



# Cryptographic Side-Channel Signaling and Authentication via Fingerprint Embedding

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# Introduction

Fingerprinting & Data Hiding

# Fingerprinting



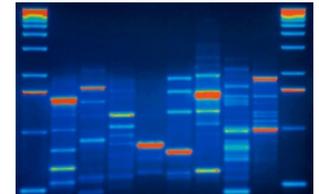
## Intrinsic Fingerprint

- *A characteristic that identifies*
- Uniqueness as a realization of a random process



## Exploit inherent randomness to develop measures of uniqueness

- Biometrics:  
fingerprints, iris scan, DNA, voice, behavioral patterns, ...
- Devices:  
Printers, cameras, scanners, microphones, recorders  
Radios, emitters, amplifiers, waveforms
- Media:  
Paper, canvass



## Desired Fingerprint Properties

- Unique, measurable (convenient & technically feasible)
- Robust to measurement noise
- Develop modeling to assess statistical reliability of ID



# Fingerprint Embedding by Design

## **Purposefully embed fingerprint** for unique ID

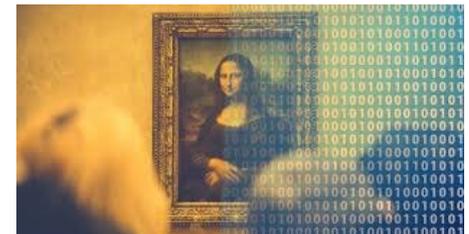
- Defeat cloning (impersonation), tampering

## **Device Manufacturing**

- Many forms for devices
- Intrinsic to randomness inherent in manufacturing
  - Example: transparent material doped with light scattering particles
  - Laser illumination yields unique speckle pattern
- Physically Clonable Function (PUF)
  - Challenge-response paradigm for authentication

## **Steganography (data hiding)**

- Convey hidden messages (Greek: concealed writing)
- Typically binary data: watermark, copyright



# Message Authentication

Classical & PHY-Based

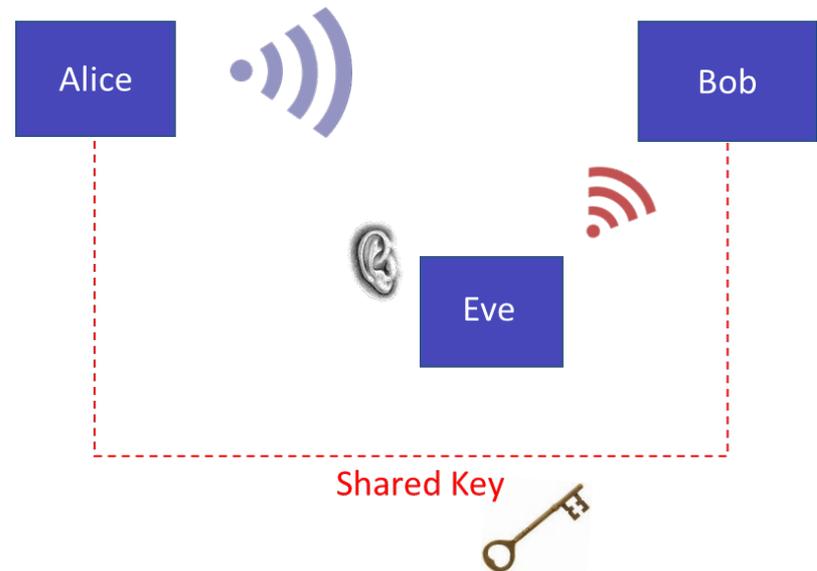
# Wireless Communications Authentication

## Eavesdropper Problem

- *Encryption* for secrecy
- *Authentication* to verify sender ID

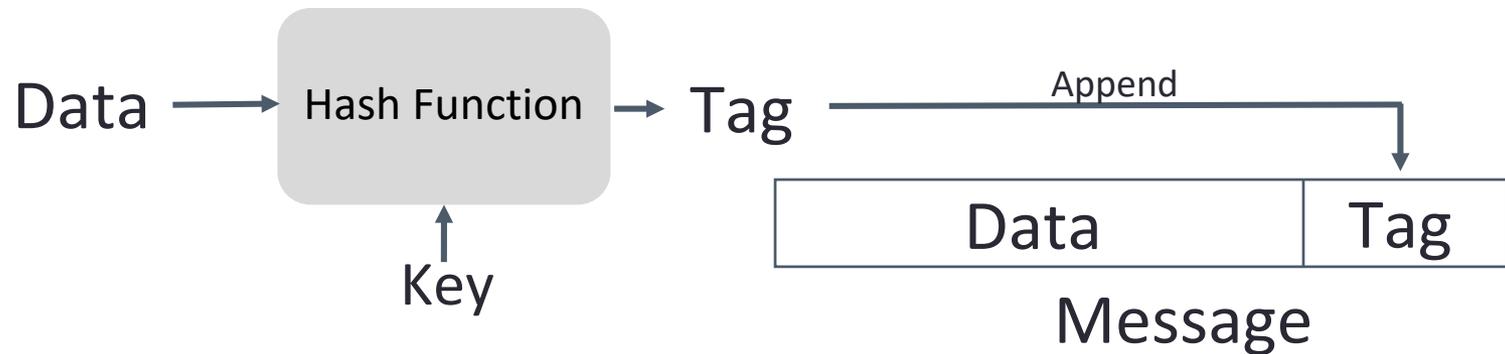
## Why Authenticate Messages?

- Verify identity of sender and safeguards message integrity
- Thwart impersonation and substitution attacks



# Classical Authentication

## Cryptographic HMAC (Hash-based Message Authentication Codes)

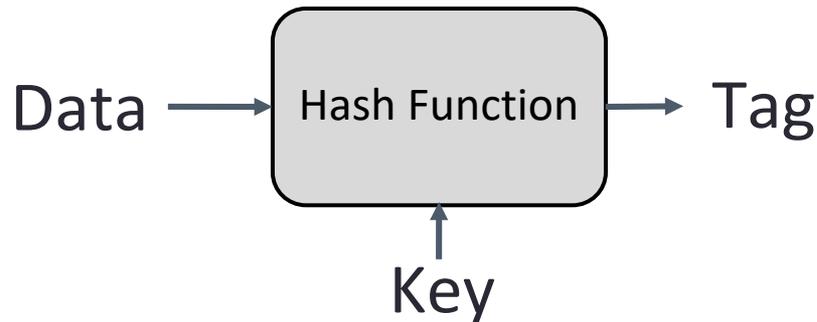


### Issues

- Requires additional bandwidth
- Provides data and tag to Eve
- Only provides computational security

# Crypto-Hash Properties

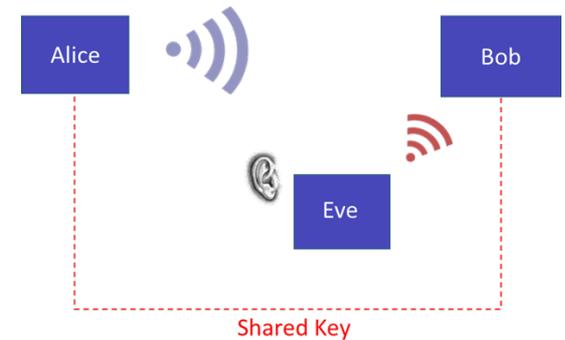
- “One-way function” infeasible to invert: requires brute force search
- Deterministic and efficient
- Resistant to collisions: behaves like a random function
- Model: Changing data or key yields random tag



# Physical Layer Authentication

## Exploit intrinsic physical layer features

- **Device fingerprint ID**
  - ADC, power amplifiers, ...
- **Channel state information (CSI)**
  - Typically: independent time-varying fading provides unique Alice-to-Bob CSI
  - Common source of randomness: Can also provide new secret key
    - Requires reconciliation protocol
- **Issues**
  - Non-tunable
  - Requires favorable channel conditions
  - Uniqueness assumptions



Polak, Dolatshahi, Goeckel, "Identifying wireless users via transmitter imperfections," *IEEE JSAC*, 2011.

