

Node-Asynchronous Spectral Clustering on Directed Graphs

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California Institute of Technology

45th International Conference on
Acoustics, Speech and Signal Processing
(ICASSP 2020)



Caltech

Outline

- 1 Graph Signal Processing and Spectral Clustering
- 2 Asynchronous Updates on Graphs
- 3 Autonomous Spectral Clustering
- 4 Conclusion

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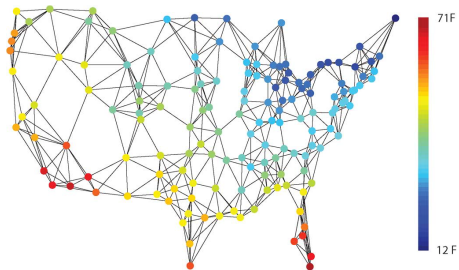
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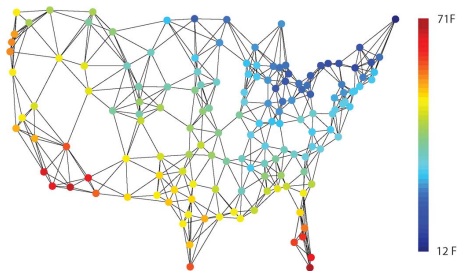
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Preliminaries



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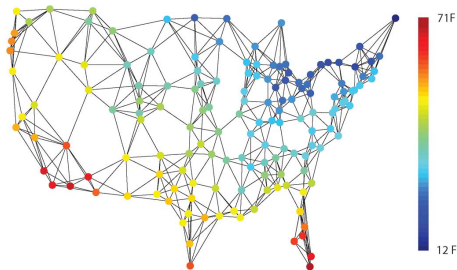
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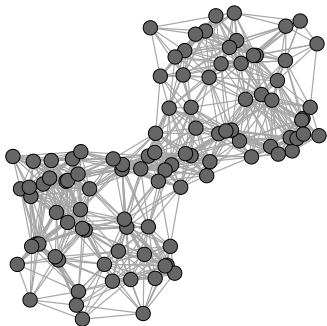
\mathbf{A} is the graph operator

Adjacency matrix¹ : \mathbf{A}
 Graph Laplacians² : \mathbf{L} , or \mathcal{L}
 Other selections³

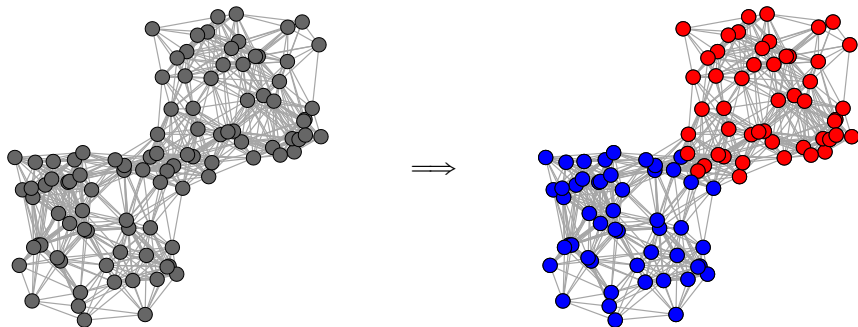
¹ Sandryhaila & Moura, "Discrete Signal Processing on Graphs," *IEEE Trans. S. P.* vol. 61, no. 7, 2013

² Shuman et al, "The emerging field of signal processing on graphs: ...," *IEEE S. P. Magazine*, vol. 30, no. 3 2013

