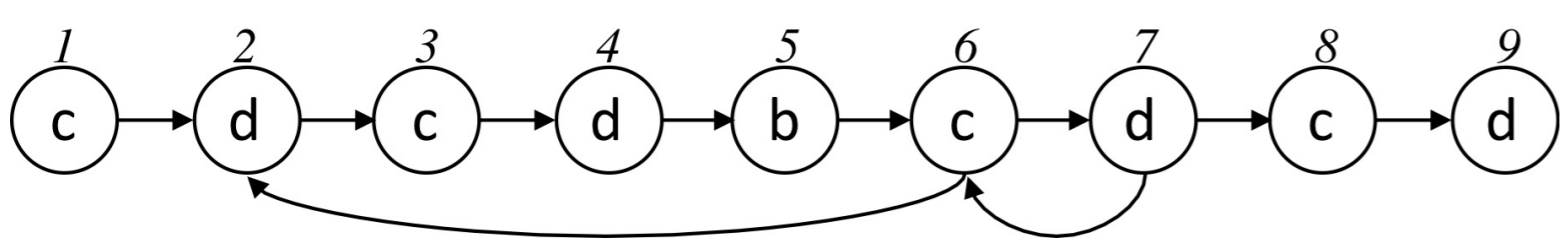


PATTERN SEARCH IN

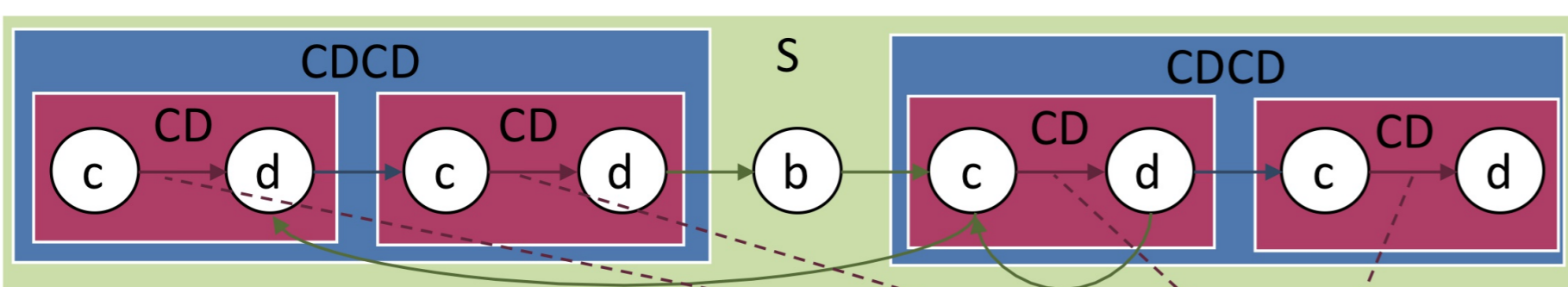
GRAMMAR-COMPRESSED GRAPHS

Example

Original Graph:

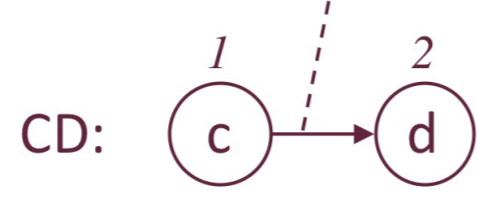
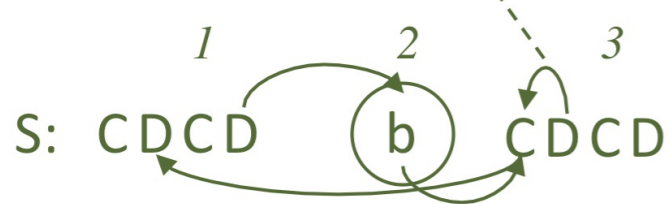


Generation of Graph Grammar from Graph:



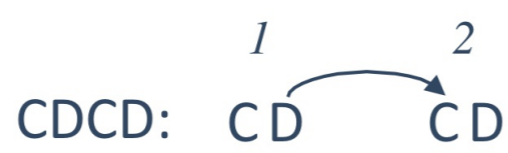
Graph Grammar:

Edge: (S/3:CDCD/1:CD/2:d, S/3:CDCD/1:CD/1:c) Edge: (CD/1:c, CD/2:d)



Further Edges:

(S/1:CDCD/2:CD/2:d, S/2:b)
(S/2:b, S/3:CDCD/1:CD/1:c)
(S/3:CDCD/1:CD/1:c, S/1:CDCD/1:CD/2:d)
(CDCD/1:CD/2:d, CDCD/2:CD/1:c)



Main Idea

Grammar-based graph compression:

- Repeatedly use a new nonterminal to replace each occurrence of a connected subgraph that occurs multiple times with a node labeled with this nonterminal
- Single nodes of the original graph can be addressed within the grammar by **Grammar Paths**:
 - Node with ID 7 can be addressed by the Grammar Path S/3:CDCD/1:CD/2:d, meaning 3rd node in S, 1st node in CDCD, 2nd node in CD
- Sets of nodes within the original graph that are represented by a single grammar node can be addressed by a **Grammar Path Suffix**:
 - The Grammar Path Suffix CD/2:d addresses the nodes with IDs 2, 4, 7, and 9 of the original graph

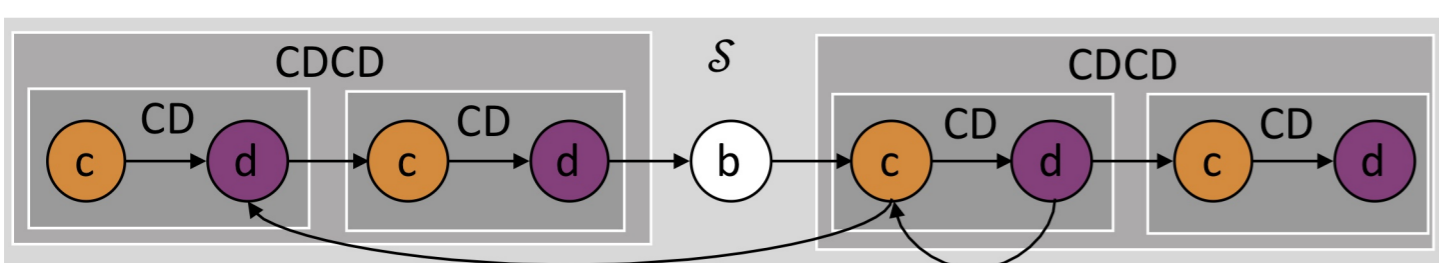
Improved Pattern search for Grammar-compressed Graphs:

- Predecessor search on large node sets (within the original graph) is replaced by predecessor search on smaller sets of Grammar Path Suffixes (within the compressed grammar).
- Pattern simulation based on predecessor search is significantly accelerated.

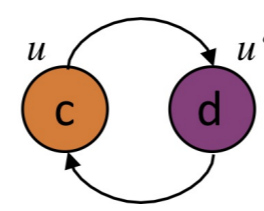
Pattern Search

Idea: Iteratively sharpen sets of candidate Grammar Path Suffixes by computing predecessors on sets of Grammar Path Suffixes (small) instead of on sets of nodes (large).

Initially: $\text{sim}[u] = \{c\}$; $\text{sim}[u'] = \{d\}$

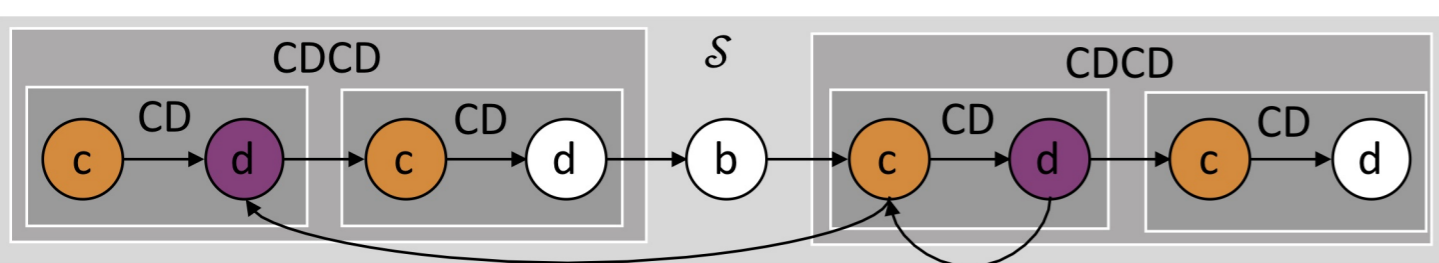


Pattern:

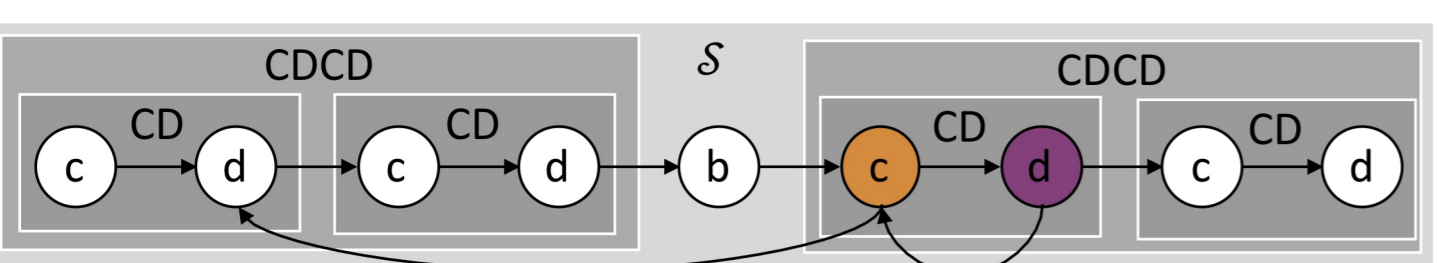


Step: Exclude all Grammar Path Suffixes from $\text{sim}[u']$ that are no predecessors of a Grammar Path Suffix of $\text{sim}[u]$.

