

# HONEY CHATTING: A NOVEL INSTANT MESSAGING SYSTEM ROBUST TO EAVESDROPPING OVER COMMUNICATION

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## INTRODUCTION

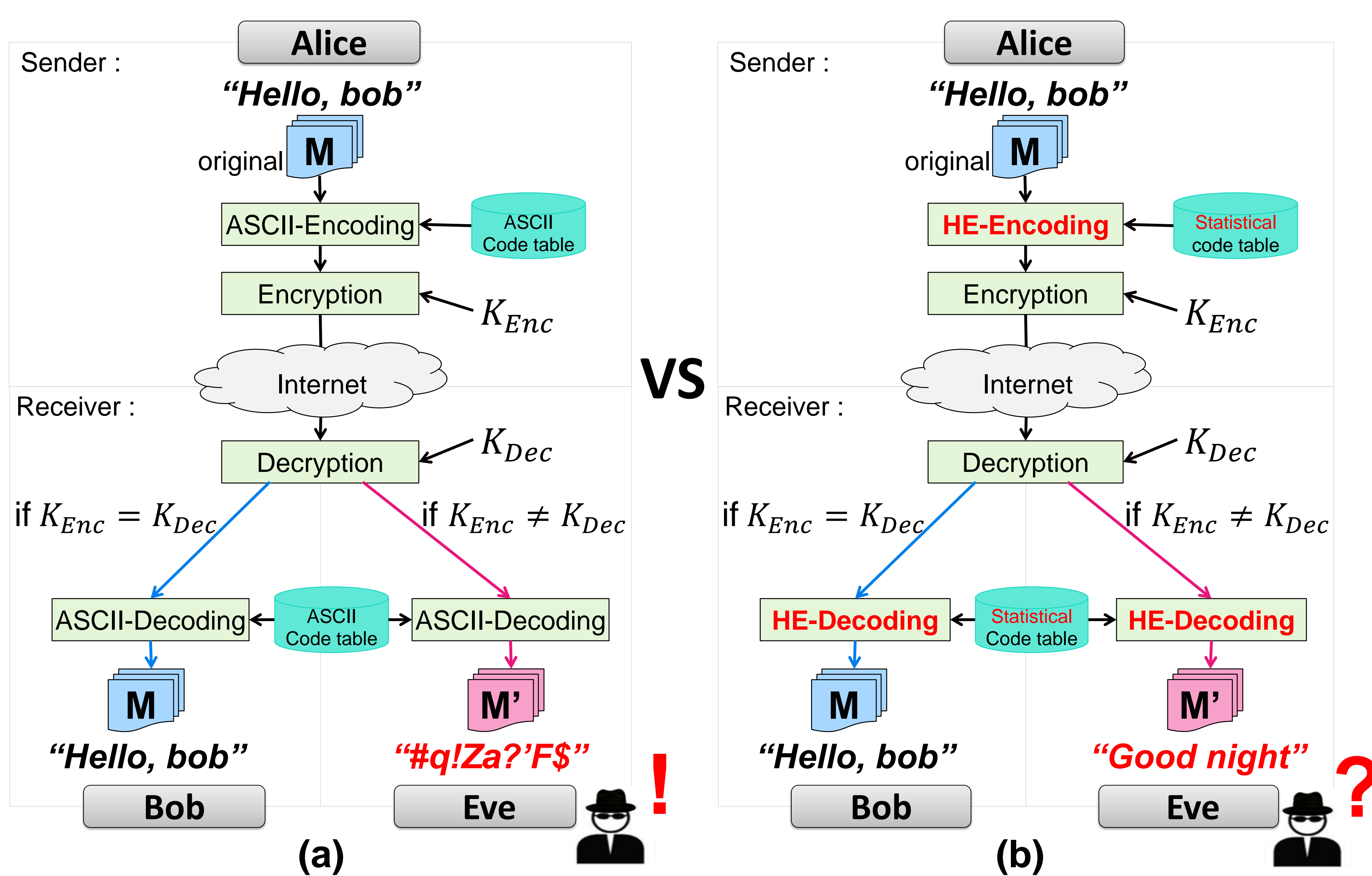
### Secure Chatting

- One of methods to strengthen the security of Instant Messaging system is the message encryption.
- However, the key used for encryption has potential vulnerability to be cracked by a brute-force attacker if the key size is not long enough.
- So, we introduce a new concept of secure chatting by applying Honey Encryption which makes decrypted texts with wrong keys hard to be distinguished from the decrypted text with a real key.