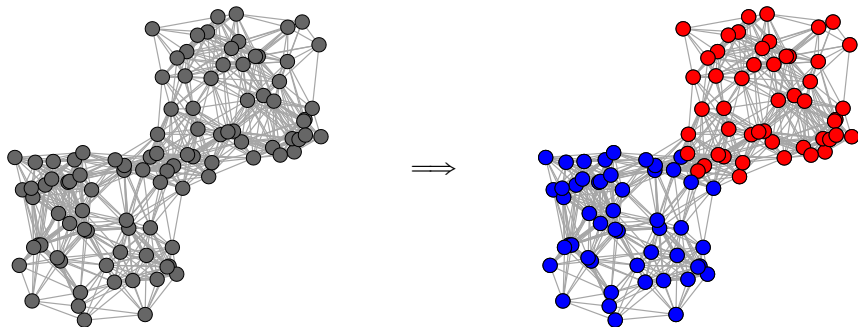
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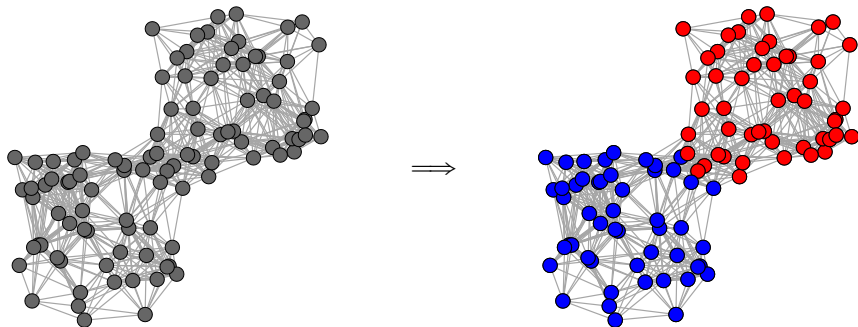


Spectral Clustering



Graph Laplacian: $\mathbf{L} = \mathbf{D} - \mathbf{A}$

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(Fiedler Value)

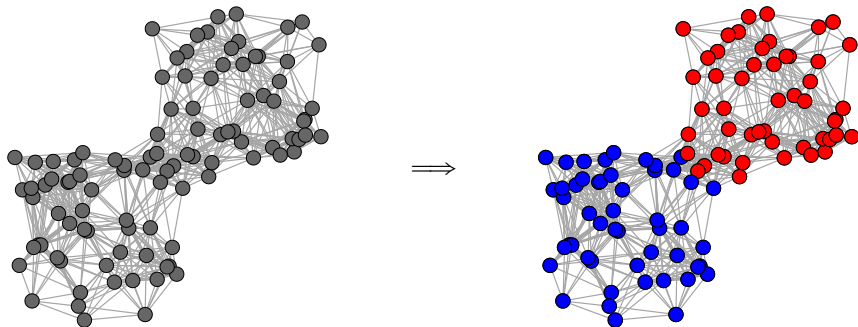
$$\mathbf{L} \mathbf{v}_2 = \lambda_2 \mathbf{v}_2$$

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Compute \mathbf{v}_2 with random asynchronous computations?

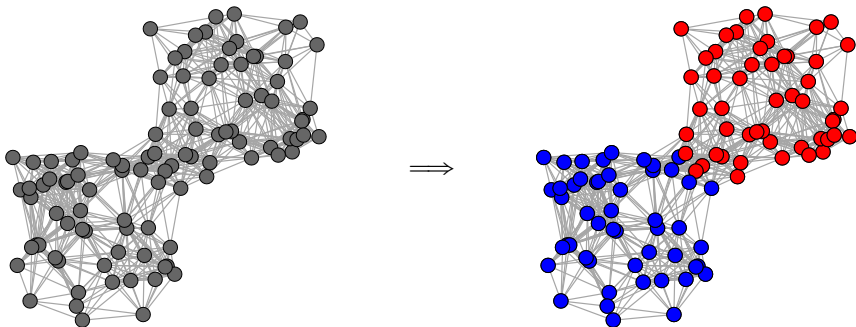
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Spectral Clustering

Here: Directed Case

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Asynchronous Fixed-Point Iterations

\mathbf{A} = Graph Operator

\mathbf{x}_k = Signal on the Graph

\mathbf{u} = Input Signal ($\mathbf{u} = \mathbf{0}$ in clustering)

