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An Optimal Symmetric Threshold Strategy for Remote Estimation over the Collision Channel

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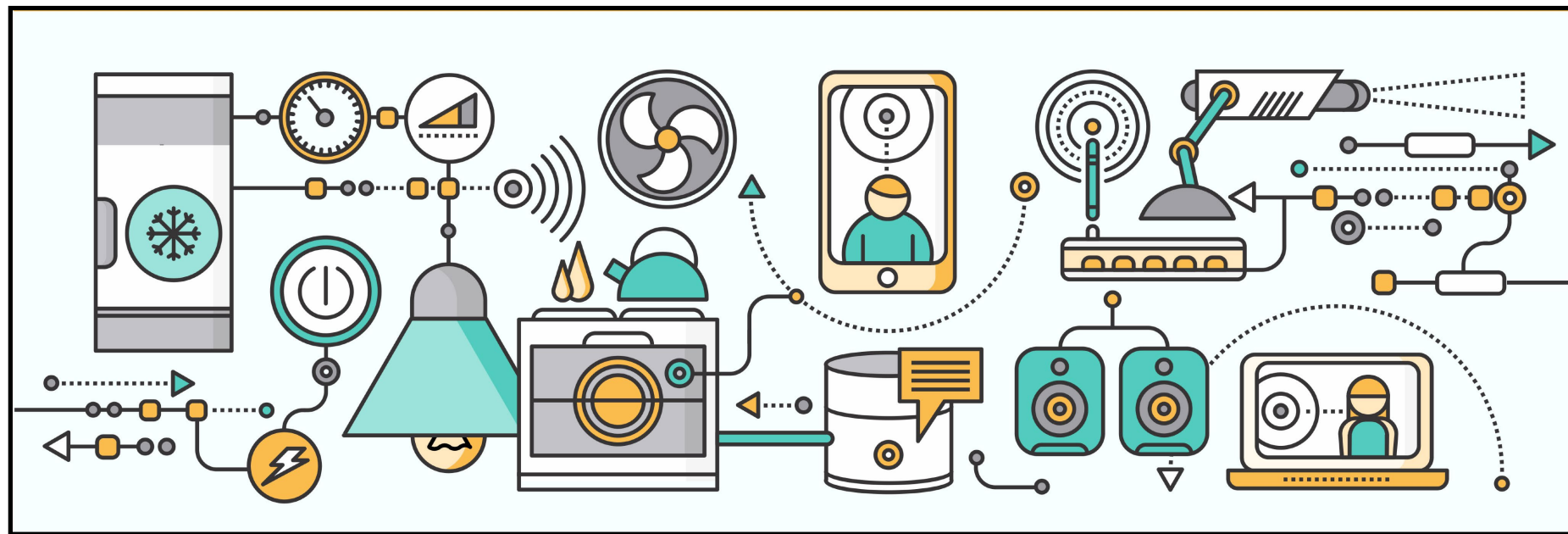
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Introduction

Motivation: Modern large-scale wireless distributed systems share limited bandwidth and should operate with minimal delay.



Channel Model: limited bandwidth \rightarrow subset of users communicate at any given time, collisions can occur

Introduction

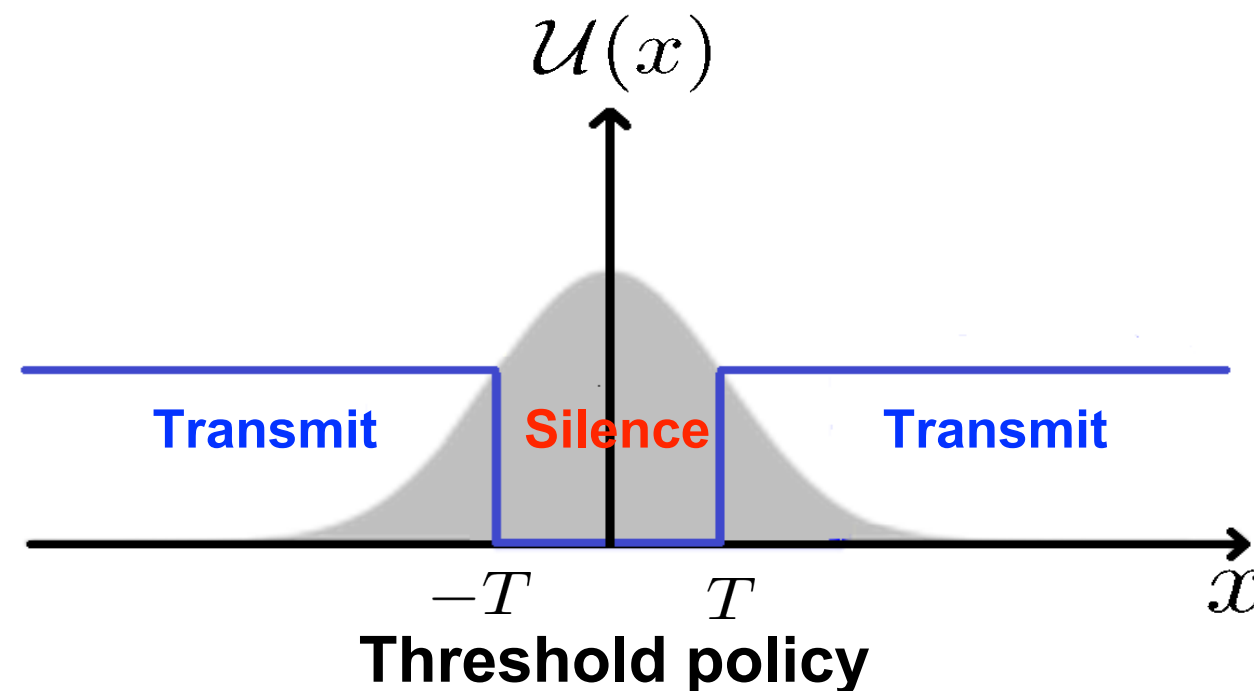
Prior work:

- systems with a single sensor
 - a limited number of transmissions [Imer&Basar'10]
 - costly communication [Lipsa&Martins'11]
 - energy harvesting sensor [Nayyar et. al'13], [Leong et. al'18]
- systems with multiple sensors
 - discrete random variables [Vasconcelos&Martins'19]
 - shared channel [Ding et. al'17]
- systems with feedback
 - based on the feedback [Huang&Dey'06], [Dogandzic&Qiu'08]
 - based on its observation and feedback quality [Michelusi&Mittra'15]

Introduction

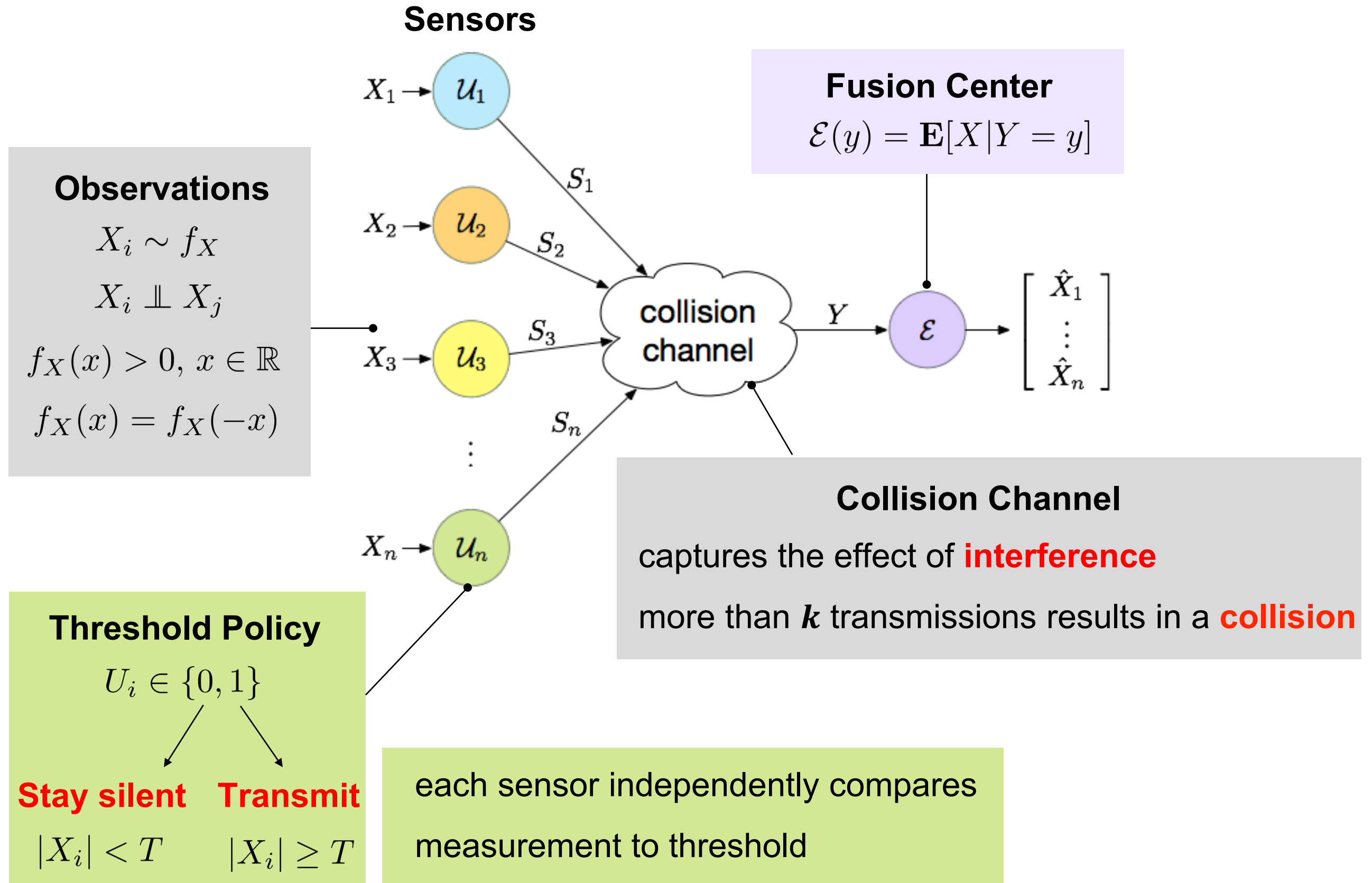
Most relevant work:

- Vasconcelos & Martins, Optimal estimation over the collision channel, IEEE TAC 2017
 - **asymmetric** threshold strategies under one-shot transmission
- Vasconcelos & Mitra, Observation-driven scheduling for remote estimation of two gaussian random variables, IEEE TCNS 2019
 - sensors with largest magnitude observations should be transmitted
 - centralized system \rightarrow lower bound of new scheme



Our Goal: Design a **decentralized threshold** strategy that minimizes the mean-squared error

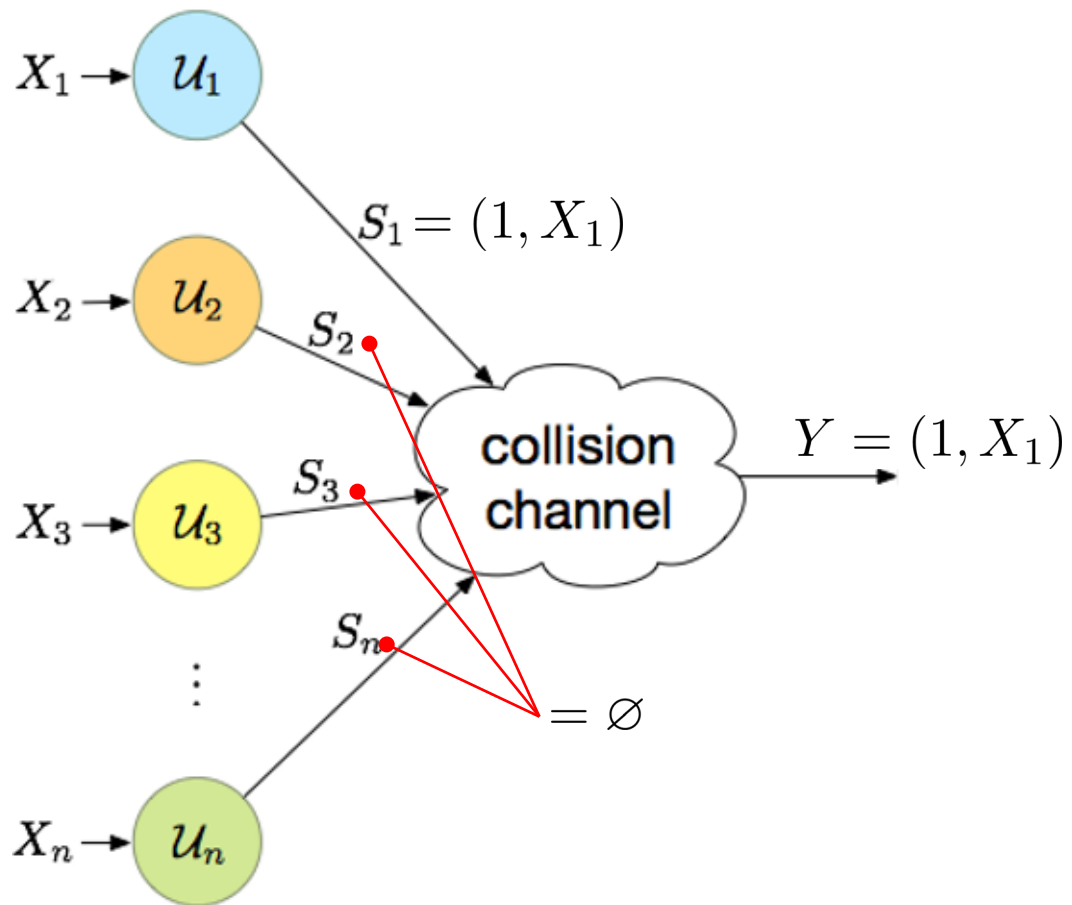
Basic framework



Collision channel

Reliable transmission

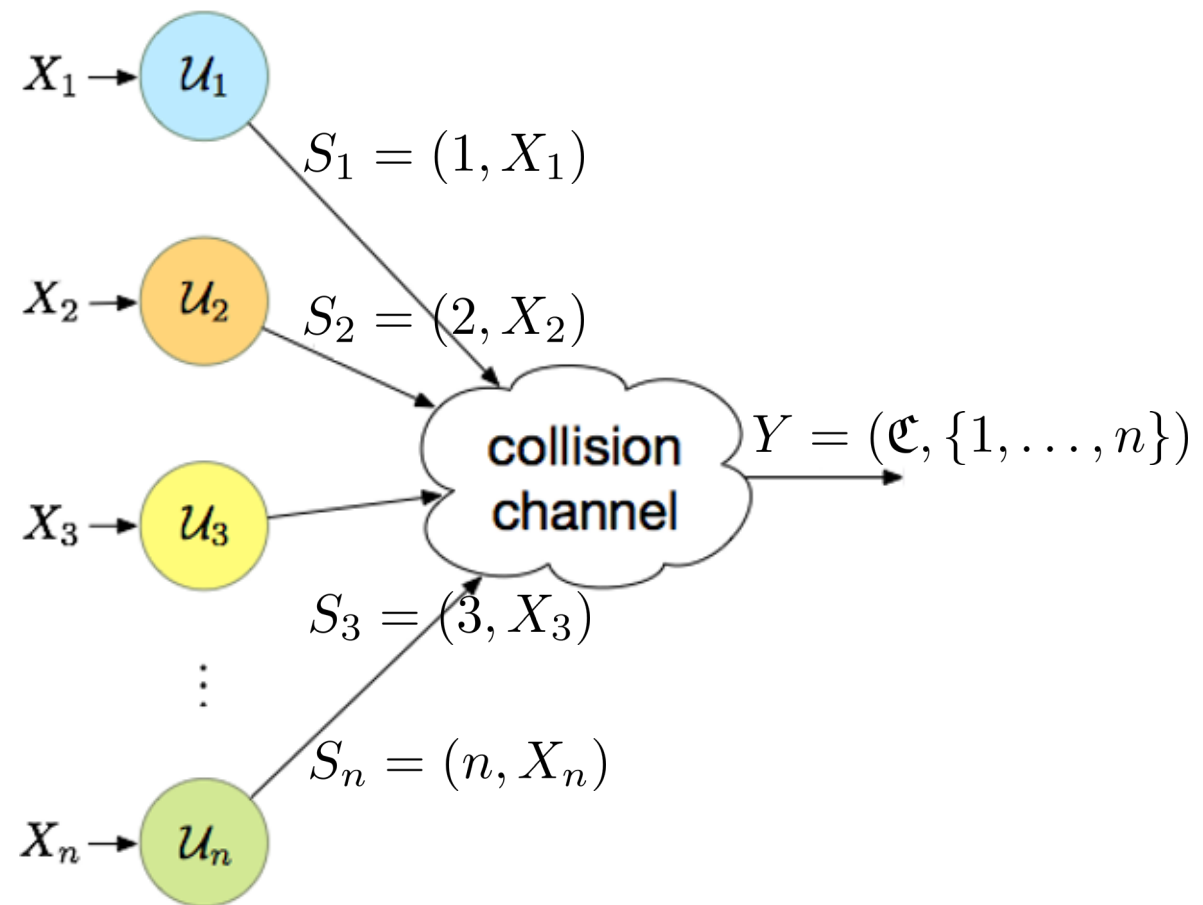
$$\sum_{i=1}^n U_i \leq k$$



$$Y = \{(i, X_i) | i \in \mathcal{S}\}$$

Collision

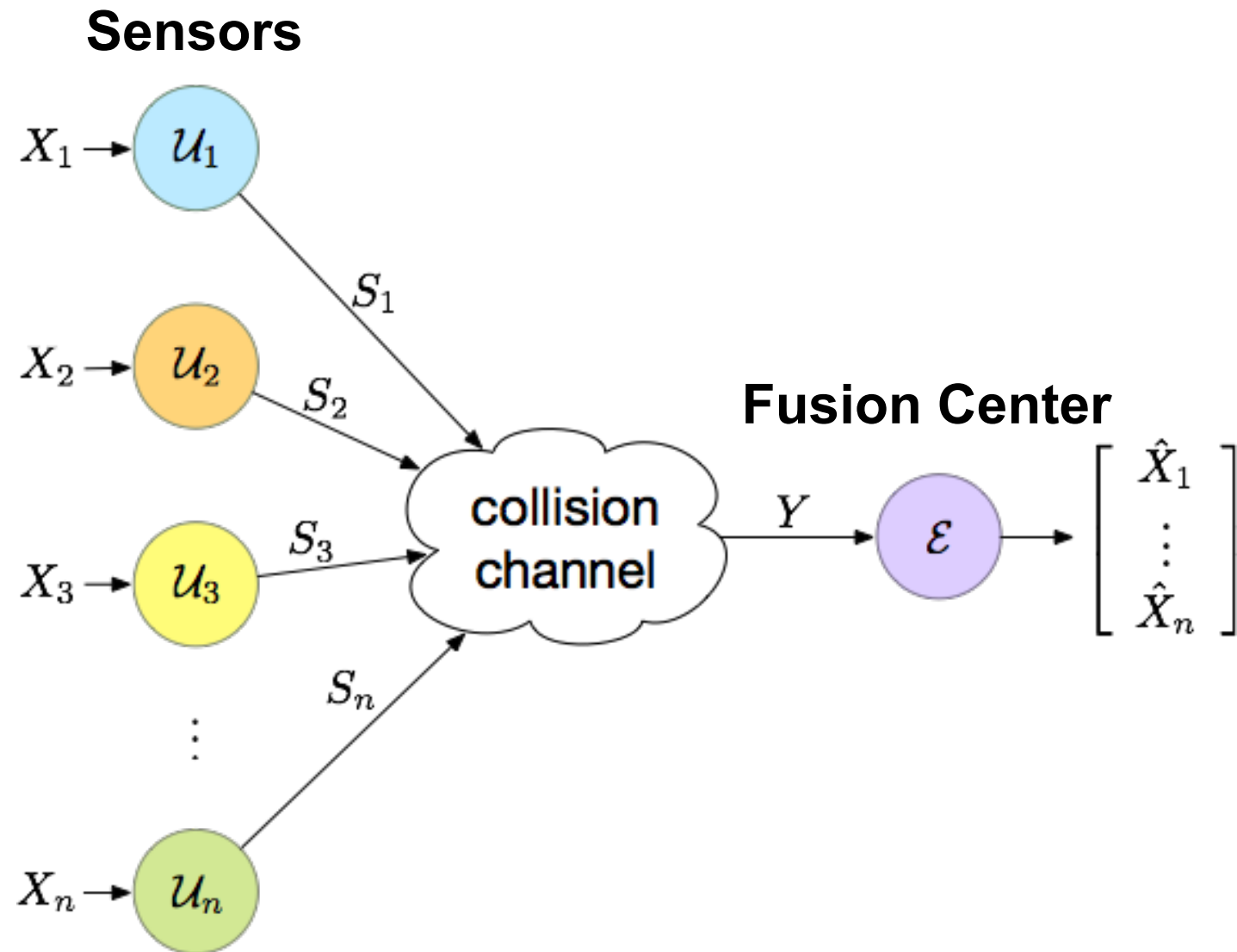
$$\sum_{i=1}^n U_i \geq k + 1$$



$$Y = (\mathfrak{C}, \mathcal{S})$$

$$\mathcal{S} = \{i | U_i = 1\}$$

Estimation over the collision channel



Problem

Find an optimal threshold T^* that minimizes MSE

MSE

