Distributed Scheduling Algorithms for Optimizing Information Freshness in Wireless Networks

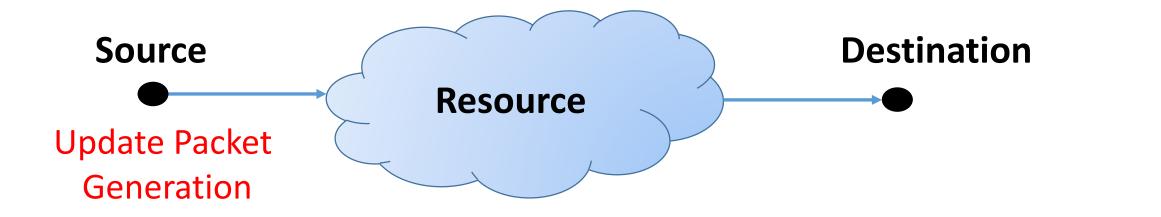
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Motivation: Information Freshness

Information Freshness Critical for Performance of UAV networks, vehicular networks, CPS

Utilizing the resource (e.g. queue, network)



Network Age Minimization

Lemma: Peak and average age are equal

$$A_e^{\rm ave} = A_e^{\rm p} = \frac{1}{\gamma_e f_e}$$

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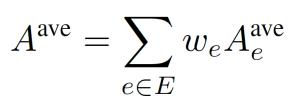
where $f_e = p_e \prod_{e' \in N_e} (1 - p_{e'})$ link activation frequency

Network age optimization problem

$$\underset{\mathbf{p}\in[0,1]^{|E|},\mathbf{f}\in\mathbb{R}^{|E|}}{\text{Minimize}} \quad \sum_{e\in E}\frac{w_e}{\gamma_e f_e}$$



SPAWC 2018



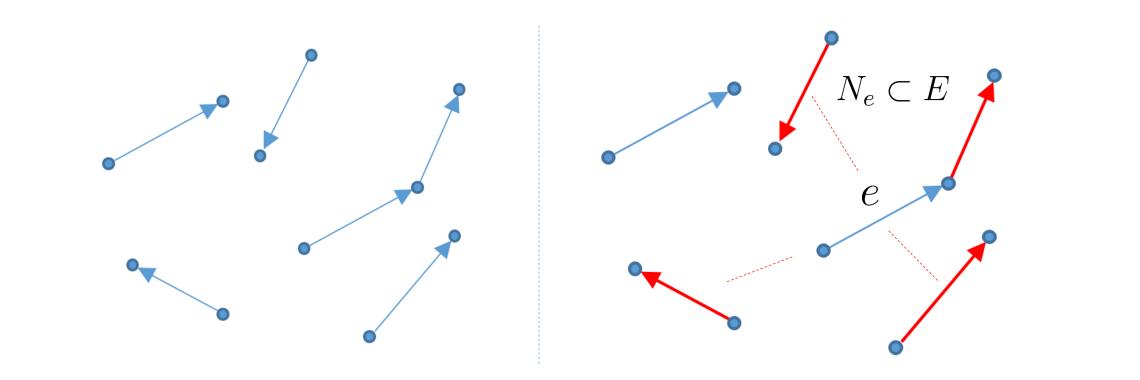
Packet generation rate too high → Congestion **Packet generation rate too low** \rightarrow Infrequent updates

Resource = Wireless Network with Interference **Contribution**: Distributed scheduling algorithm for age

System Model

minimization

Network G = (V, E) Directed Links **Source-destination pair** for every $e \in E$