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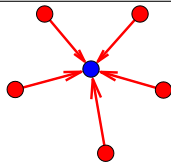
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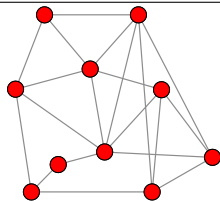
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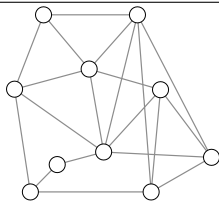
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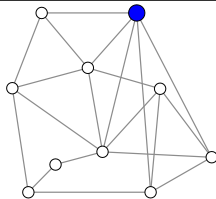
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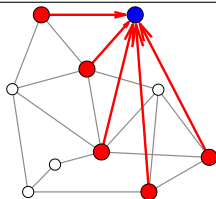
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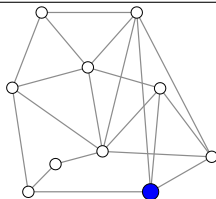
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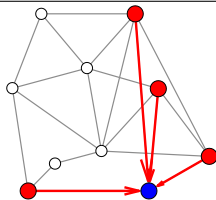
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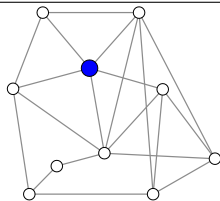
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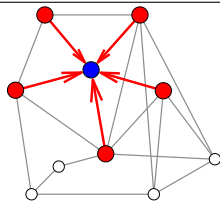
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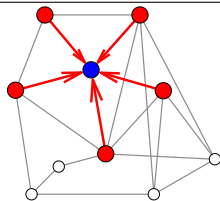
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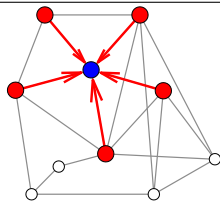
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Recurrent NN
(Hopfield Model)

$$x_k[i] = \theta(\mathbf{a}_i \mathbf{x}_{k-1} + u_i)$$

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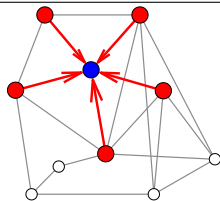
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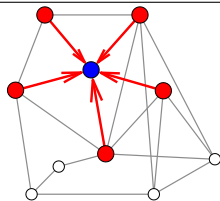
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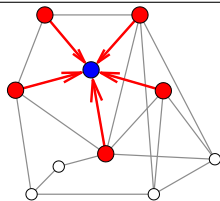
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Random Asynchronous case:

To be discussed next ...

Mean-Squared Convergence of the Updates

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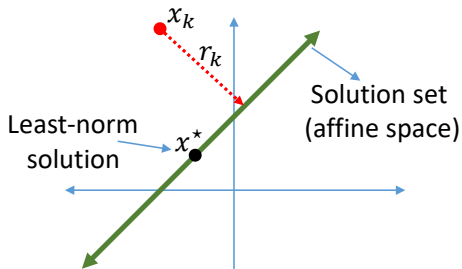
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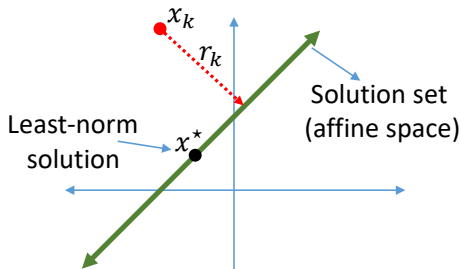


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$$\boxed{\mathbf{r}_k = \mathbf{Q}(\mathbf{x}_k - \mathbf{x}^*)}$$

\mathbf{Q} : Projection on $\text{null}^\perp(\mathbf{I} - \mathbf{A})$

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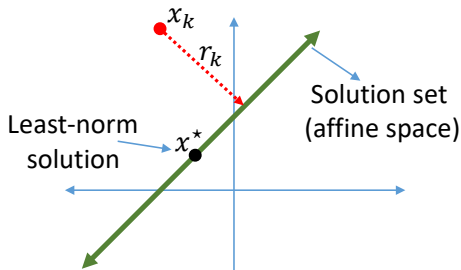
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$$\mathbf{r}_k = \mathbf{Q}(\mathbf{x}_k - \mathbf{x}^*)$$

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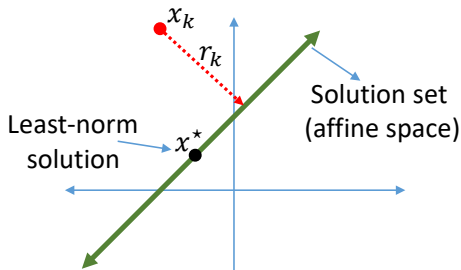
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When does it converge?