Xiao, Greenstein, Mandayam, Trappe, "Using the physical layer for wireless authentication in time-variant channels," *IEEE TWC*, 2008.

Wang, Hao, Hanzo, "Physical-layer authentication for wireless security enhancement: current challenges and future developments," *IEEE Comm Mag*, 2016.

# Fingerprint Embedding Authentication

Tag Embedding

# Our Approach:

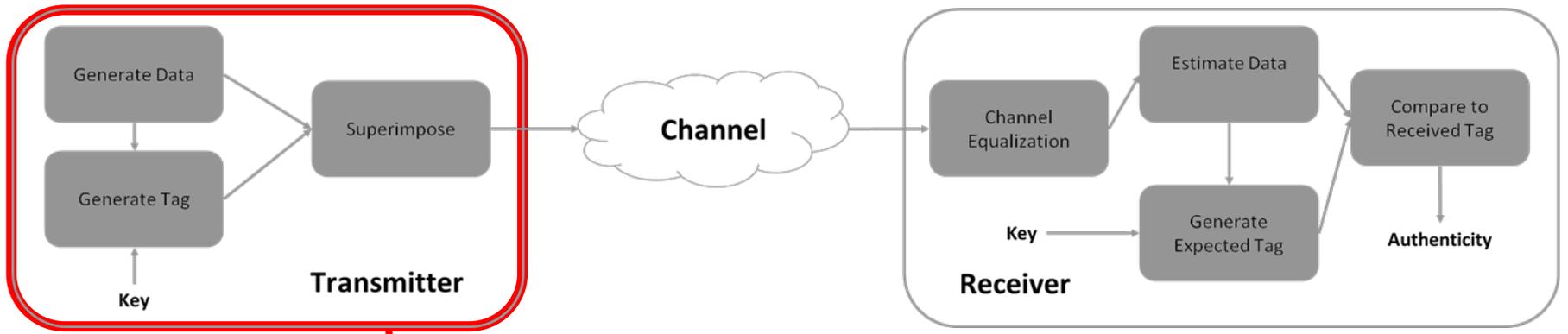
- **Design & Embed cryptographic fingerprint in wireless communications physical layer**
- **Goals:**
  - Secrecy – difficult to detect
  - Security – difficult to estimate and exploit fingerprint
  - Self interference – minimal impact on communications
  - Low complexity – easy to implement
- Enhances information theoretic security (manage key leakage)
- Enhances computational security (raises Eve's complexity)

Does not assume:

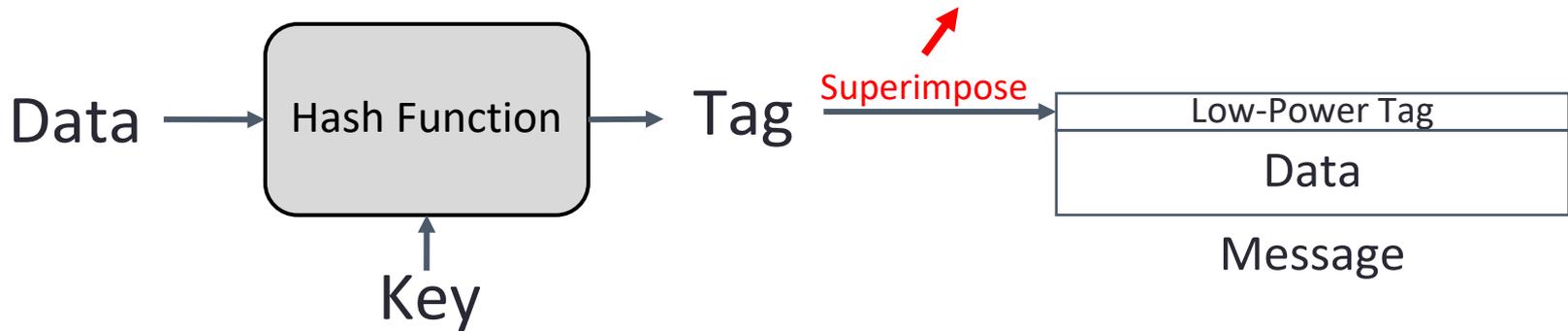
Eve's channel has lower SNR

Alice knows Eve's channel

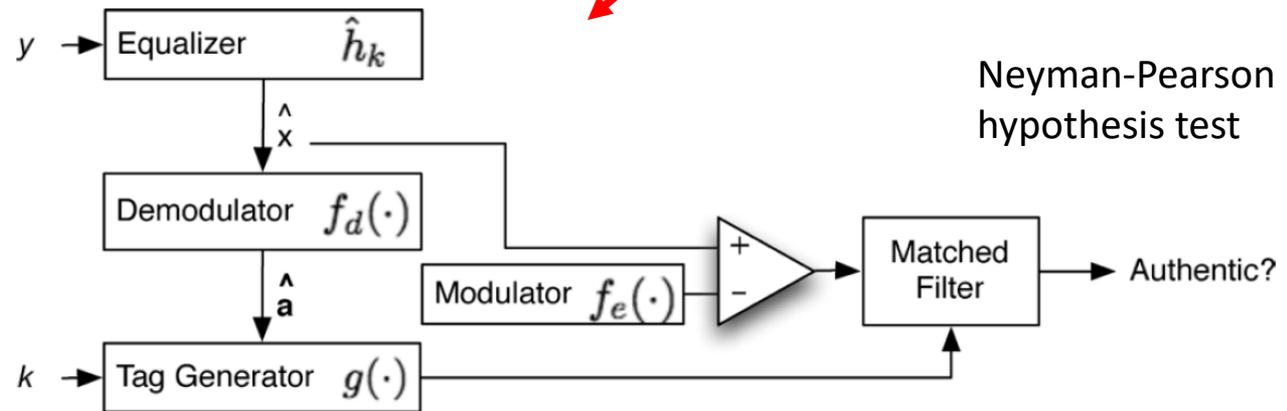
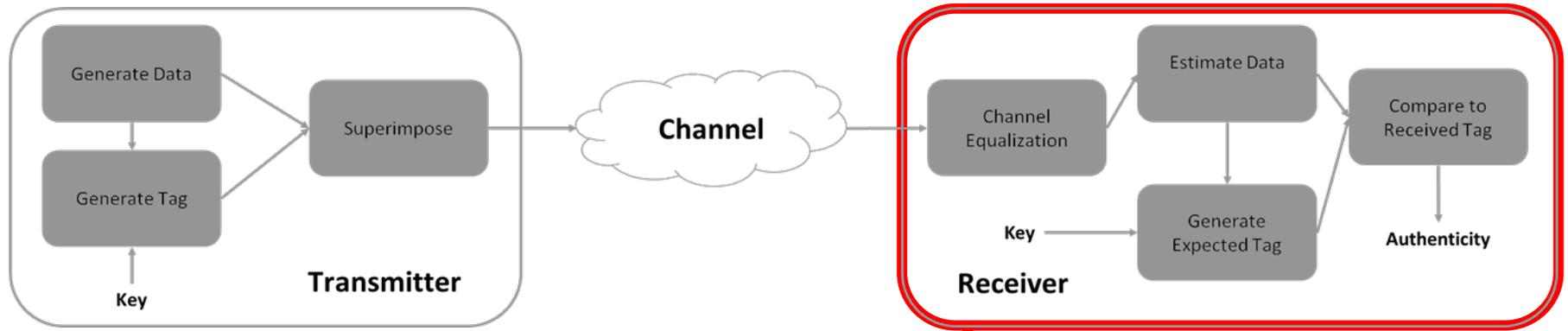
# Tag Embedding



