

Performance of the Asynchronous Consensus Based Bundle Algorithm in Lossy Network Environments

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The Intersection of Planning and Communications

- Without reliable wireless communications, drones cannot coordinate
- Consequences of network losses
 - Failed delivery of sensed data to processing nodes
 - Insufficient situational awareness for effective in-field planning
 - Delayed/lost command and control messages (focus of this work)
 - Worst case: mission failure!
- Despite these adverse effects, most planning literature assumes perfect communication among nearby agents
- This has led to various techniques to maintain network connectivity



Related Work

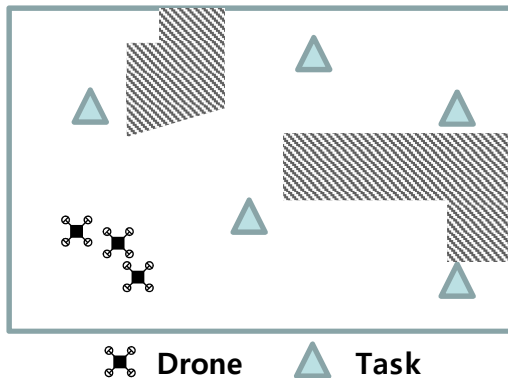
- "Binary" connectivity
 - Connectivity-as-a-service [Cornejo, '09]
 - Refine arbitrary motion plan to preserve network connectivity and meet goals
 - Control-theoretic connectivity [Zavlanos, '11]
 - Convex optimization and subgradient descent algorithms to maximize network's algebraic connectivity
 - Potential fields to control network topology
 - Connectivity-aware task allocation [Ponda, '12]
 - Extend the well-known Consensus Based Bundle Algorithm (CBBA [Choi, '09]) to include planning for relays
- Connectivity with variable reliability
 - BER- and throughput-aware task allocation [Kopeikan, '12]
 - Extend CBBA with relays to meet BER and throughput constraints

Prior work investigates how planning affects communication, but not how *unreliable communication affects planning*



The Task Allocation Problem

- **Given**
 - A set of drones
 - A set of tasks
- **Goal**
 - Allocate tasks to drones (at most one drone per task)
 - Maximize sum utility



Example environment (with obstacles)

We investigate the effect of realistic network environments on the Asynchronous CBBA (ACBBA [Johnson, '10], [Johnson, '11])



Problem Formulation: Notation

- $A = \{1, 2, \dots, N_a\}$: Set of agents (drones)
 - $i \in A$: Specific agent
- $T = \{1, 2, \dots, N_t\}$: Set of tasks
 - $j \in T$: Specific task
- $\mathbf{x}_i = (x_{i1}, x_{i2}, \dots, x_{iN_t})$: agent i 's assignment vector
 - $x_{ij} = 1$ if agent $i \in A$ is assigned task $j \in T$
 - $x_{ij} = 0$, otherwise
- \mathbf{p}_i : ordered sequence of tasks assigned to agent i
- $u_{ij}(\tau_{ij}(\mathbf{p}_i))$: agent i 's utility for completing task j at time $\tau_{ij}(\mathbf{p}_i)$
 - $u_{ij}(\tau_{ij}(\mathbf{p}_i)) = r_j \lambda^{\tau_{ij}(\mathbf{p}_i)}$, where r_j is reward for task j and $\lambda \in [0, 1)$



Problem Formulation: Optimization

$$\max \sum_{i \in A} \left(\sum_{j \in T} u_{ij}(\tau_{ij}(\mathbf{p}_i)) x_{ij} \right)$$

$$\text{subject to } \sum_{j \in T} x_{ij} \leq L_t, \quad \forall i \in A$$

$$\sum_{i \in A} x_{ij} \leq 1, \quad \forall j \in T$$

$$x_{ij} \in \{0,1\}, \quad \forall (i,j) \in A \times T$$

If agents form a *connected network* and there are *no transmission errors*, then the CBBA guarantees a non-conflicting task assignment. This assignment achieves within 50% of the optimal utility.



