

Traffic-Aware Association in HetNet

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Problem Statement

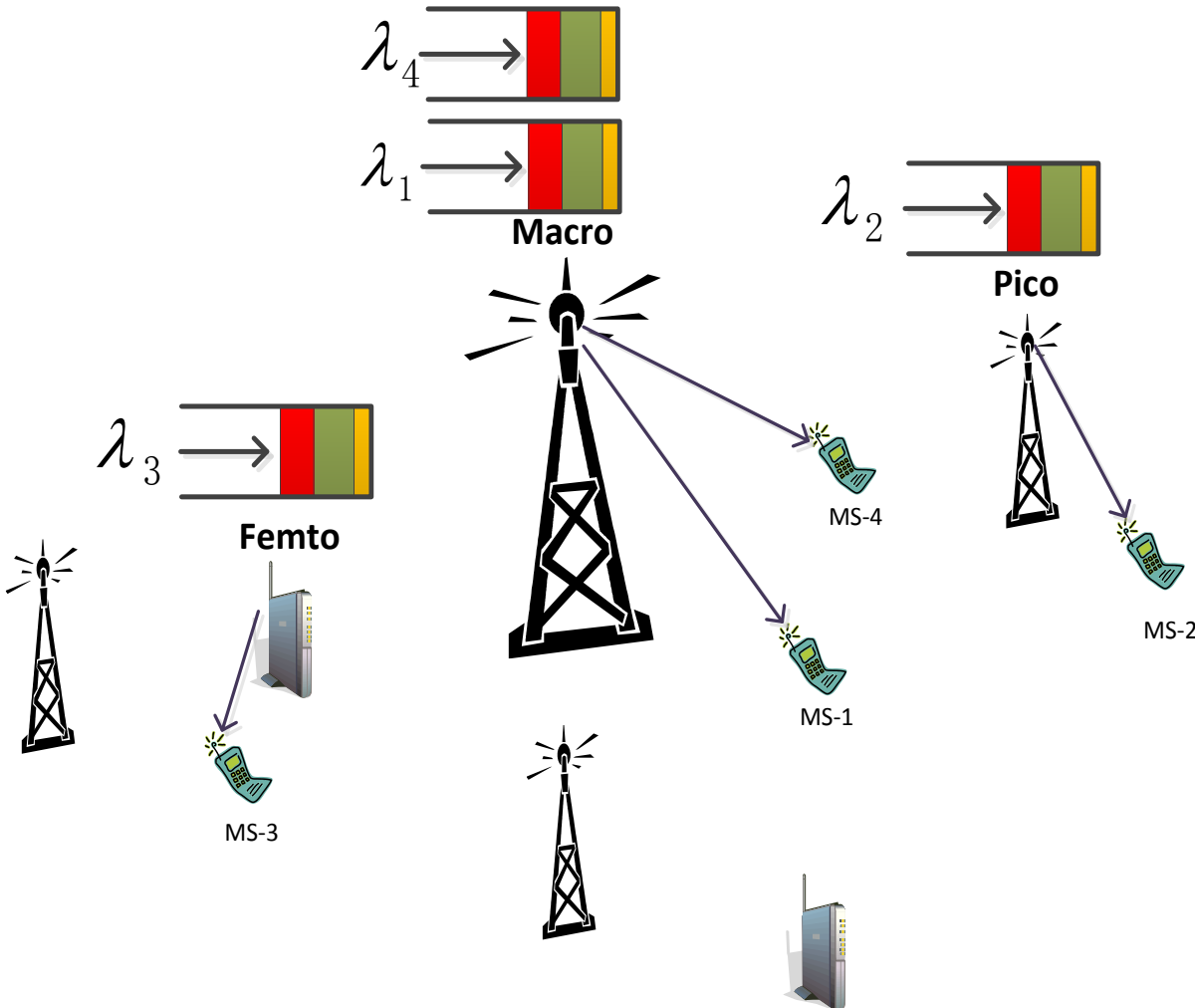
□ Conventional User Association in HetNet:

- Max-DL-SINR association as in 2G/3G/4G networks
- Cell Range Expansion [Sesia et al'11, Andrews et al'14]
- Joint optimization of association and resource allocation
 - Assuming Best-Effort (BE) traffic
 - Performance metrics
 - the sum of log-rate [Fooladivanda et al'13, Ye et al'13, Deb et al'14]
 - number of admitted users [Li et al'12]
 - sum of the inverse of the per-user throughput [Chen'11]

□ Our proposed association rule:

- ✓ Model the DL QoS traffic explicitly
- ✓ Optimize the network-wide packet delay performance

HetNet with QoS Traffic



Definitions:

B : system BW in Hz

P_n : txmn pwr of BS- n

λ_k : pkt arrival rate of MS- k

L_k : avg pkt length in bits of MS- k

$h_{k,n}$: channel btw MS- k and BS- n

$x_{k,n}$: assoc indx of MS- k to BS- n

$y_{k,n}$: resource alloc

Assumptions:

i.i.d. exp inter-arrival times

i.i.d. exp packet lengths

BS always ON

full freq. reuse

Problem Formulation

- Avg rate of MS-k associated with BS-n occupying all resources:

$$R_{k,n} = B \log \left(1 + \frac{P_n |h_{k,n}|^2}{\sum_{l=1, l \neq n}^N P_l |h_{k,l}|^2 + \sigma^2} \right)$$

- Service time for the traffic towards MS-k is i.i.d. exp with mean:

$$t_{k,n} = \frac{L_k}{y_{k,n} R_{k,n}} = \frac{1}{y_{k,n} r_{k,n}}$$

- Avg delay for a packet in the M/M/1 queue:



$$\tau_{k,n} = \frac{1}{1/t_{k,n} - \lambda_k} = \frac{1}{y_{k,n} r_{k,n} - \lambda_k}$$

Problem Formulation



- Optimal association rule to minimize the average pkt delay across the whole network:

$$\begin{aligned} & \underset{x_{k,n}, y_{k,n}}{\text{minimize}} && \frac{1}{\sum_{k=1}^K \lambda_k} \sum_{n=1}^N \sum_{k=1}^K \frac{x_{k,n} \lambda_k}{y_{k,n} r_{k,n} - x_{k,n} \lambda_k} \\ & \text{subject to} && \sum_{n=1}^N x_{k,n} = 1.0, \forall k = 1, \dots, K \\ & && \sum_{k=1}^K y_{k,n} \leq 1.0, \forall n = 1, \dots, N \\ & && x_{k,n} \in \{0, 1\}, \forall k = 1, \dots, K, n = 1, \dots, N \\ & P1 && y_{k,n} > \frac{x_{k,n} \lambda_k}{r_{k,n}}, \forall k = 1, \dots, K, n = 1, \dots, N. \end{aligned}$$

- A classic knapsack problem is NP-hard [Bertsekas'99]
- **We try to find low-complexity approx. solutions!**

Step 1: Opt Resource Alloc

□ Proposition 1:

At BS-n, given user association, when the following feasibility condition is met:

$$\sum_{k=1}^K \frac{x_{k,n} \lambda_k}{r_{k,n}} < 1.0,$$

the optimal resource allocation minimizing the average pkt delay is as follows:

$$y_{k,n} = \frac{x_{k,n} \lambda_k}{r_{k,n}} + \frac{1 - \sum_{u=1}^K \frac{x_{u,n} \lambda_u}{r_{u,n}}}{\sum_{u=1}^K \sqrt{\frac{x_{u,n} \lambda_u}{r_{u,n}}}} \sqrt{\frac{x_{k,n} \lambda_k}{r_{k,n}}}$$

$$y_{k,n} = \frac{1}{\sum_{u=1}^K \frac{x_{u,n} \lambda_u}{r_{u,n}}} \frac{x_{k,n} \lambda_k}{r_{k,n}}$$