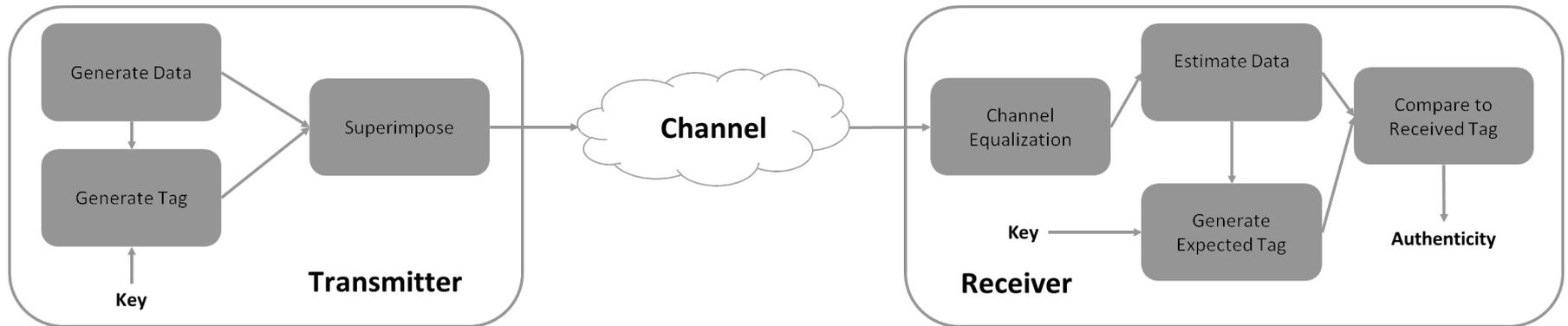
$$X = p_s S + p_t T \quad \text{where } p_s^2 + p_t^2 = 1 \quad p_t \ll p_s$$



# Authentication Hypothesis Test



# Authentication via Fingerprint Embedding

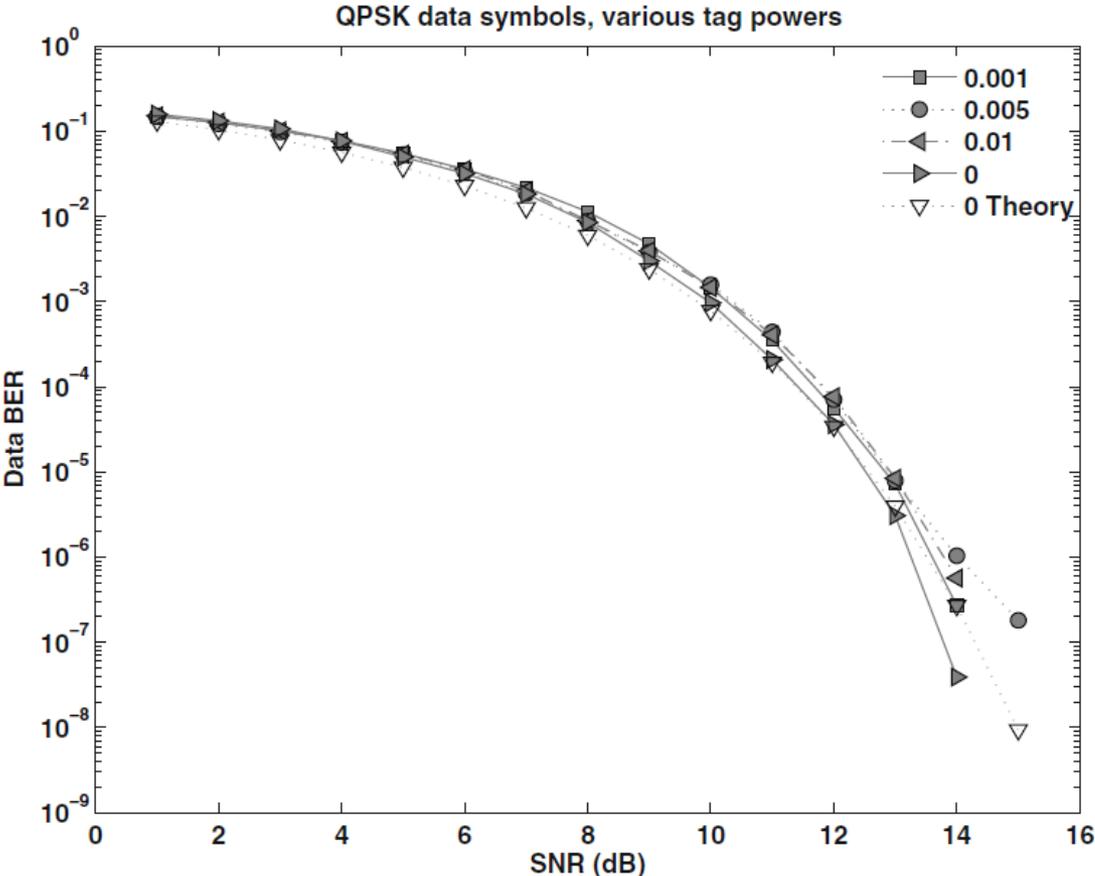


- No additional bandwidth
- Symbol synchronous, low complexity
- Many variations possible, e.g.,
  - Coupling with other security methods
  - Nonlinear embedding