Internal State Information in the CBBA

Each agent $i \in A$ maintains the following five internal state vectors

- **Bundle vector \mathbf{b}_i**
 - Element $b_{in} \in T$ corresponds to the n th task assigned to agent i
 - Tasks are ordered based on when they are "won"
- **Path vector \mathbf{p}_i**
 - Contains same tasks as bundle, but ordered based on when they will be completed
- **Winning agent vector \mathbf{z}_i**
 - Element $z_{ij} \in A$ indicates who agent i believes has highest bid for task j
- **Winning bid vector \mathbf{y}_i**
 - Element $y_{ij} \in R_+$ corresponds to agent z_{ij} 's winning bid for task j
- **Timestamp vector \mathbf{t}_i**
 - Element $t_{ij} \in R_+$ indicates when agent z_{ij} placed bid y_{ij} on task j

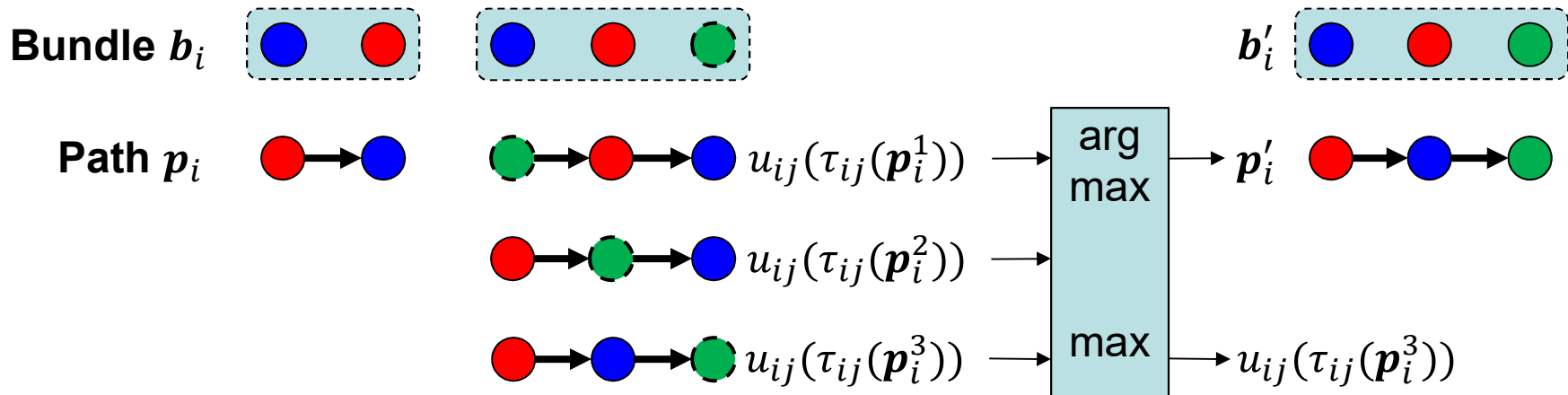


CBBA Iterations

The CBBA iterates among three phases

- **Bundle construction phase**

- Each agent adds tasks to its bundle in a *sequential greedy* fashion



- **State exchange phase**

- Each agent communicates its winning agent vector z_i , winner bid vector y_i , and timestamp vector t_i