$$\mathcal{J}_{n,k}(T) \triangleq \frac{1}{n} \mathbf{E} \left[\sum_{i=1}^n \left(X_i - \hat{X}_i \right)^2 \right]$$

Main results

Estimator

$$\hat{X}_i = \begin{cases} X_i \mathbf{1}(|X_i| \geq T) & \text{if } \sum_{i=1}^n U_i \leq k \\ 0 & \text{if } \sum_{i=1}^n U_i \geq k + 1 \end{cases}$$

Cost Function

$$\mathcal{J}(T) = \mathbf{E} [X^2] - \mathbf{E} [X^2 \mathbf{1}(|X| \geq T)] F_{n,k}(T)$$

cdf for an order statistic



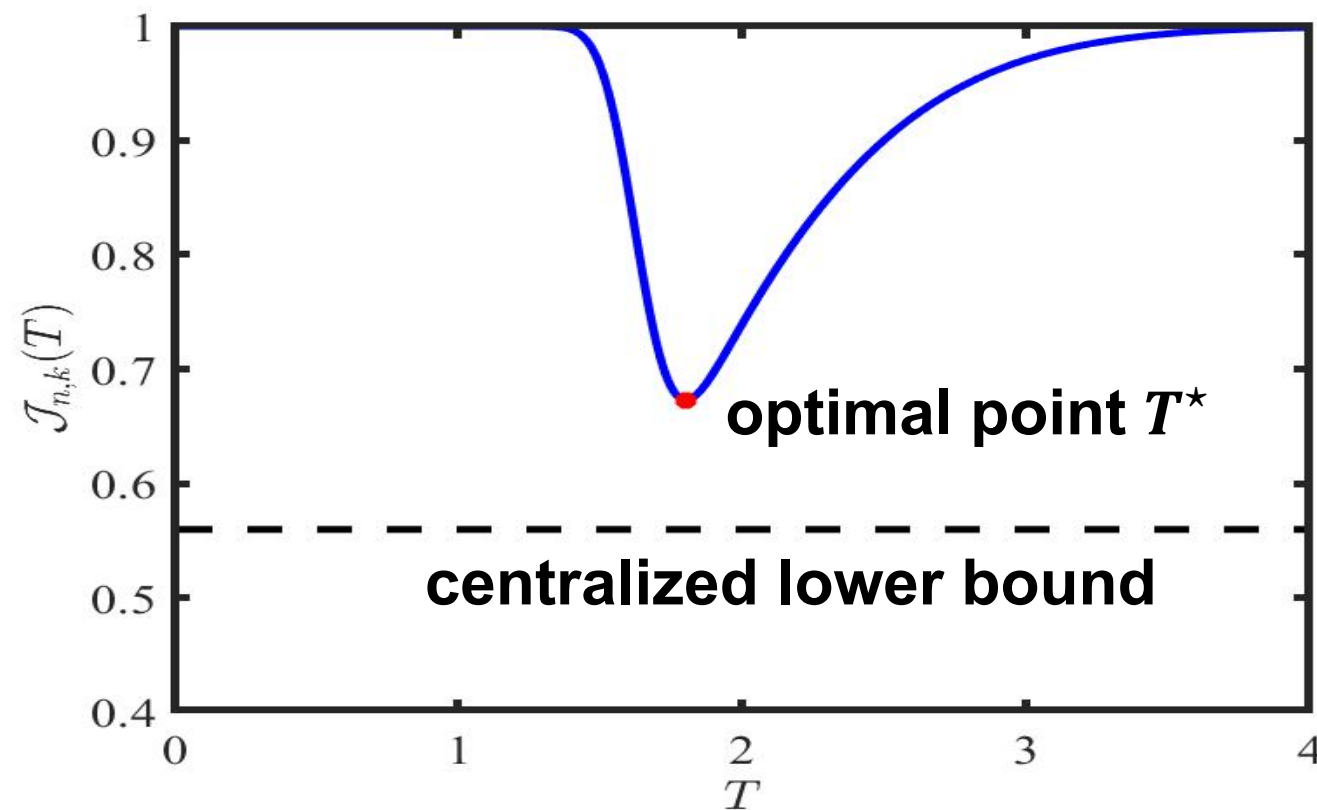
$$F_{n,k}(T) = \sum_{j=0}^{k-1} \binom{n-1}{j} (F_Z(T))^{n-j-1} (1 - F_Z(T))^j, \quad Z = |X|$$