# SDR SISO Experiment



- Minimal impact of ~1% tag power on receiver BER



# SDR SISO Experiment

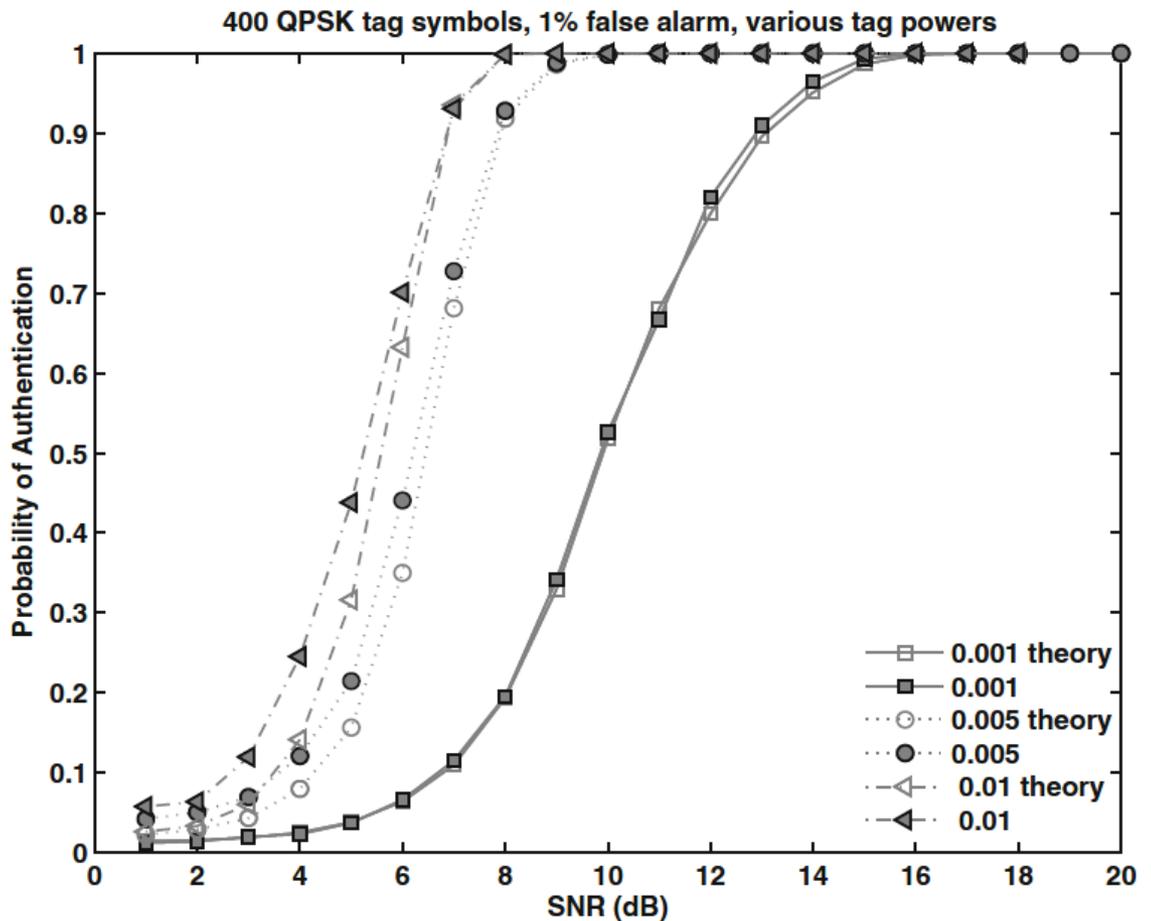


- Tag power tradeoffs

- Enhances authentication performance

*versus*

- Higher SNR for Eve's tag estimate
- Small decrease in Bob's SNR



# MIMO Authentication

- Known channel state info (CSI)

Pre-coding  $\mathbf{X} = \gamma_S \mathbf{F}_S \mathbf{P}_S^{\frac{1}{2}} \mathbf{S} + \gamma_T \mathbf{F}_T \mathbf{P}_T^{\frac{1}{2}} \mathbf{T}$

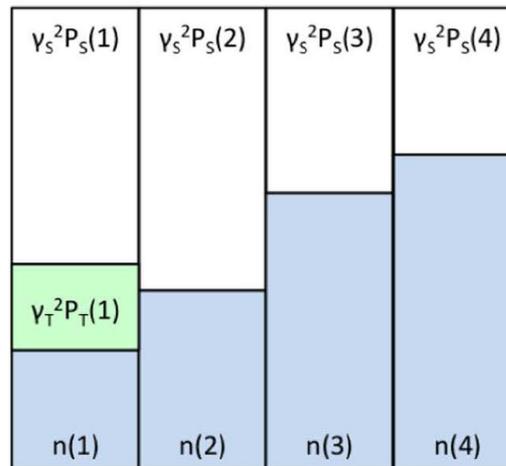
Received  $\mathbf{Y} = \sqrt{g} \mathbf{H} \mathbf{X} + \mathbf{W}$

Residual

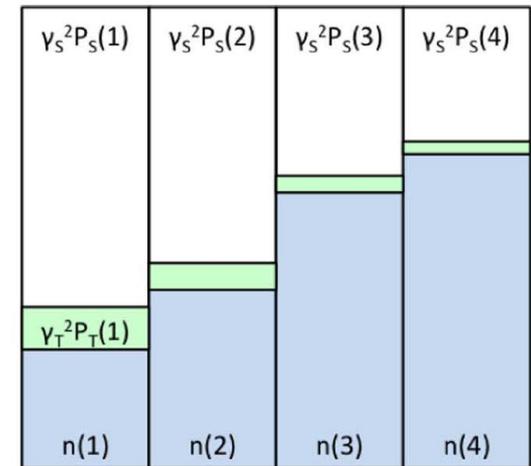
$$\hat{\mathbf{Q}} = \sqrt{g} \gamma_T \hat{\mathbf{H}} \mathbf{F}_T \mathbf{P}_T^{\frac{1}{2}} \tilde{\mathbf{T}}$$

Test Statistic

$$\tau = \Re[\text{Tr}(\hat{\mathbf{Q}}^\dagger \mathbf{Q})]$$



Strongest mode only

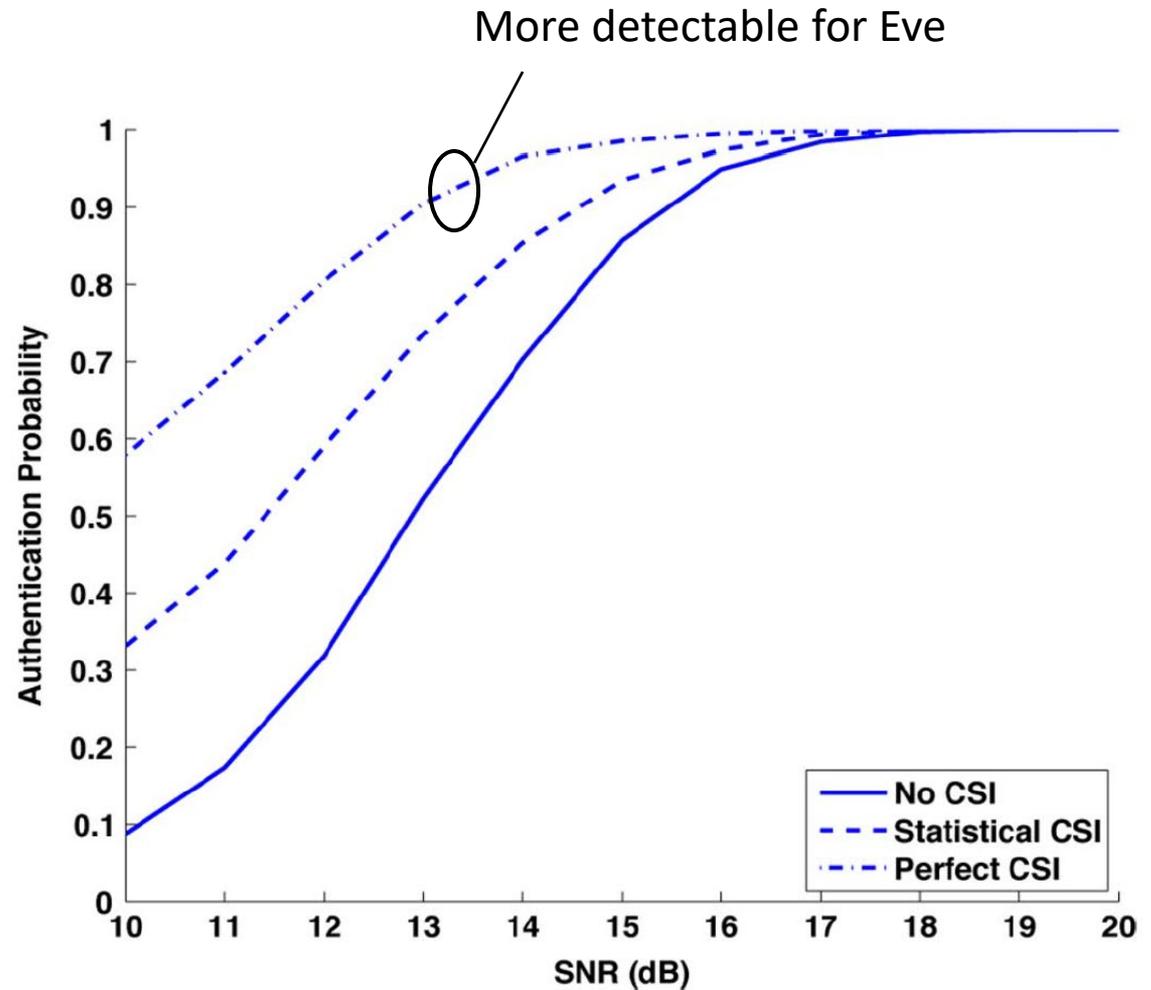


All modes proportionally

# MIMO Authentication

4x4 MIMO Simulation:

- 4 x 256 symbols
- Rayleigh fading
- Multi-mode tagging



# Security

Key Information Leakage

# Key Information Leakage

## Conditional Entropy:

- *Equivocation* (calling two different things by the same name)
- Assume Eve knows architecture, parameters, and hash function
  - Zero equivocation in noise free case & if hash is uniquely invertible

$$H(k|Y, \theta) = \sum_{s \in \mathcal{S}, t \in \mathcal{T}} p(s, t) H(k|s, t)$$