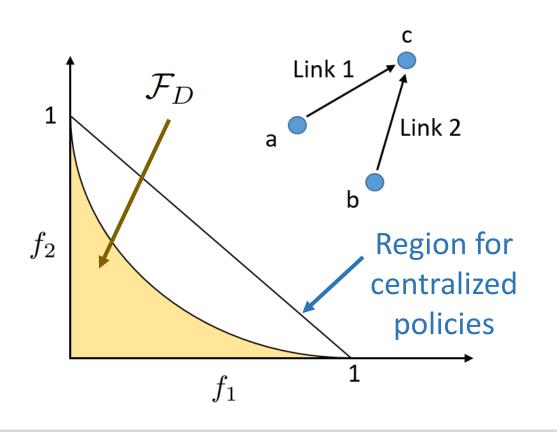


subject to $f_e = p_e \prod (1 - p_{e'}) \quad \forall e \in E$ $e' \in N_e$

Non-convex constraint set:

$$\mathcal{F}_D = \left\{ \mathbf{f} \in \mathbb{R}^{|E|} \mid f_e = p_e \prod_{e' \in N_e} (1 - p_{e'}) \text{ and } 0 \le p_e \le 1 \right\}$$

Theorem: $p_e^* = \frac{w_e A_e^*}{w_e A_e^* + \sum_{e' \in C} w_e' A_{e'}^*}$



Requires centralized computation

Distributed Computation

Theorem: Network age problem is equivalent to

$$\underset{\lambda_{e} \geq 0}{\text{Maximize}} \sum_{e \in E} \left(\lambda_{e} + \sum_{e': e \in N_{e'}} \lambda_{e'} \right) H\left(\frac{\lambda_{e}}{\lambda_{e} + \sum_{e': e \in N_{e'}} \lambda_{e'}} \right) + \sum_{e \in E} \lambda_{e} \left[1 + \log\left(\frac{w_{e}}{\lambda_{e}}\right) \right]$$

where $H(p) = p \log\left(\frac{1}{p}\right) + (1-p) \log\left(\frac{1}{1-p}\right)$ is the entropy function.

• λ_e is proxy for weighted age $w_e A_e^*$ $p_e^* = \frac{\lambda_e^*}{\lambda_e^* + \sum_{e':e \in N} \lambda_{e'}^*}$

Interference

- Interfering subset $N_e \subset E$ for each link e
- Popular *k*-hop interference model is a special case

Slotted time

Channel Process

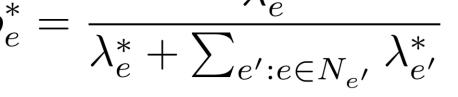
- Channel state for link $e: S_e(t) \in \{0, 1\}$ i.i.d. across t
- Independent across links

 $\gamma_e = \mathbb{P}\left[S_e(t) = 1
ight]$ Channel Statistic

time *i*

Distributed Stationary Scheduling Policy

Link e attempts transmission with probability p_e



Projected gradient descent has a distributed implementation

• We use the proxy for weighted age λ_e and iterate over time frames m

• Proxy for age:
$$\lambda_e(m)$$
 and $\theta_e(m) = \sum_{e' \in N_e} \lambda_{e'}(m)$

- Attempt probability: $p_e(m) = \frac{\lambda_e(m)}{\lambda_e(m) + \theta_e(m)}$
- Symmetric Interference Case
 - $N_e = \{ e' \in E \mid e \in N_{e'} \}$

• Local Updates: $\lambda_e(m+1) \leftarrow \Pi_\epsilon \left\{ \lambda_e(m) + \eta_m \left\{ \log \left(\frac{w_e}{\lambda_e(m)} \right) + \log \left(1 + \frac{\theta_e(m)}{\lambda_e(m)} \right) + \sum_{e' \in \mathcal{W}} \log \left(1 + \frac{\lambda_{e'}(m)}{\theta_{e'}(m)} \right) \right\} \right\}$

Links exchange $\lambda_e(m)$ and $\theta_e(m)$

Active Sources and Age Evolution

Active source: generates fresh updates for every Tx **Buffered source**: can only transmit queued up updates

Studied in [Talak et al. Mobihoc'18]

Our Recent Work

Age can be improved by centralized scheduling algorithms. Centralized policies and buffered sources studied in [Talak et al. Mobihoc 2018]