When \mathbf{A} is normal [1]

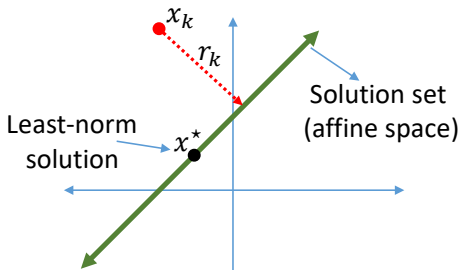
Mean-Squared Convergence of the Updates

$$x_k[i] = \begin{cases} \mathbf{a}_i \mathbf{x}_{k-1} + u_i, & \text{w.p. } p_i, \\ x_{k-1}[i], & \text{w.p. } 1-p_i. \end{cases}$$

$$\lim_{k \rightarrow \infty} \mathbf{x}_k = ?$$

Where does it converge?

$$\mathbf{x} = \mathbf{A}\mathbf{x} + \mathbf{u} \Rightarrow \boxed{\mathbf{u} \in \text{range}(\mathbf{I} - \mathbf{A})}$$

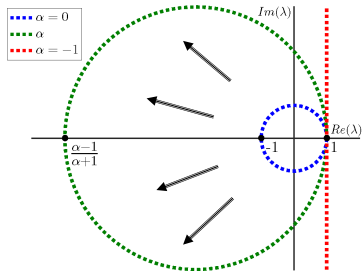


$$\boxed{\mathbf{r}_k = \mathbf{Q}(\mathbf{x}_k - \mathbf{x}^*)}$$

\mathbf{Q} : Projection on $\text{null}^\perp(\mathbf{I} - \mathbf{A})$

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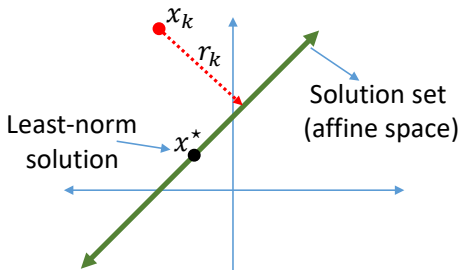
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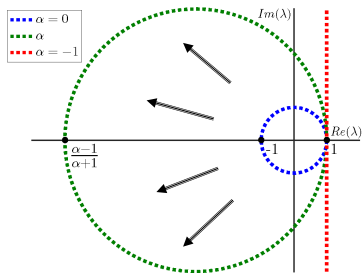


$$\mathbf{r}_k = \mathbf{Q}(\mathbf{x}_k - \mathbf{x}^*)$$

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When does it converge?

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$$\lim_{k \rightarrow \infty} \mathbb{E}[\|\mathbf{r}_k\|_2^2] = 0$$

Mean-Squared Convergence of the Updates - Cont.

Theorem (The necessary and sufficient condition)

Mean-Squared Convergence of the Updates - Cont.

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$$\lim_{k \rightarrow \infty} \mathbb{E}[\|\mathbf{r}_k\|_2^2] = 0 \quad \iff \quad \rho(\mathbf{\Theta S}) < 1$$

Mean-Squared Convergence of the Updates - Cont.

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$$\mathbf{S} = \bar{\mathbf{A}}^* \otimes \bar{\mathbf{A}} + \left((\mathbf{I} - \mathbf{P}) \otimes \mathbf{P} \right) \mathbf{J} \left((\mathbf{A}^* - \mathbf{I}) \otimes (\mathbf{A} - \mathbf{I}) \right),$$

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- ✓ Robust to input noise [1, 2]
- ✓ Can compute eigenvectors (to be discussed next)

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Computing Eigenvectors

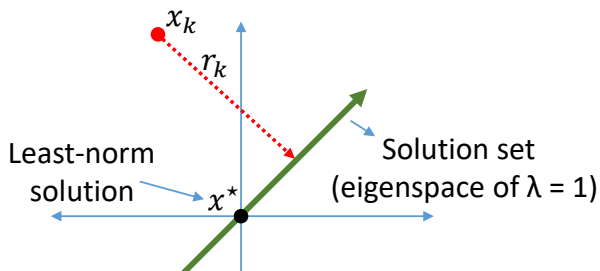
$$\mathbf{u} = \mathbf{0} \quad \Longrightarrow \quad x_k[i] = \begin{cases} \mathbf{a}_i \mathbf{x}_{k-1}, & \text{w.p. } p_i, \\ x_{k-1}[i], & \text{w.p. } 1-p_i. \end{cases}$$