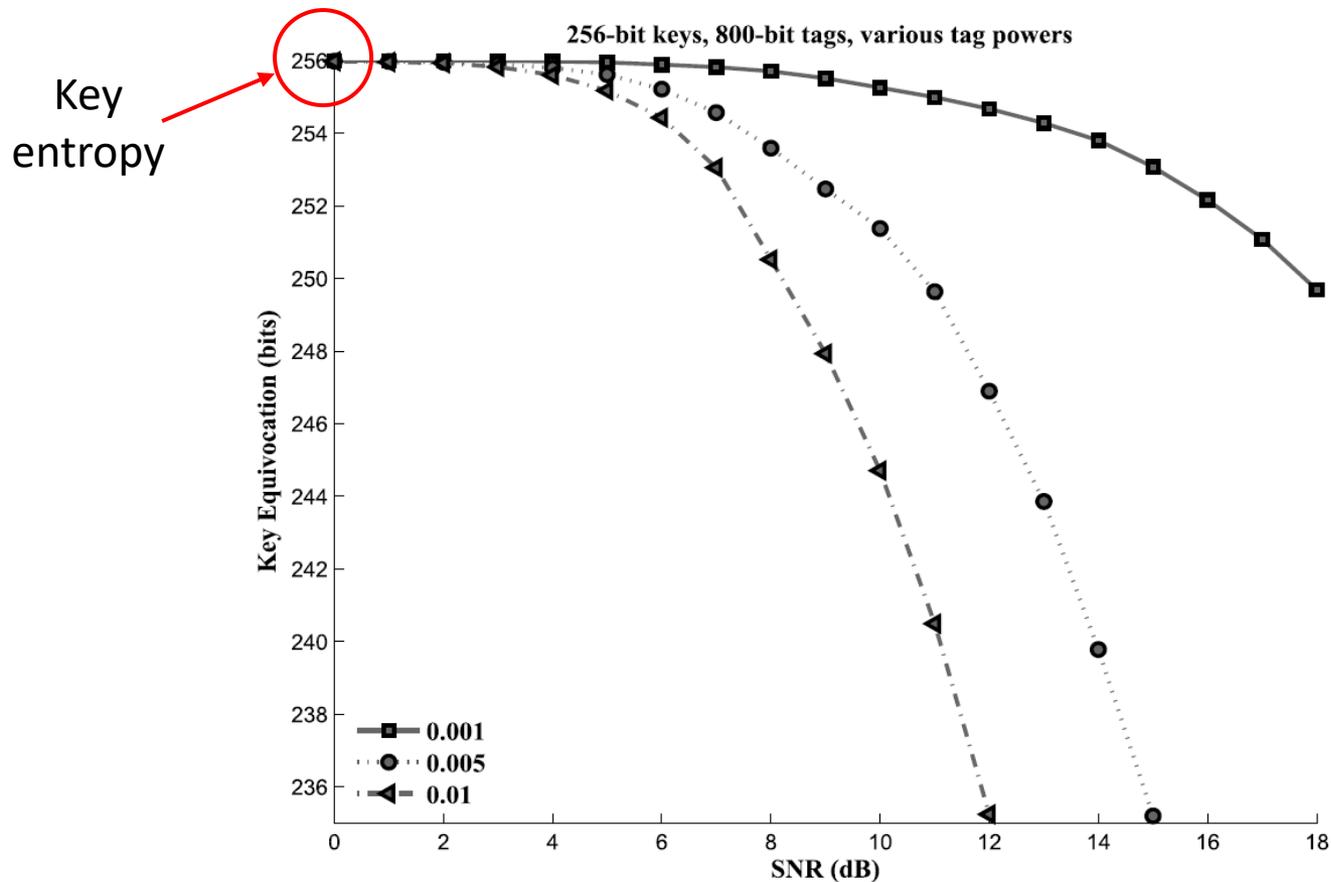
$$H(k|Y) \cong \frac{|\mathcal{K}|}{|\mathcal{T}|} \sum_{i=0}^{\log |\mathcal{T}|} \binom{\log |\mathcal{T}|}{i} H\left(\frac{|\mathcal{T}|}{|\mathcal{K}|} p_e^i (1-p_e)^{\log |\mathcal{T}| - i}\right)$$

Randomness through Eve's bit error probability

# Key Information Leakage



- SISO Conditional Entropy (single Tx)  
Provides insight into key update strategy

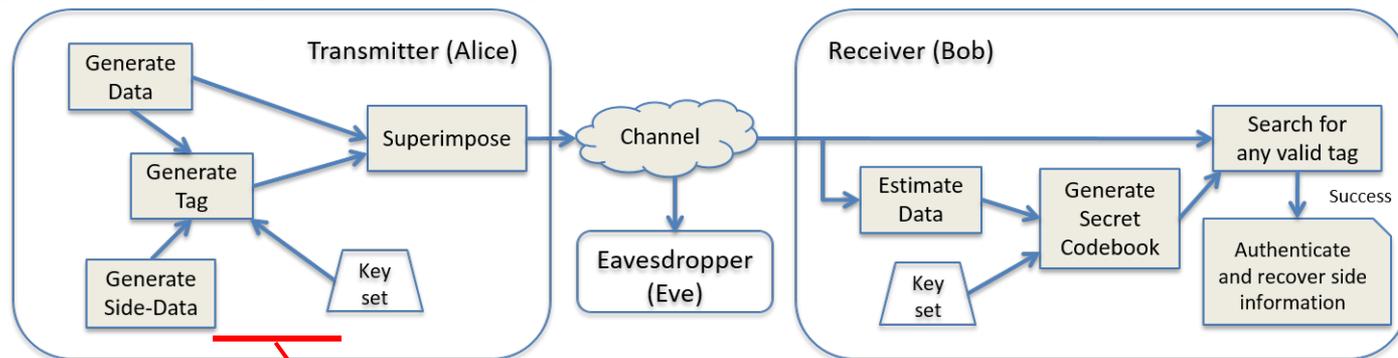


# Communications in the Side-Channel

Creating a Secret Codebook of Tags

# Authentication + Side-Channel Comms

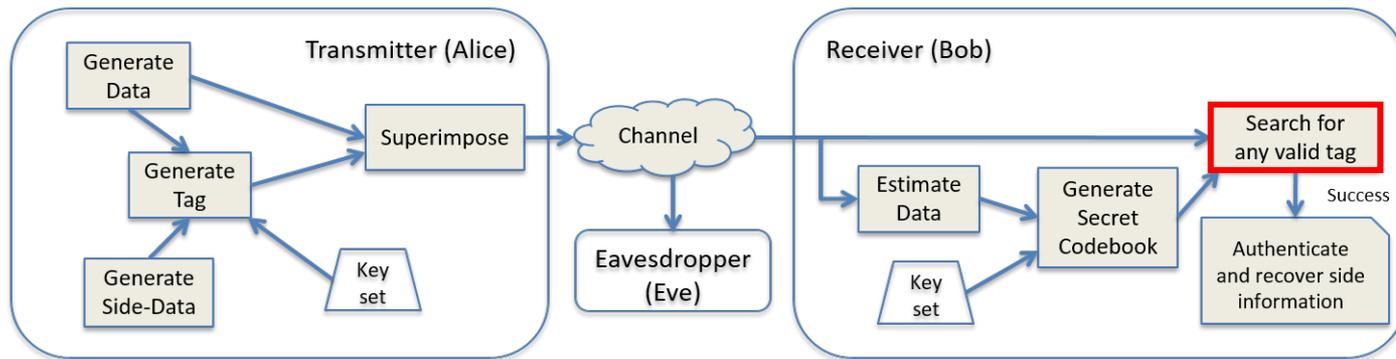
Block Diagram of Multi-Key Authentication System