results based on **order statistics**, hold for **arbitrary** symmetric distributions

Main results

Theorem

There exists a unique threshold: $T^* = \operatorname{argmin}_{T \in \mathbb{R}} \mathcal{J}_{n,k}(T)$



$$n = 200, k = 20$$

$$X_i \sim \mathcal{N}(0, 1)$$

Sketch of Proof:

Step 1: Compute derivative

Step 2: Prove there is a unique local minimum

Sketch of Proof

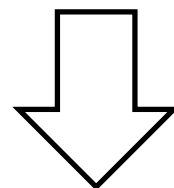
Step 1

$$\mathcal{J}'(T) = f_{n,k}(T) \underbrace{\left[\frac{T^2 f_Z(T) F_{n,k}(T)}{f_{n,k}(T)} - \mathbf{E}[X^2 \mathbf{1}(|X| \geq T)] \right]}_{\doteq h(T)}$$

Step 2

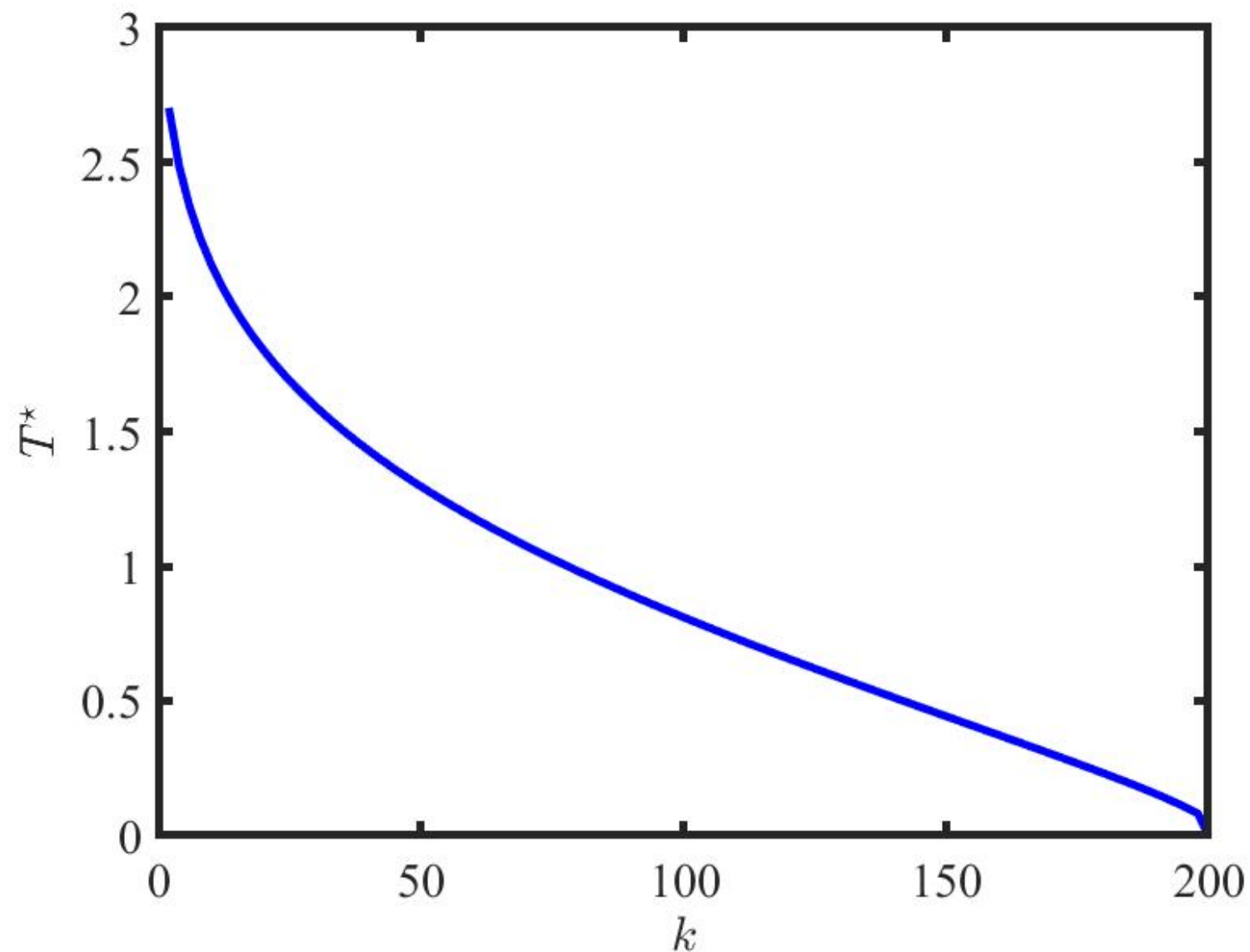
$$1) f_{n,k}(T) = k \binom{n-1}{k} (F_Z(T))^{n-k-1} (1 - F_Z(T))^{k-1} f_Z(T) > 0$$