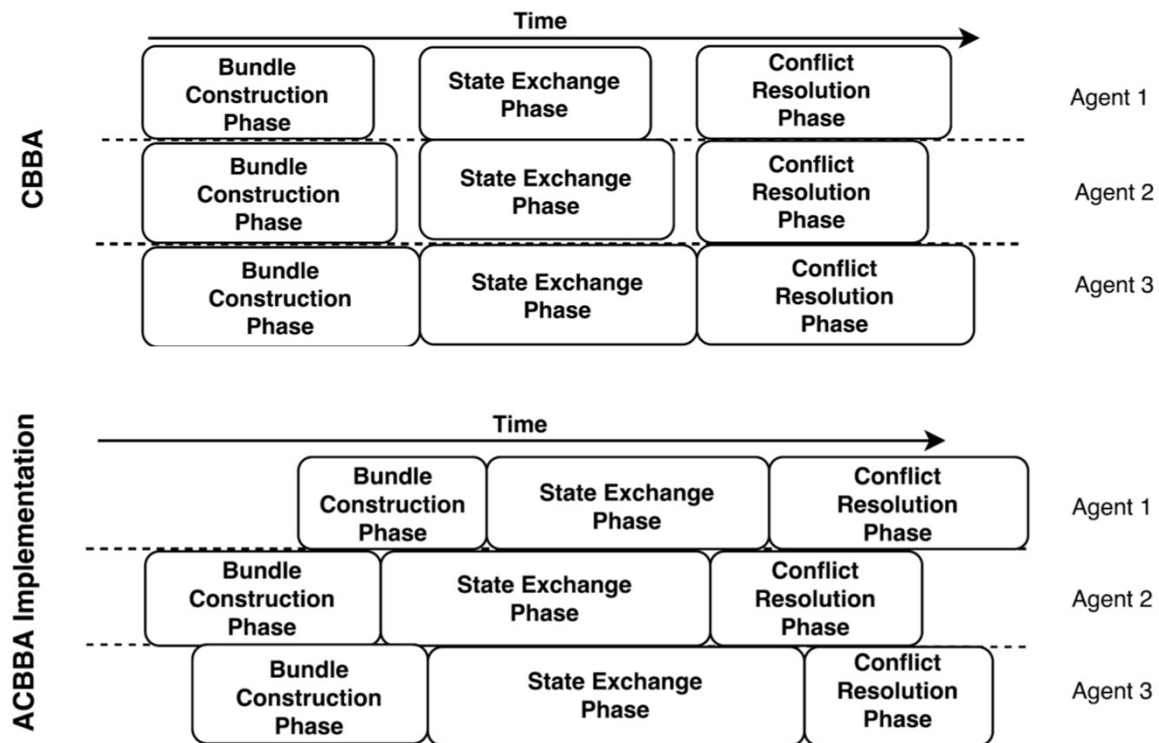
- **Conflict resolution phase**

- Each agent releases tasks it was outbid on and tasks added thereafter



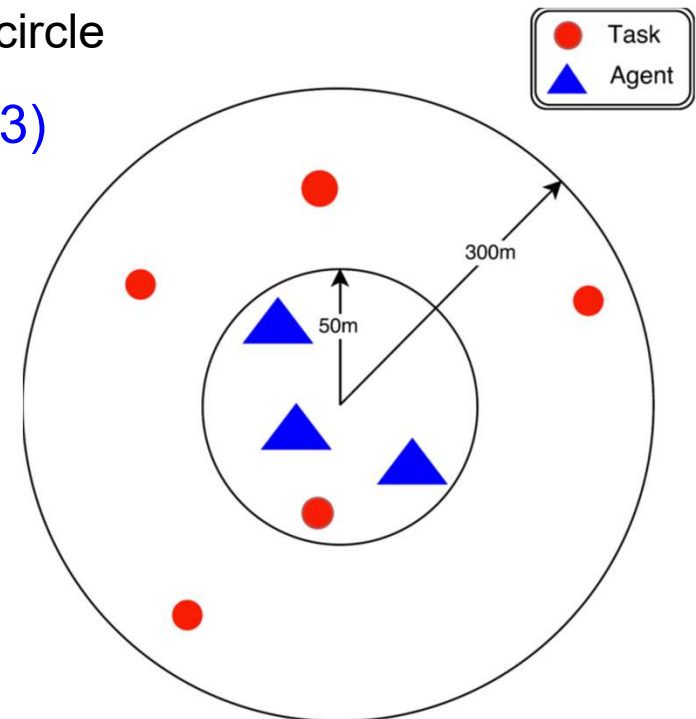
Asynchronous CBBA (ACBBA) vs. CBBA

- ACBBA is conceptually similar to the CBBA, but
 - Each agent builds its bundle and performs consensus asynchronously
 - Each agent only transmits the winning agent, winning bid, and timestamp for a *single task* at a time (less bandwidth required)



ACBBA Simulation Setup

- 100 ACBBA simulation scenarios
 - $N_t = 1, 2, \dots, 10$ tasks and $N_a = 1, 2, \dots, 10$ drones
- Each scenario executed 100 times
 - Drones randomly dropped in 50 m radius circle (ensures connectivity)
 - Tasks randomly dropped in 300 m radius circle
- IEEE 802.11b Wi-Fi broadcast mode (ns-3)
 - No ACKs
 - No retransmissions
 - No exponential backoff
- UDP (ns-3)
 - Connectionless transport protocol



Evaluation Metrics

- **Redundant task assignments**
 - If n_j agents are assigned the same task $j \in T$, then there are $n_j - 1$ redundant assignments of task j
 - Total number of redundant task assignments $n_r := \sum_{j \in T} \max(n_j - 1, 0)$
- **Total number of transmission/reception events**
 - $n_{TX} := \sum_{i \in A} n_{TX,i}$, where $n_{TX,i}$ is the number of times agent i broadcasts its state information
 - $n_{RX} := \sum_{i \in A} n_{RX,i}$, where $n_{RX,i}$ is the number of times agent i receives state information
- **Fraction of received packets**
 - $f_{RX} := n_{RX} / [n_{TX} \cdot (N_a - 1)]$
- **Negotiation time**
 - Elapsed time from the first bundle construction phase to the last conflict resolution phase