### Our Work

- We develop a Instant Messaging system(Honey Chatting) robust to eavesdropping by using the basic idea of Honey Encryption.
- In our system, we generate plausible-looking but fake plaintexts by using statistical encoding/decoding scheme to confuse the brute-force attacker.
- Through simple experiments, we show the difference between a real message and fake messages by calculating the entropy of texts in the decrypted message.

## STRUCTURE SUMMARY



### Overall Procedure

- The sender's message  $M$  is encoded using the code table and encrypted with  $K_{Enc}$ .
- It passes through the communication channel such as Internet.
- The receiver decrypts it with  $K_{Dec}$  and decodes it using the same code table. If  $K_{Enc} = K_{Dec}$ , the receiver can obtain a true message in both cases. Else If  $K_{Enc} \neq K_{Dec}$ ,  $M'$  is become false message in (a) or plausible-looking fake message in (b).

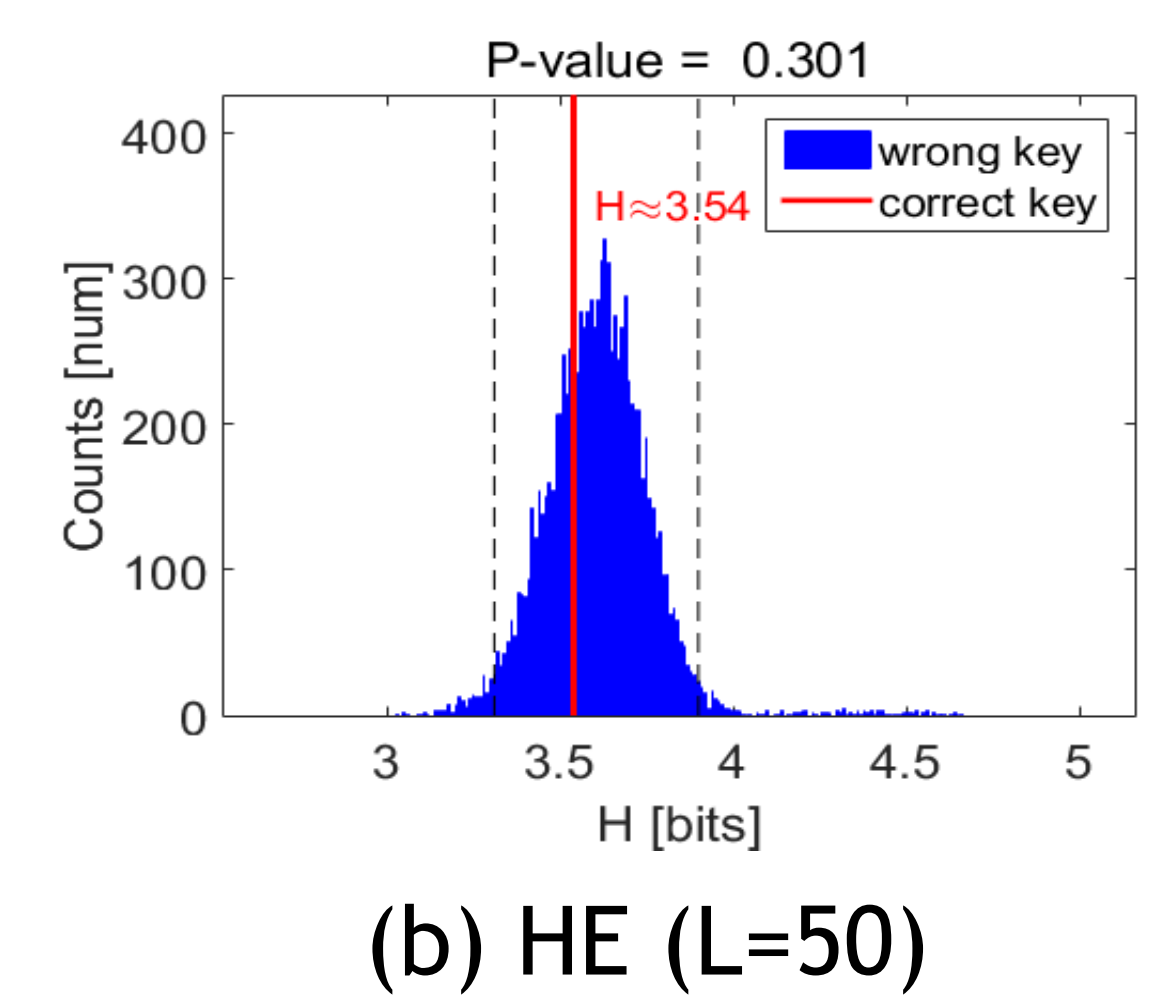
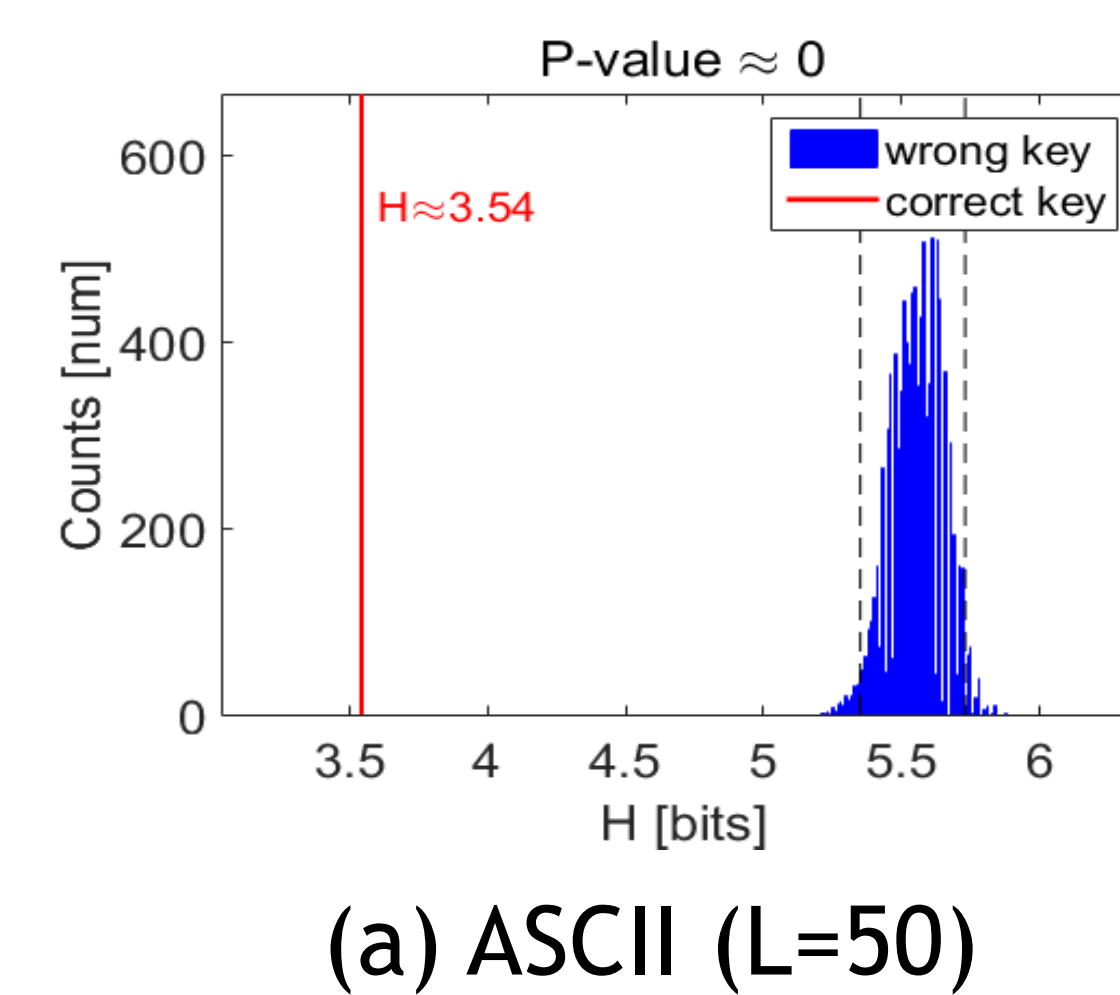
→ Therefore, Eve(brute-force attacker) would confuse to find real message.