2) $h(T)$ is a strictly increasing continuous function of T for $T \in [0, +\infty)$
 $h(0^+) < 0$ and $h(\infty) > 0$



There is a unique threshold such that $\mathcal{J}'(T^*) = 0$ and $T^* = \operatorname{argmin}_T \mathcal{J}(T)$

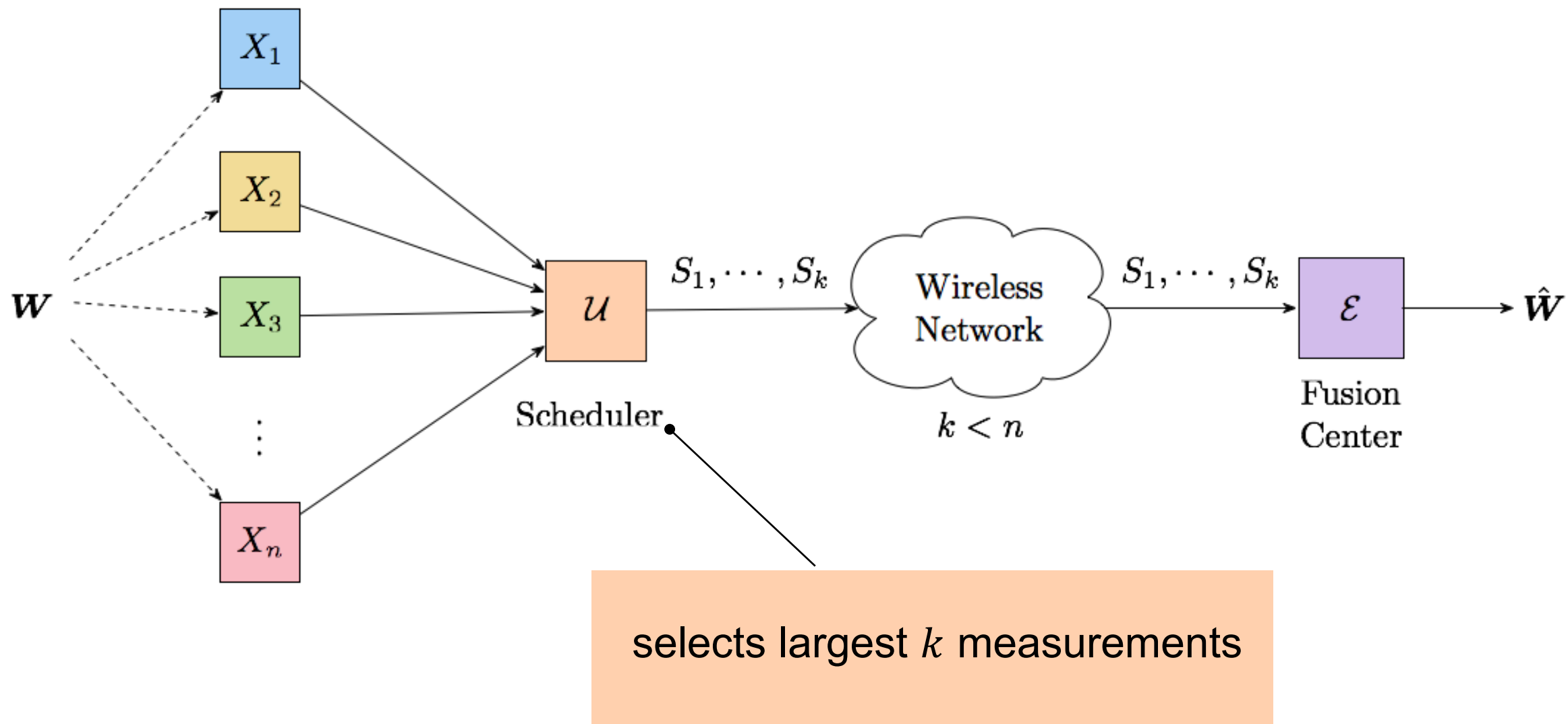
Optimal threshold vs channel capacity



$n = 200$
 $X_i \sim \mathcal{N}(0, 1)$

- Optimal threshold T^* decreases as collision capacity k increases
- When channel supports all packets, optimal threshold is zero