Simulation Results (1/3)

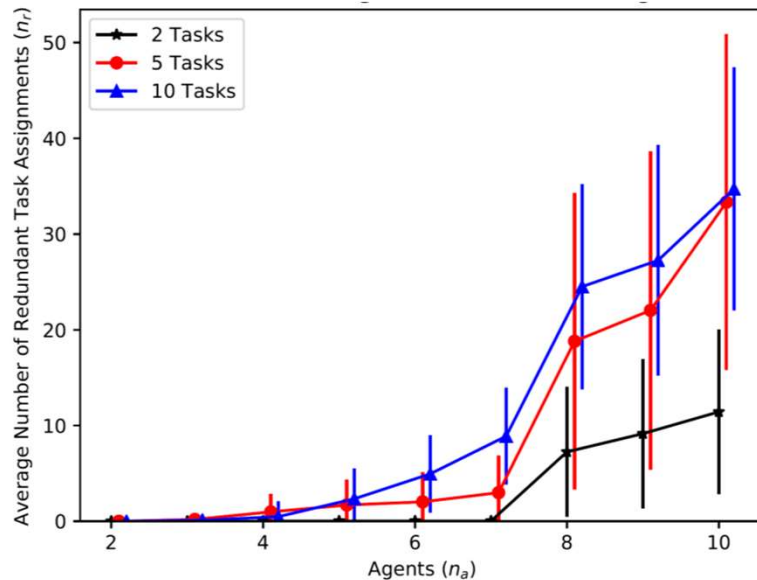


Fig A. Redundant task assignments n_r vs. number of agents N_a .

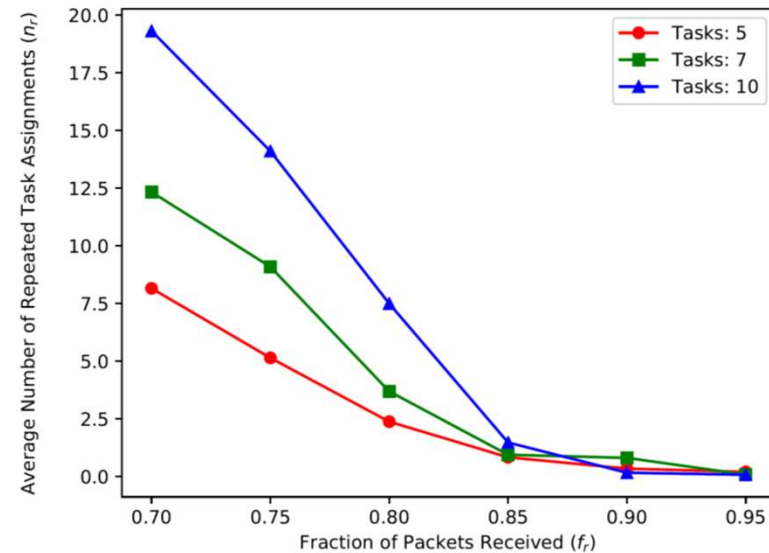


Fig B. Redundant task assignments n_r vs. fraction of packets received f_{RX} .

- Number of redundant task assignments increases with number of agents
- Number of redundant task assignments is negatively correlated with the fraction of packets received
- *What causes this? Channel errors and/or collisions?*



Simulation Results (2/3)