❖ Here, the **Statistical code table** is made from the **statistical coding scheme** using text corpus in advance, and the sender and the receiver **share** it.

## EXPERIMENT

### The Difference of Entropy

- We conduct a significance test(hypothesis test) to show difference between decrypted text with wrong and real key when applying HE scheme.
- In (a), P-value is significantly small which means there are clear distinction between  $M$  and  $M_s$ .
- In (b), moderately large P-value shows that observed data  $M$  is agreed with  $M_s$ . It means that  $M$  is similar with  $M_s$ , so the brute-force attacker could not notice his success.



### Statistical Coding Scheme

- Chat messages can be represented by N-gram language model, so we get the probability of consecutive characters in a sentence.
- We construct the **cumulative massive function(CMF)** based on the N-gram language model. CMF is used as statistical code table for HE-Encoding and HE-Decoding.
- The CMF for  $i$ -th character of message :

$$p_{cmf}^{(i)}(c_k) = \frac{\sum_{k=0}^S p(x_i = c_k | \mathbf{x}_{i-1:i-n})}{\sum_{j=0}^S p(x_i = c_j | \mathbf{x}_{i-1:i-n})}$$

- $S$  is the number of possible character set // a-z, space, comma, period
- $n$  is the order of markov process. //  $n=5$  in our application
- $p(x_i | \mathbf{x}_{i-1:i-n})$  is the  $i$ -th character influenced by previous  $n-1$  characters

## HONEY CHATTING SIMULATION

- Situation:** While Alice and Bob enjoy chatting (share a real password), a malicious Eve is trying to eavesdrop their chat messages (try to enter wrong passwords).

|                    |         |
|--------------------|---------|
| UserID:            | alice   |
| Passwd:            | realkey |
| Choose File:       | movies  |
| Connect Disconnect |         |

[alice] hello, bob, how's it going  
[alice] it is simulation of our chatting program.  
[alice] alice and bob, who shared same secret key,  
[alice] can see real plain text message.  
[bob] however, malicious user eve will get fake message  
[bob] if eavesdrops their communication message.

|                    |         |
|--------------------|---------|
| UserID:            | bob     |
| Passwd:            | realkey |
| Choose File:       | movies  |
| Connect Disconnect |         |

[alice] hello, bob, how's it going  
[alice] it is simulation of our chatting program.  
[alice] alice and bob, who shared same secret key,  
[alice] can see real plain text message.  
[bob] however, malicious user eve will get fake message  
[bob] if eavesdrops their communication message.

|                    |          |
|--------------------|----------|
| UserID:            | eve      |
| Passwd:            | wrongkey |
| Choose File:       | movies   |
| Connect Disconnect |          |

[alice] the door opens to the table  
[alice] i don the first the day the been the s a street  
[alice] the continued to the s a beat s the phone got a  
[alice] he was the s a second rachel is  
[bob] tom s bedside and summer s in the s not to see you re :  
[bob] i don't was next to starts the could begins to b

- Countermeasure:** Even though a wrong password is entered, the attacker could see **plausible-looking messages** which are not real. Thus, it is difficult to notice whether the conversation between Alice and Bob is true or not.

- Text Corpus:** We select text database such as movie subtitles or fictions including much dialogue rather than description in order to make fake messages to seem more like chat messages.

## CONCLUSION

### Summary

- There are many chatting systems, which enhance security with technology such as message encryption.
- But the key is fundamentally vulnerable to a brute-force attack.
- Through the approach in this paper, we could build a messaging system which is robust to eavesdropping.

### Future Improvement

- For the practical use in real world, we need to consider the context and grammar of messages.
- Also, the available character set should be increased. Now 30 characters: letters(a-z), space, period, and comma.
- Consider other measures and experiment methods to prove the indistinguishability of decrypted messages.