Centralized lower bound



Optimal performance of this system is a **lower bound** to the **decentralized problem**

Centralized lower bound

$$\mathcal{J}_L = \frac{1}{n} \sum_{i=k+1}^n \mathbf{E} \left[Z_{(i)}^2 \right]$$

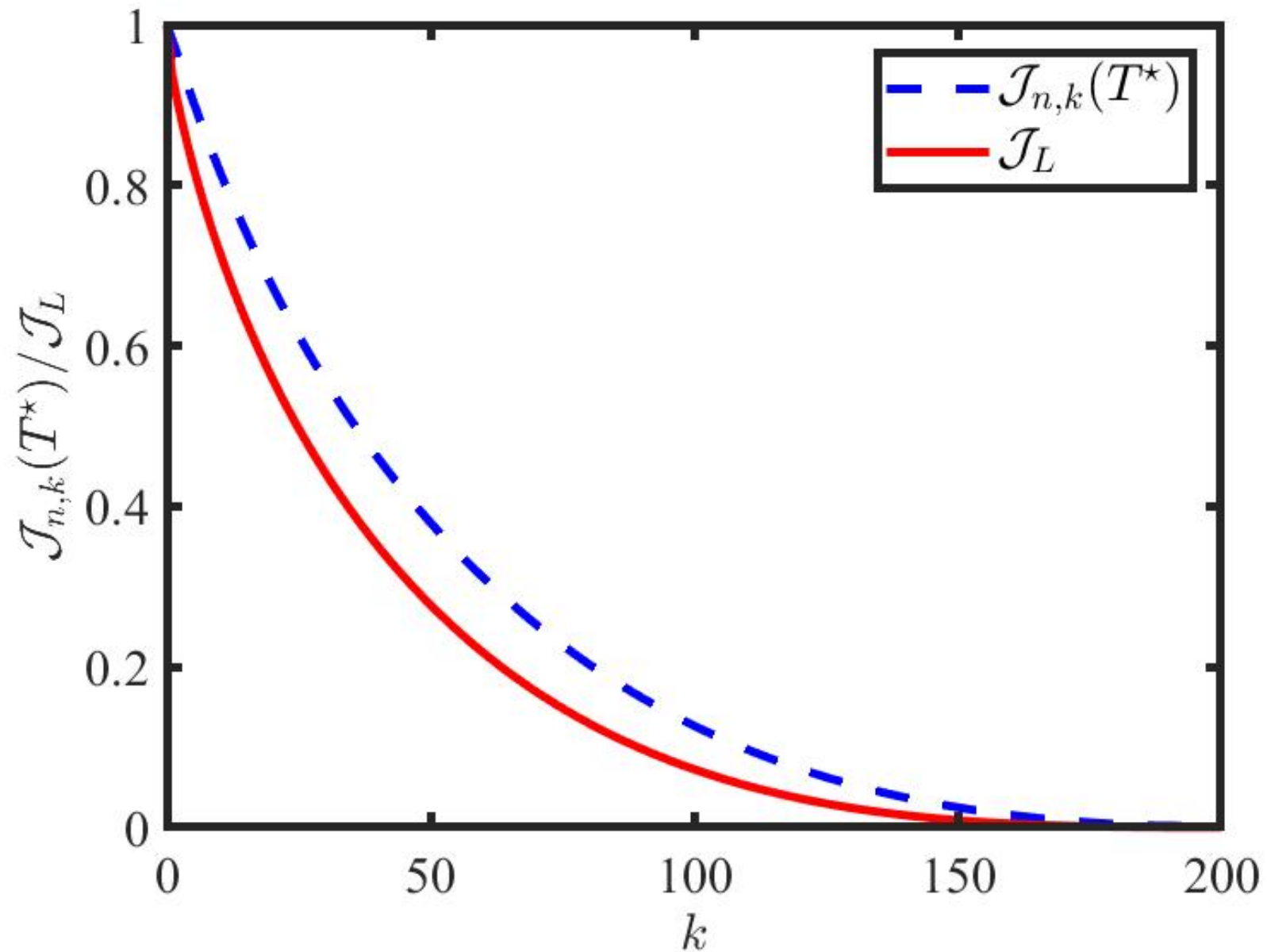
$Z_{(i)}$ denotes the i -th largest value in $\{|X_i|\}_{i=1}^n$

$$\mathbf{E} \left[Z_{(t)}^2 \right] = \frac{\int_0^\infty z^2 F_Z(z)^{n-t} (1 - F_Z(z))^{t-1} f_Z(z) dz}{\mathbf{B}(n - t + 1, t)}$$

$Z = |X|$

beta function

Decentralized vs Centralized



$n = 200$
 $X_i \sim \mathcal{N}(0, 1)$

- As capacity k increases, MSE decreases
- Decentralized scheme performs close to centralized scheme

Summary & ongoing work

1. Summary:

Analyzed the optimal symmetric threshold policy for remote estimation problem

Provided existence of **unique** optimal threshold

2. Ongoing work:

Consider local communication

Analysis for **unknown** (or imprecise) probabilistic models

- Quantile-based decentralized scheme

Thank you!