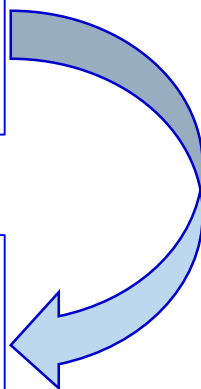
linear approx

Simplified Problem

□ Optimal Association:

minimize $\sum_{k=1}^K \sum_{n=1}^N \frac{x_{k,n} \lambda_k}{r_{k,n}} \frac{1}{1 - \sum_{u=1}^K \frac{x_{u,n} \lambda_u}{r_{u,n}}}$
 subject to $\sum_{n=1}^N x_{k,n} = 1, \forall k = 1, \dots, K$
 $x_{k,n} \in \{0, 1\}, \forall k = 1, \dots, K, n = 1, \dots, N$
P2 $\sum_{k=1}^K \frac{x_{k,n} \lambda_k}{r_{k,n}} < 1, \forall n = 1, \dots, N$

minimize $f(\{x_{u,n}\}) = \sum_{k=1}^K f_k(\{x_{u,n}\})$
 subject to $\sum_{n=1}^N x_{k,n} = 1, \forall k = 1, \dots, K$
 $x_{k,n} \in [0, 1], \forall k = 1, \dots, K, n = 1, \dots, N$
P3 $\sum_{k=1}^K \frac{x_{k,n} \lambda_k}{r_{k,n}} < 1, \forall n = 1, \dots, N$



**convex
relaxation**

$$f_k(\{x_{u,n}\}) = \sum_{n=1}^N \frac{x_{k,n} \lambda_k}{r_{k,n}} \frac{1}{1 - \sum_{u=1}^K \frac{x_{u,n} \lambda_u}{r_{u,n}}}$$

Step 2: Opt Assocication

□ Proposition 2

For MS-k, given others' association, the optimal association of MS-k minimizing the objective function in problem P3 is as follows:

$$x_{k,n} = \max \left\{ 0, \frac{r_{k,n} \delta_{k,n}}{\lambda_k} - \alpha \sqrt{\frac{r_{k,n}}{\lambda_k}} \right\}$$

chosen such that $\sum x_{k,n} = 1$

$$\delta_{k,n} := 1 - \sum_{u=1, u \neq k}^K \frac{x_{u,n} \lambda_u}{r_{u,n}}$$

available load in BS-n for MS-k

Traffic-Aware Assoc Algorithm



1. Feasible start: (LP)

$$\begin{aligned} & \underset{x_{k,n}, s}{\text{minimize}} && s \\ & \text{subject to} && \sum_{n=1}^N x_{k,n} = 1, \forall k \\ & && x_{k,n} \in [0, 1], \forall k, n \\ & && \sum_{k=1}^K x_{k,n} \frac{\lambda_k}{r_{k,n}} - 1 \leq s, \forall n. \end{aligned}$$

2. Iterations:

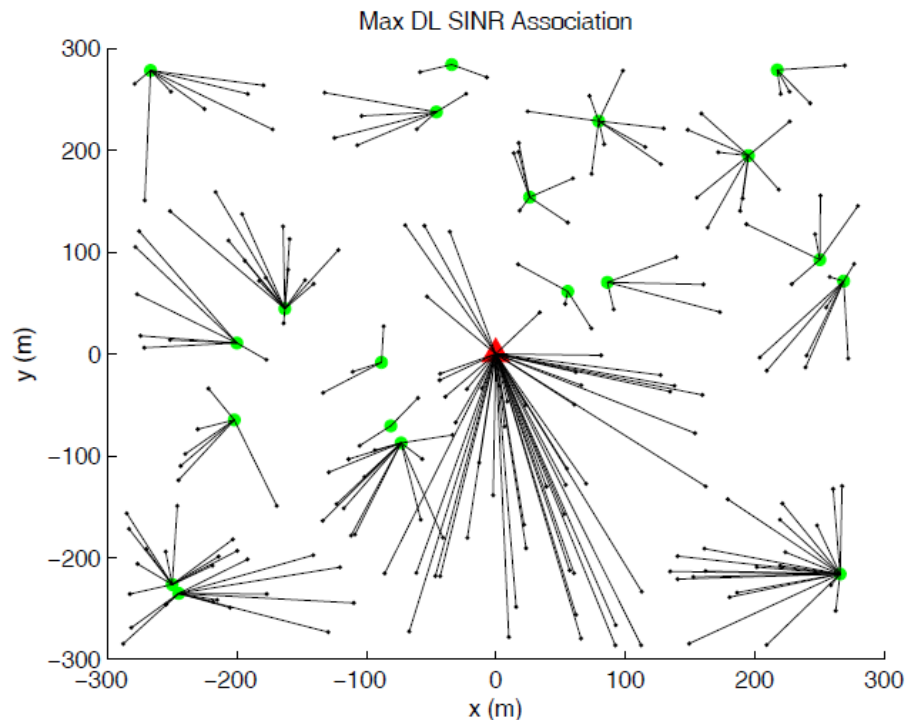
Update the association pattern of each MS 1-by-1 with the rule in Proposition 2.