Communicate via Key Choice

# Authentication + Side-Channel Comms

Block Diagram of Multi-Key Authentication System



$$\text{Test over codebook entries} \left\{ \begin{array}{l} H_0 \quad \text{No valid tag was sent} \\ H_1 \quad t_1^{\text{valid}} \text{ was sent, } m = 1 \\ \vdots \\ H_{N_k} \quad t_{N_k}^{\text{valid}} \text{ was sent, } m = N_k \end{array} \right.$$

Authenticates & recovers side-channel symbol

# Secret **Random** Codebook: 2 Designs

0. Key is partitioned into  $N_k$  sub-keys

## 1. Simple Codebook Construction

- One sub-key per symbol
- **$\log_2 N_k$  bits communicated**

## 2. Linear Codebook Construction

- $N_k$  possible tags are rows in generator matrix  $G$
- Transmit  $\mathbf{m}$  by linear combination of possible tags
- **$N_k$  bits communicated**

$$\overline{\mathbf{G}} = \begin{bmatrix} \mathbf{t}_1^{\text{valid}} \\ \mathbf{t}_2^{\text{valid}} \\ \vdots \\ \mathbf{t}_{N_k}^{\text{valid}} \end{bmatrix}$$

$$\begin{aligned} \mathbf{t}^{\text{xmit}} &= \mathbf{m}\overline{\mathbf{G}} \\ &= \sum_{j=1}^{N_k} \mathbf{m}_j \mathbf{t}_j^{\text{valid}} \end{aligned}$$

# Authentication Performance

$$\Pr \text{Decide } H_1 | H_1 = \int_{\tau_{1,0}}^{\infty} \Phi^{N_k-1} \left( \frac{z}{\sqrt{\frac{L}{2} + \sigma_{\tilde{w}}^2}} \right) \phi \left( \frac{z-L}{\sigma_{\tilde{w}}} \right) F_{\tau_1}(z) dz,$$

WLOG assumes H1 true

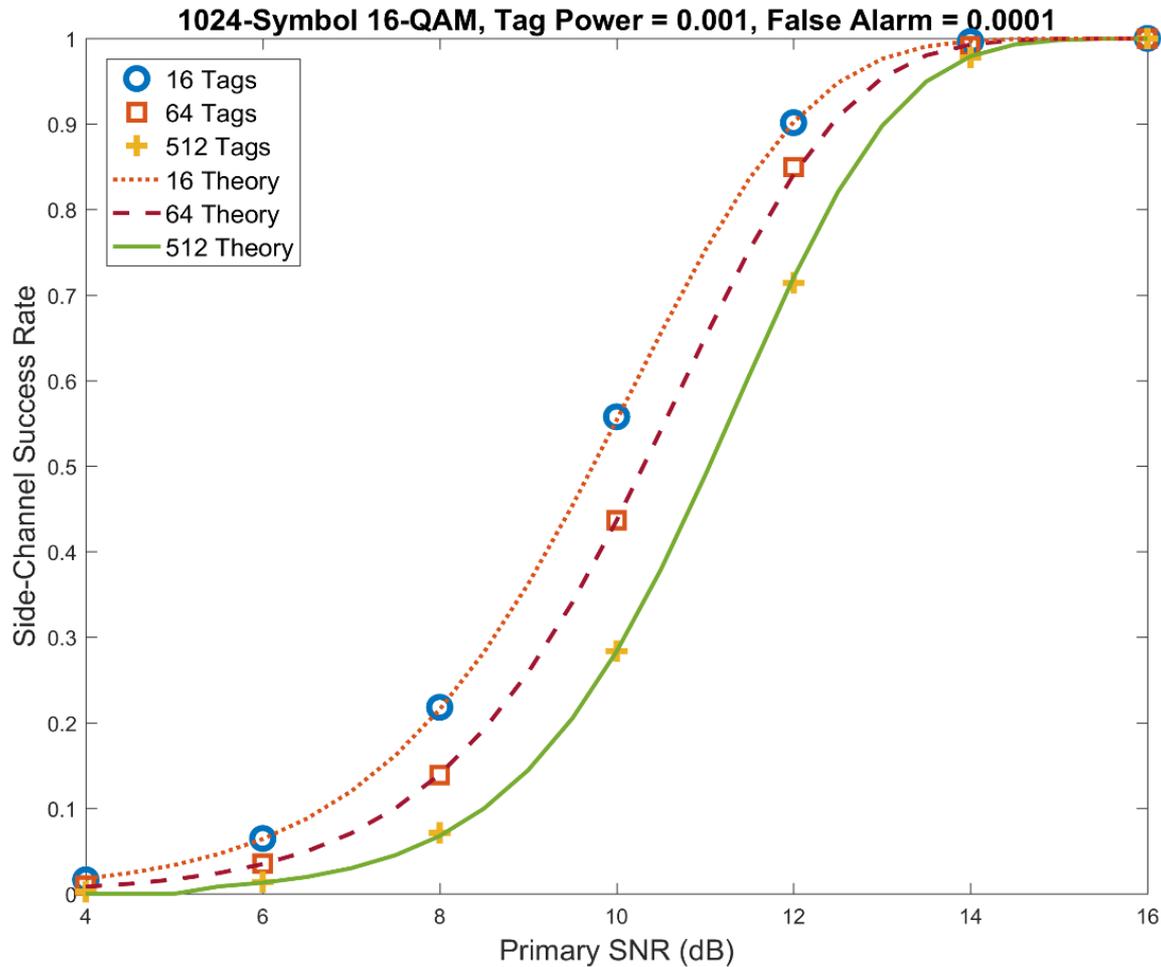
$\tau_{1,0} \triangleq \tau_1 | H_0$  and  $F_{\tau_1}(z) = \Pr \tau_1 < z$  is the CDF of  $\tau_1$

Threshold under  
H0 is constant

$$\tau_i | H_j (\neq i) = \min_{\tau} \text{ s.t. } \Pr Z_i(R | H_j) > \tau < \alpha$$