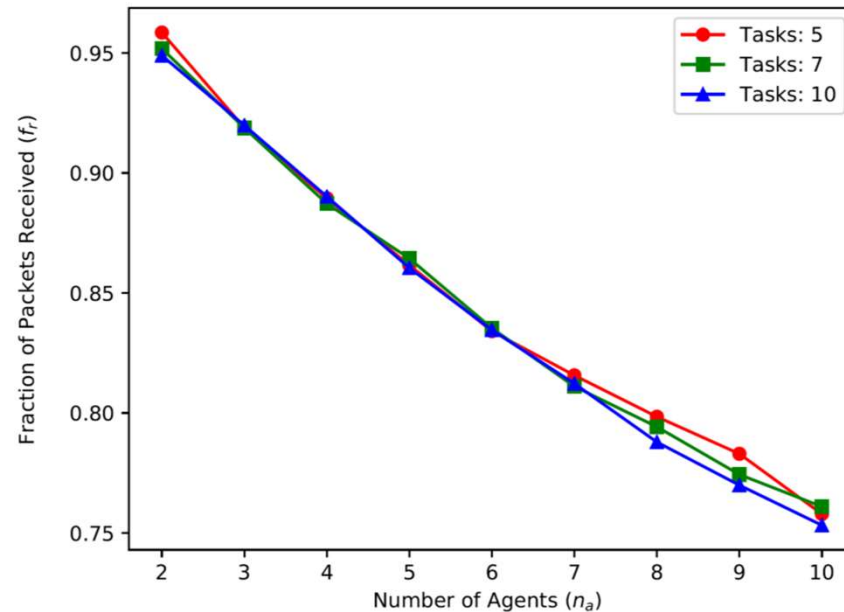


Fig C. Fraction of packets received f_{RX} vs. number of agents N_a

- The fraction of received packets decreases with the number of agents and is approximately invariant in the number of tasks
- Simulation channel errors are independent of number of agents → Performance degradation is primarily due to collisions
- *Why are collisions so problematic?*



Simulation Results (3/3)

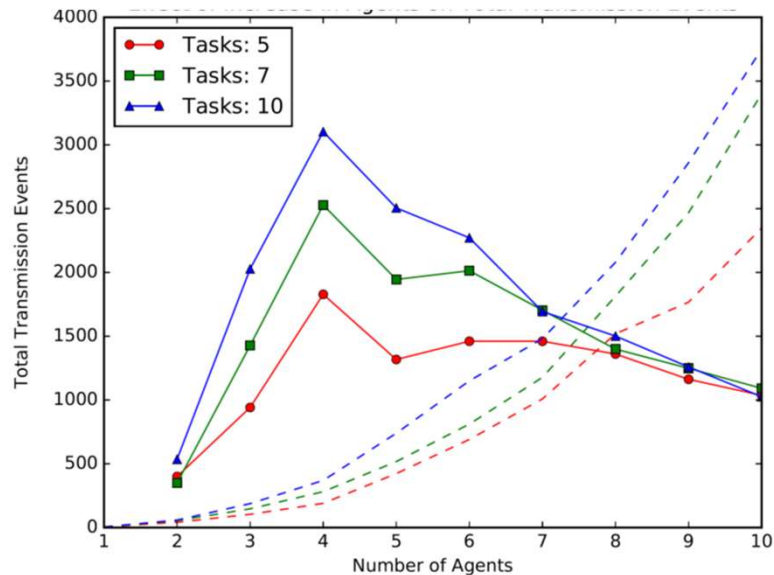


Fig D. Total transmissions n_{TX} vs. number of agents N_a . Solid and dashed lines show lossy and lossless results, respectively.

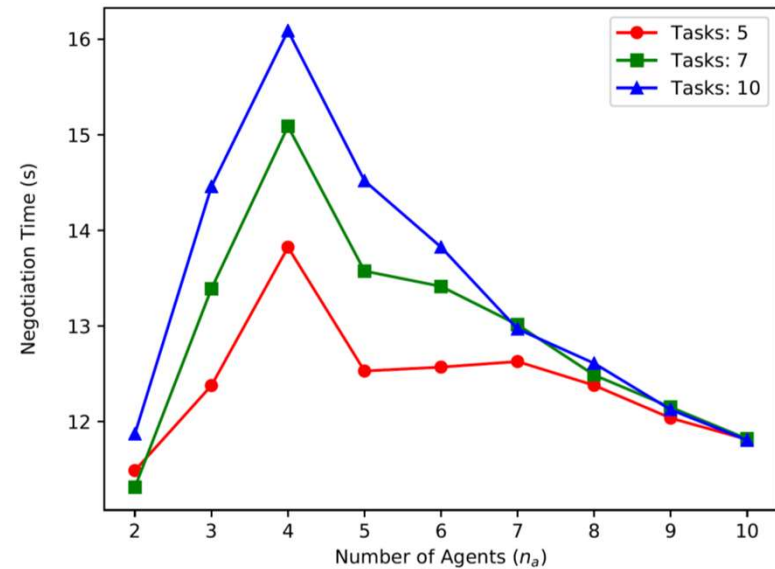


Fig E. Negotiation time (s) vs. number of agents N_a .