3. Finalizing:

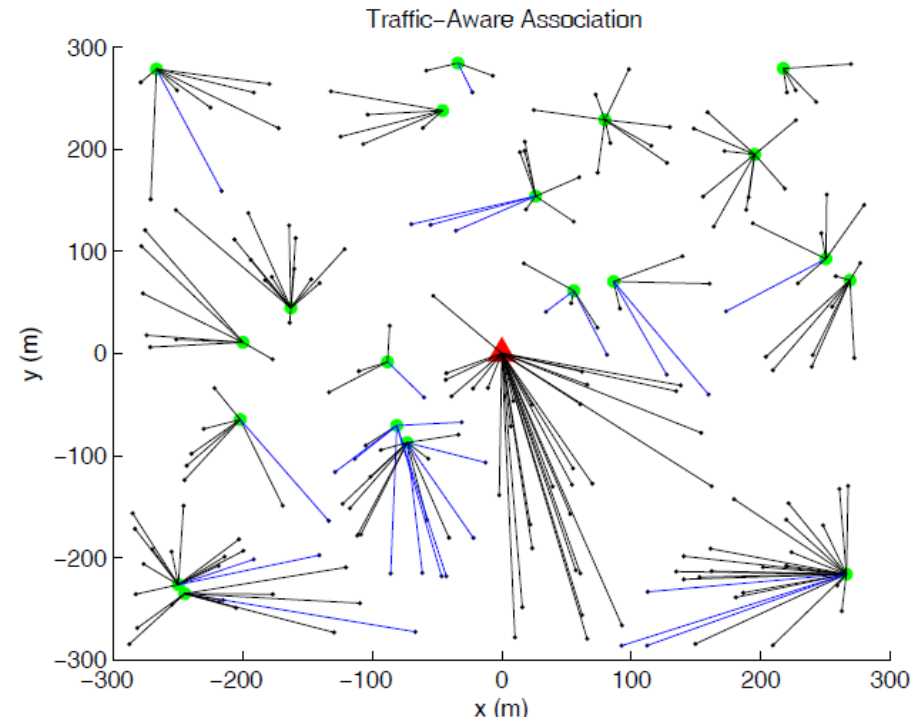
Make the converging association pattern a practical one:

$$x_{k,n}^{(f)} = 1 \{ x_{k,n} > x_{k,l}, \forall l \neq n \}$$

Simulations

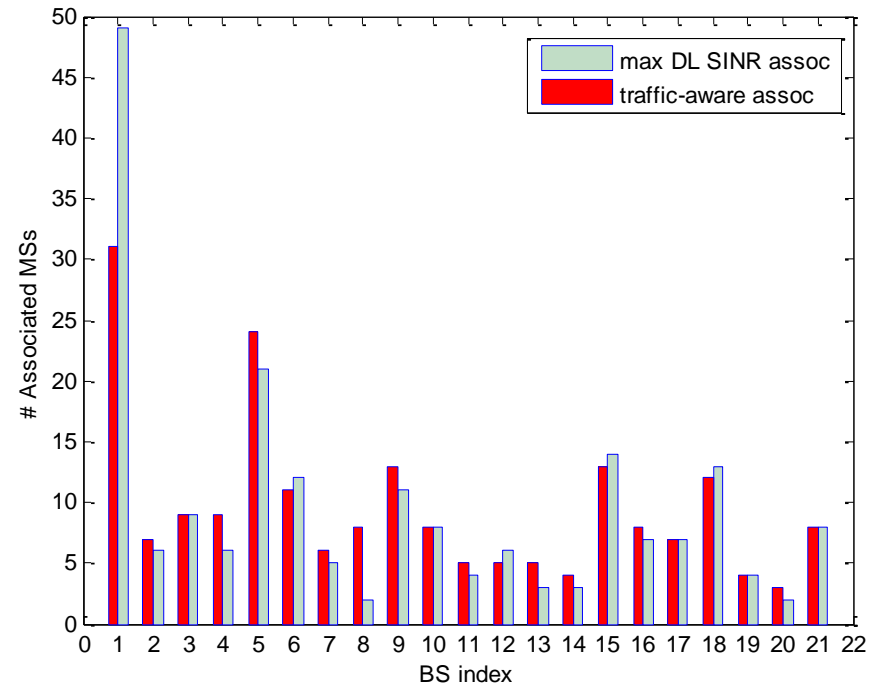
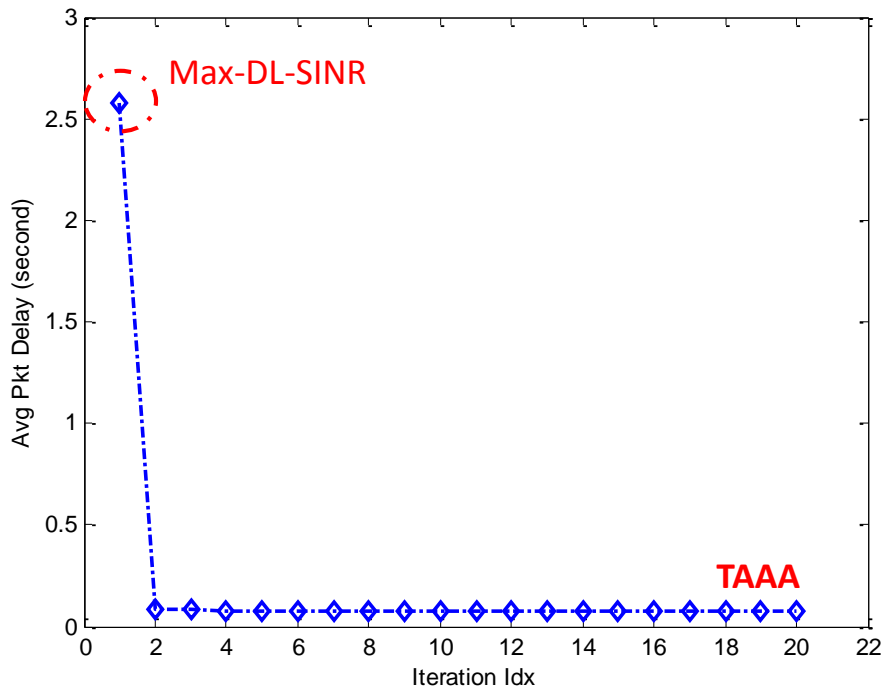


Maximum DL SINR association (Red Triangle: Macro-BS; Green Circle: Pico-BS; Black Dot: MS).



Traffic-aware association (Blue links indicate the association changes w.r.t. the Max-DL-SINR rule).

Simulations



Average packet delay and load of each BS

Conclusion



- **In heterogeneous networks with QoS traffic:**
 - ✓ Closed-form optimal resource allocation when users' association pattern is fixed
 - ✓ Closed-form optimal association scheme for one MS when given others' association
 - ✓ Our proposed low-complexity association algorithm: TAAA enjoys fast convergence and provides significant performance gain