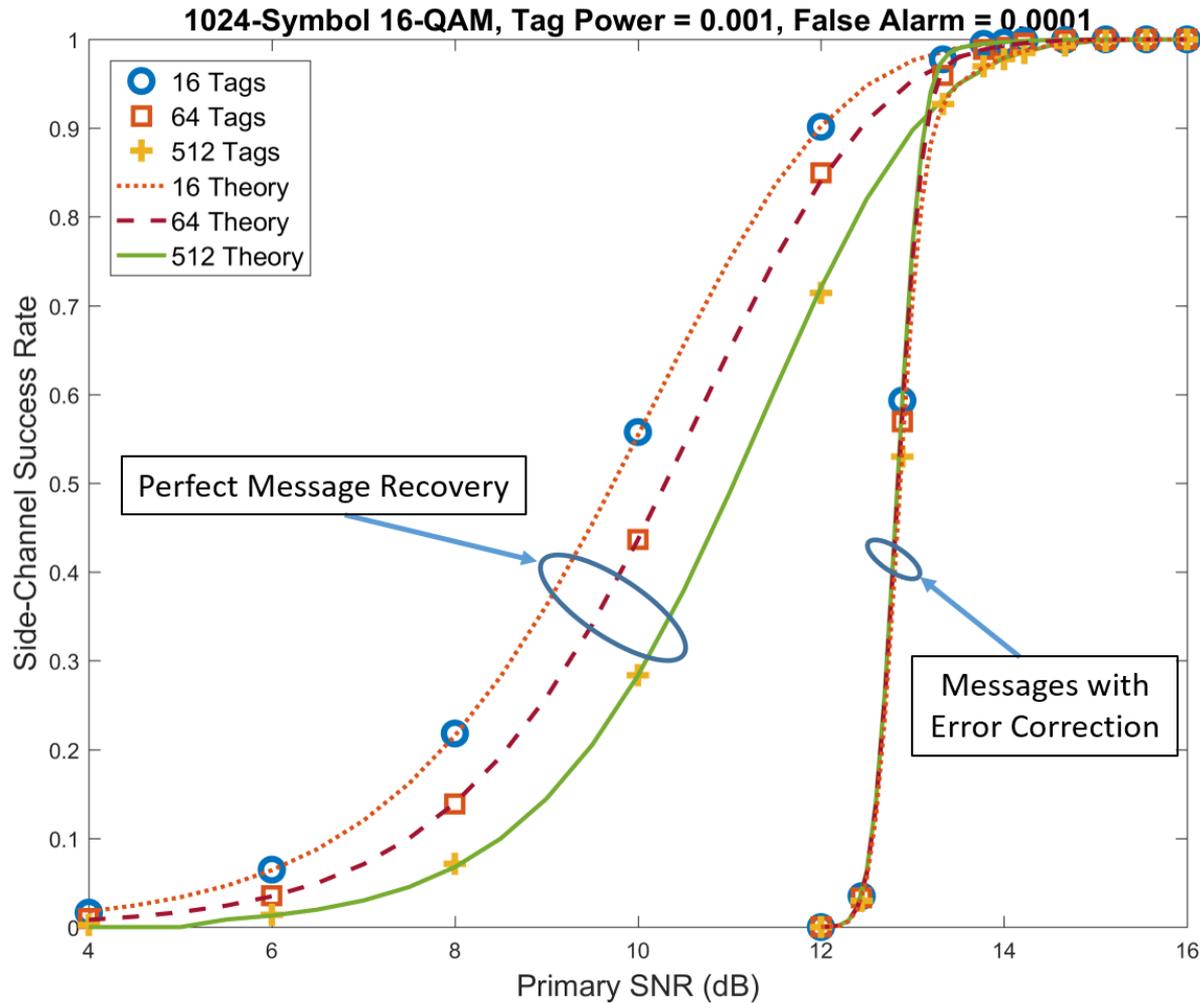
Thresholds are recalculated by Bob for each transmission  
(New Random Codebook)

# Side Channel Performance: No Data EC Coding



Assumption: Bob correctly reconstructs secret codebook  
(Primary message obtained without error)

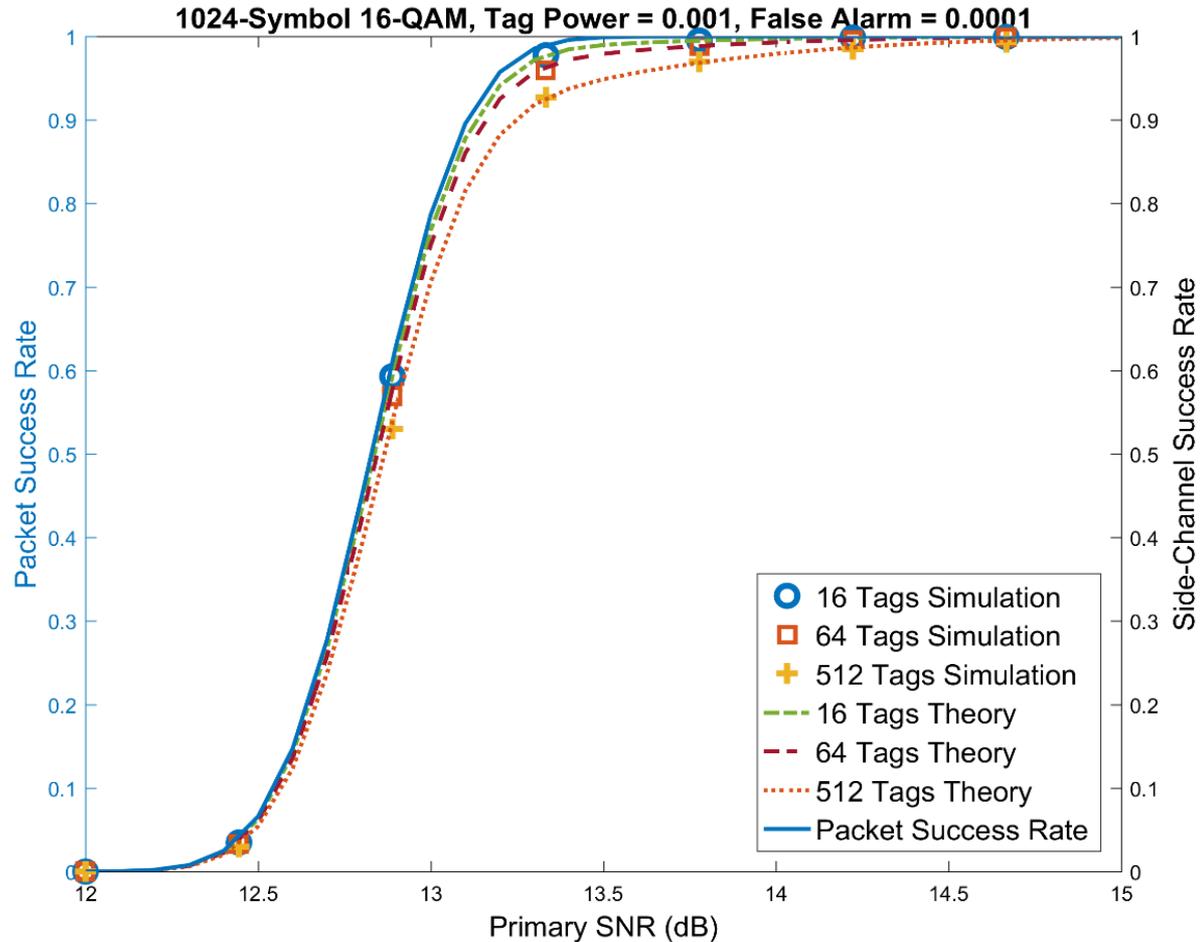
# Performance w/ Data Error Correction Coding