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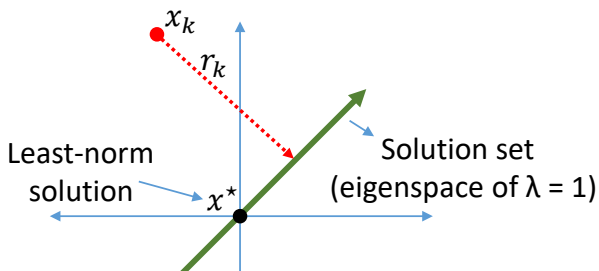
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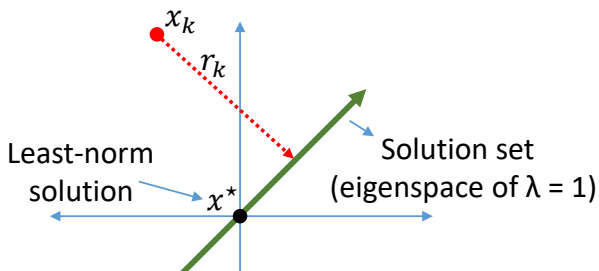
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$$\lim_{k \rightarrow \infty} \mathbb{E}[\|\mathbf{r}_k\|_2^2] = \mathbf{0}$$



\mathbf{x}_k converges to
an eigenvector of $\lambda = 1$

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Use of Graph Polynomials

A
Asynchronous

Use of Graph Polynomials

$$\lambda = 1$$

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$$\xrightarrow[\text{Asynchronous}]{\mathbf{A}}$$

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$$H(\mathbf{A}) = \sum_{k=0}^L h_k \mathbf{A}^k$$

Asynchronous

Use of Graph Polynomials

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Asynchronous

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(Fiedler Vector)

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$$H(\mathbf{A}) \sim \mathbf{A}$$

L^{th} order $\Rightarrow L$ -hop neighborhood

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Linear Programming:

$$\begin{aligned} \max_{\mathbf{h}} \quad & c \quad \text{s.t.} \quad \phi_2 \mathbf{h} = 1 \\ & c \geq 0 \quad |\bar{\Phi} \mathbf{h}| \leq (1 - c) \mathbf{1}_{N-1} \end{aligned}$$

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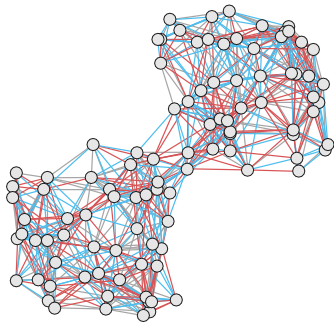
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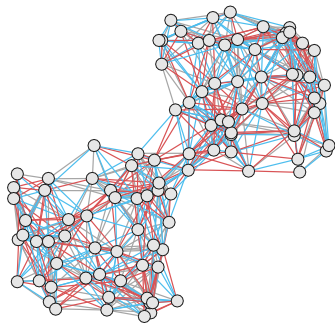
$L = 2$ works in practice

A Numerical Application



A Numerical Application

(L)