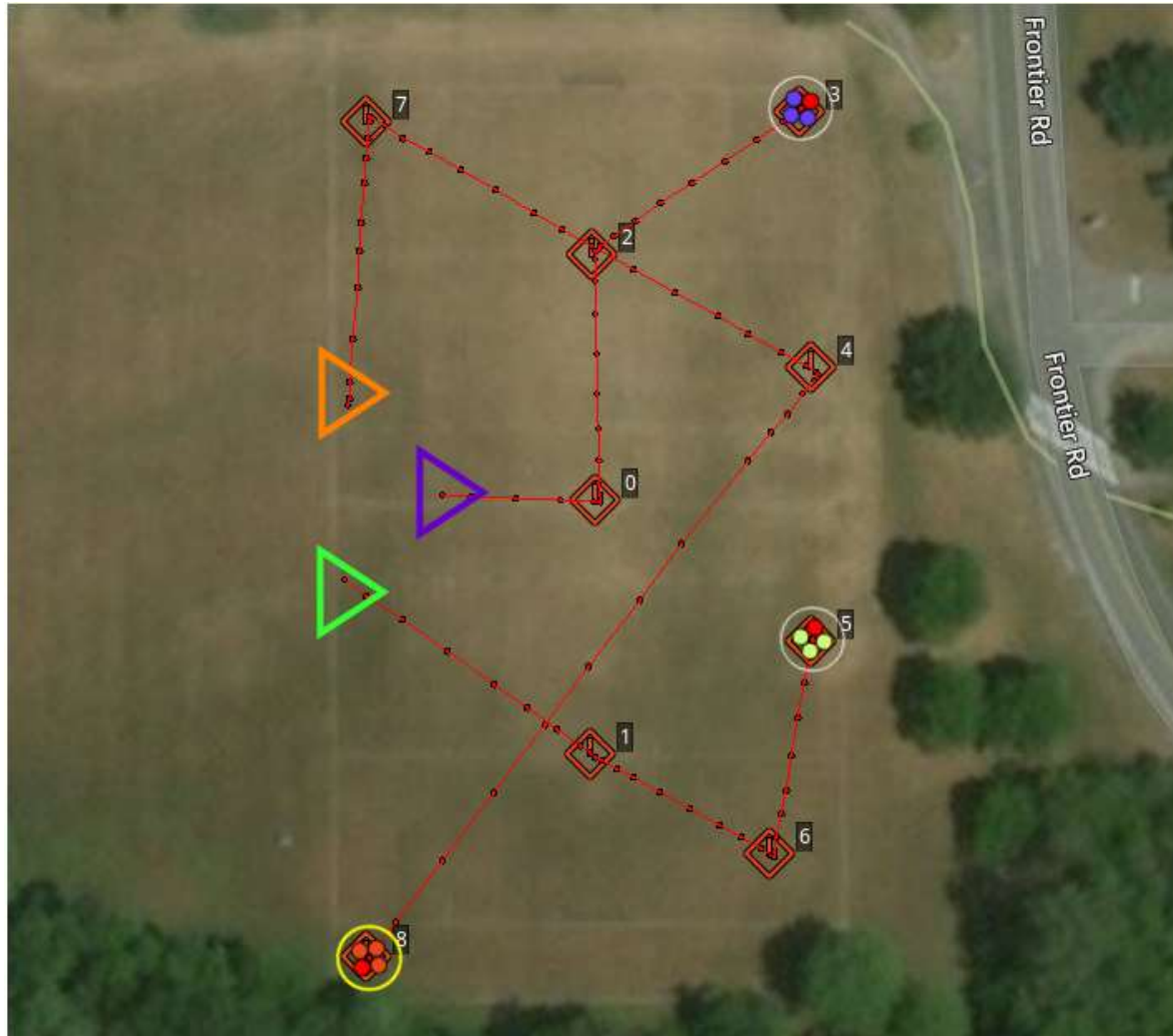
- *Lossless communication:* TX events increase with number of agents
- *Lossy communication:* TX events initially increase with number of agents, but eventually decline due to effect of collisions
- Decline in transmission events is related to shortened negotiation time



- ACBBA yields inefficient task assignments in lossy networks
- Agents mistakenly attribute absence of new messages in network to reaching consensus, when actually due to lost packets
 - Collisions have more significant impact than channel errors
- UDP + IEEE 802.11 broadcast mode provides insufficient QoS
- Ongoing work:
 - Study performance of ACBBA under other network configurations
 - UDP + IEEE 802.11 unicast mode
 - TCP + IEEE 802.11 unicast mode
- Future work:
 - Make ACBBA more robust to network disruptions
 - Study interaction of planning and communications for other applications
 - Swarming, formation control, etc.



Example Execution of ACBBA in the UB-ANC Emulator



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