BCH: 10 bit EC  
512 block

Bit error causes random codebook mismatch

# Performance w/ Data Error Correction Coding



Performance dominated by packet success rate

# Security

Multi-Key Codebook Scheme

# Key Information Leakage

## Conditional Entropy:

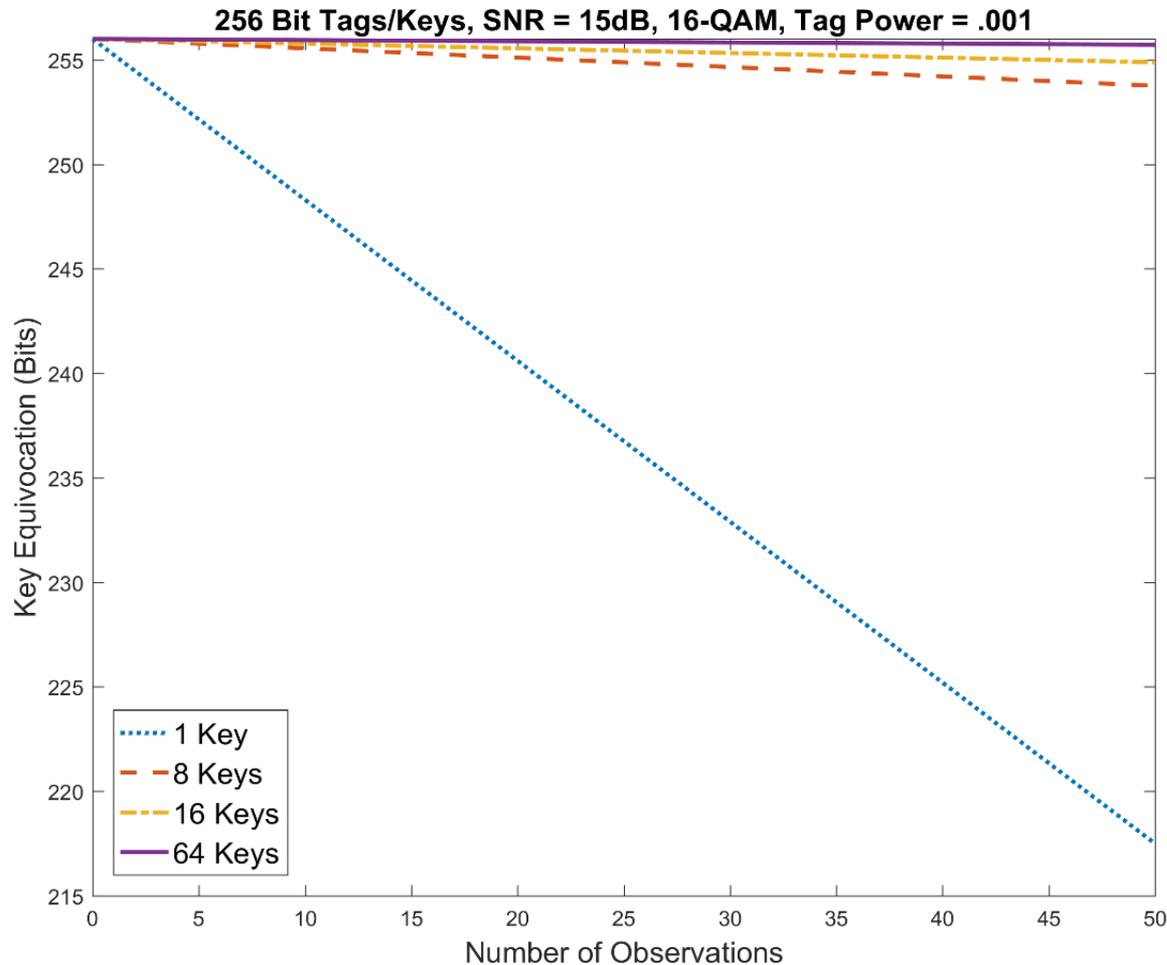
$$H(k|Y, \theta) = \sum_{s \in \mathcal{S}, t \in \mathcal{T}} p(s, t) H(k|s, t)$$

$$H(k|Y^n; \theta) \cong \frac{|\mathcal{K}|}{|\mathcal{T}|^{\frac{N}{N_k}}} \sum_{i=0}^{\frac{N}{N_k} \log_2 |\mathcal{T}|} \binom{\frac{N}{N_k} \log_2 |\mathcal{T}|}{i} H \left( \frac{|\mathcal{T}|^{\frac{N}{N_k}}}{|\mathcal{K}|} p_e^i (1 - p_e)^{\frac{N}{N_k} \log_2 |\mathcal{T}| - i} \right)$$

## Computational Security:

Multi-key attribution problem increases Eve's search space  
Much worse for linear codebook

# Key Leakage

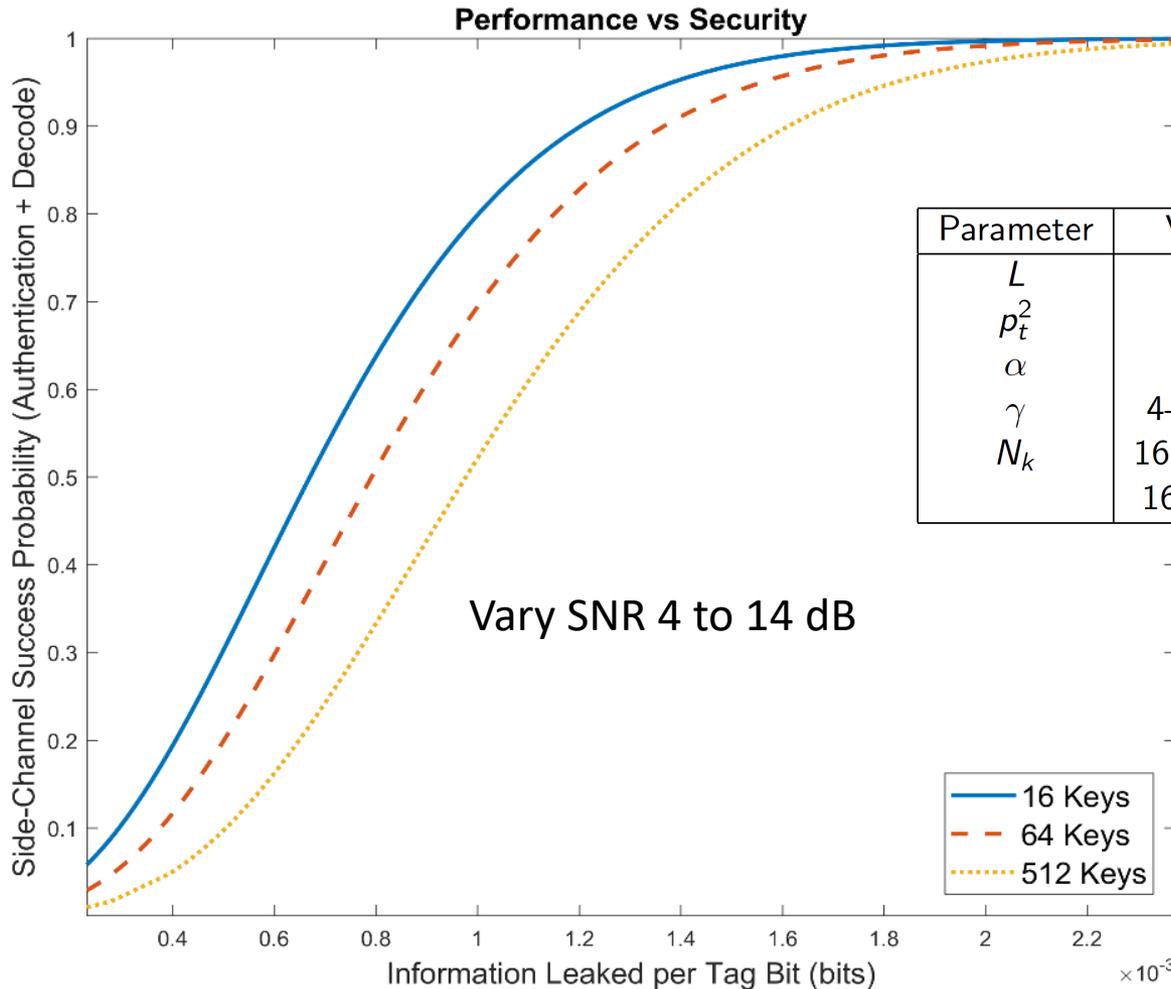


Simple Codebook

Assume Eve knows key assignment

Eve needs more observations to obtain information about a sub-key

# Security - Performance Trade-off: Side-Channel Success vs Key Leakage



# Conclusion

- Design framework yields good tradeoffs in secrecy, security, self-interference, and complexity

## Going Further :

- Couple approach with PHY layer encryption & jamming, active & passive techniques
  - MIMO, directional modulation, beamforming
- Networking & broadcast authentication
- Key evolution using the side-channel

# References

- J. B. Perazzone, P. L. Yu, B. M. Sadler, R. S. Blum, “Cryptographic side-channel signaling and authentication via fingerprint embedding,” *IEEE Transactions on Information Forensics and Security*, vol. 13, no. 9, pp. 2216--2225, 2018.
- P. L. Yu, B. M. Sadler, G. Verma, J. S. Baras, “Fingerprinting by design: embedding and authentication,” in *Digital Fingerprinting* (Springer, 2016), C. Wang, R. M. Gerdes, Y. Guan, S. K. Kasera, editors.
- P. L. Yu, G. Verma, B. M. Sadler, “Wireless physical layer authentication via fingerprint embedding,” *IEEE Communications Magazine, special issue on Wireless Physical Layer Security*, pp. 48--53, June 2015.

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