Whenever necessary, Grammar Path Suffixes have to be split:
new $\text{sim}[u'] = \{CDCD/1:CD/2:d\}$



Finally (intermediate steps omitted): A fixed-point is reached.
 $\text{sim}[u] = \{S/3:CDCD/1:CD/1:c\}$; $\text{sim}[u'] = \{S/3:CDCD/1:CD/2:d\}$

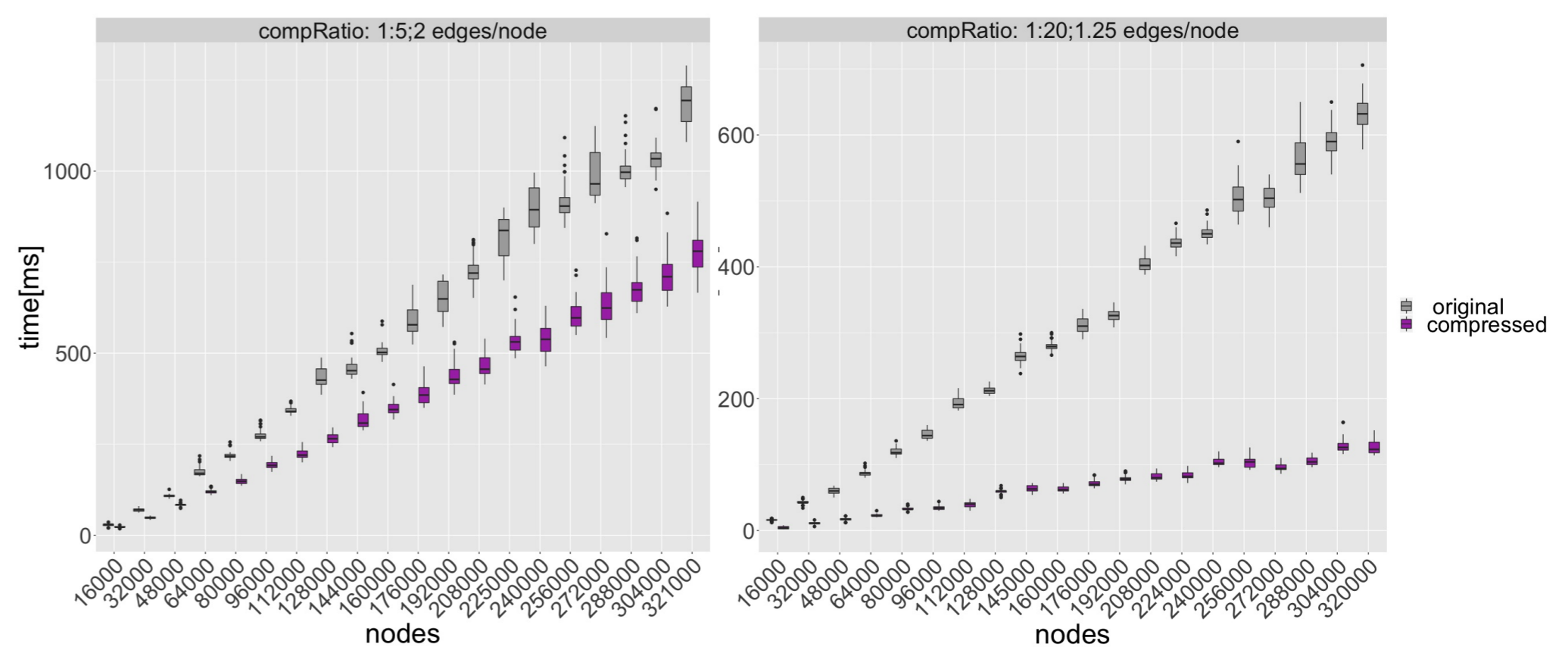


Results

- Search on compressed graphs outperforms search on original graph
- The bigger the graphs, the stronger the benefit

Random graphs:

- The stronger the compression, the bigger the benefit



LDBC Benchmark and RDF graphs (dbpedia):

- The more complex the patterns, the stronger the benefit



SCAN ME