T}_{2}|s,t) p(t) dt \neq \mathbb{P}(\mathscr{T}_{1}|s) \mathbb{P}(\mathscr{T}_{2}|s).$

Thus, we aim to find the *Joint Maximum Likelihood Estimator*:

 $(\hat{s}_{\mathsf{JMLE}}, \hat{t}_{\mathsf{JMLE}}) = \underset{(v,t)\in\mathcal{V}\times\mathbb{R}_+}{\arg\max} L(v,t|G_I), L(v,t|G_I) \triangleq \Pr(\mathcal{G}_I^s(t) = G_I|s = v,t).$

The more similar a graph is to a tree, the more powerful the BP

algorithm will be.

Reference:

[1] D. Shah and T. Zaman, "Rumors in a network: Who's the culprit?" *IEEE Transactions on Information Theory*, vol. 57, no. 8, pp. 5163–5181, Aug 2011. [2] W. Luo, W. P. Tay, M. Leng, and M. K. Guevara, "On the universality of the jordan center for estimating the rumor source in a social network," in 2015 *IEEE International Conference on Digital Signal Processing (DSP), July 2015, pp. 760–764.* [3] A. Y. Lokhov, M. M'ezard, H. Ohta, and L. Zdeborov'a, "Inferring the origin of an epidemic with a dynamic message-passing algorithm," Phys. Rev. E,

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