$$0 = \lambda_1 < |\lambda_2| \leq \dots \leq |\lambda_N|$$

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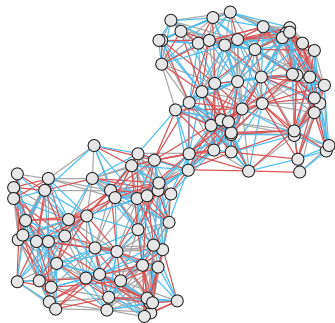
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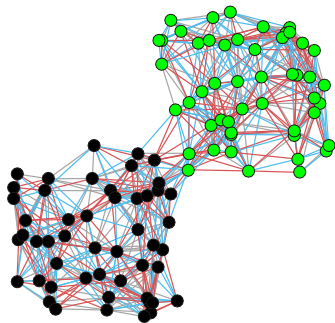
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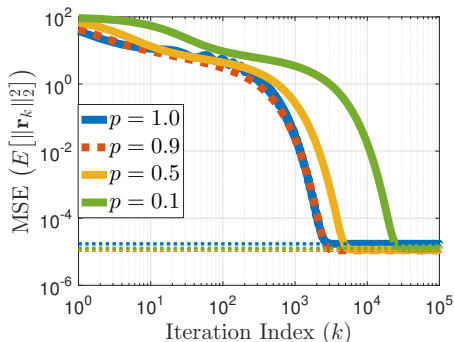
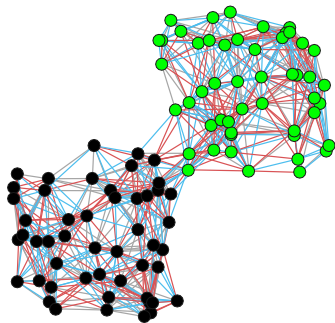
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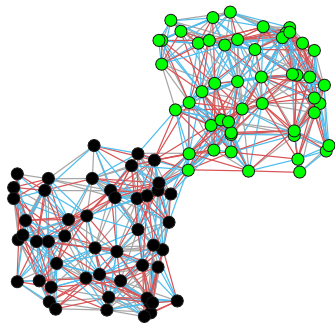
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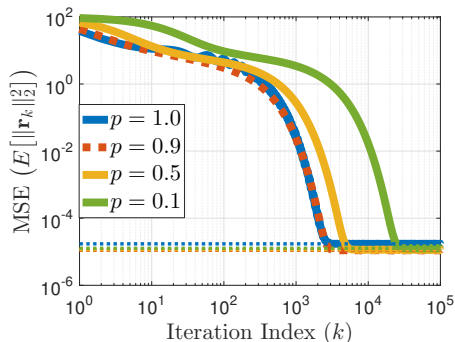
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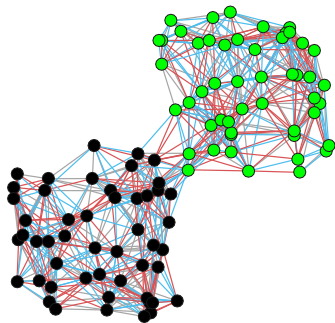
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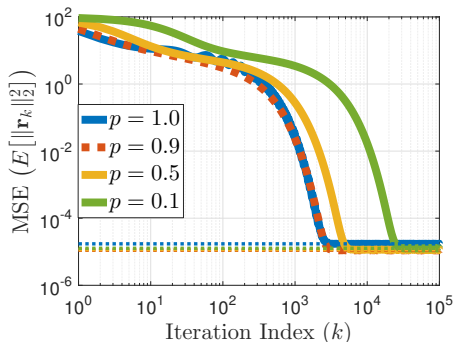
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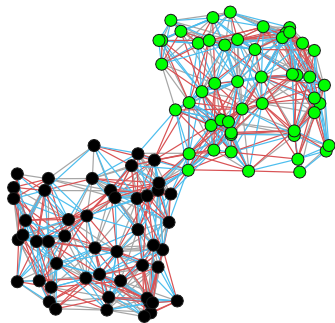
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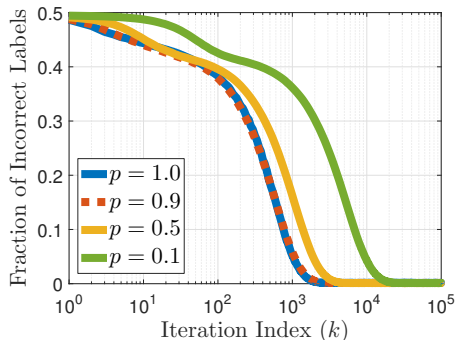
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Thank you!

New Horizons ...

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Randomized numerical linear algebra

- *Emerging field!*
- *Approximate v.s. Exact*

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² Martinsson & Tropp, "Randomized Numerical Linear Algebra: Foundations and Algorithms," *arXiv*, 2020

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Randomization opens up a new dimension!

How can we exploit it for better?

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