

Structured Support Exploration For Multilayer Sparse Matrix Factorization

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Multilayer sparse matrix factorization problem

- Given a matrix A , we would like to find J matrices S_1, S_2, \dots, S_J satisfying

$$A \approx S_1 S_2 \dots S_J$$

such that S_1, S_2, \dots, S_J are sparse matrices.

Motivation: Faster linear operator

- Given a matrix $A \approx S_1 S_2 \dots S_J$ and a vector x , then

$$Ax \approx S_1(S_2(\dots(S_J x))) \quad (1)$$

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- If S_1, S_2, \dots, S_J are **sparse** matrices, evaluating the RHS of equation (1) must be faster than directly evaluating the LHS.
- Application: the Fast Fourier Transform, the Fast Hadamard Transform, etc.



The factorization of the Discrete Fourier Transform.

- 1 Introduction and motivation
- 2 Main approach and tools**
- 3 Contribution 1: Support Exploration For Sparse Matrix Factorization
- 4 Contribution 2: Promotion of k -regular sparse matrix
- 5 Conclusion

Sparse Factorization Problem Formulation

Given a matrix A , $J \in \mathbb{N}$ and $\delta_{\mathcal{E}_j}$ as indicator function of set \mathcal{E}_j , solve:

$$\text{Minimize}_{S_1, \dots, S_J} \underbrace{\|A - \prod_{j=1}^J S_j\|_F^2}_{\text{Data fidelity}} + \underbrace{\sum_{j=1}^J \delta_{\mathcal{E}_j}(S_j)}_{\text{Sparsity-inducing penalty}} \quad (2)$$

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Choice of matrices set \mathcal{E} :

- $\mathcal{E}_{row}^k = \{S \mid \|S_{i,\bullet}\|_0 \leq k\}$: at most k nonzero entries per row.
- $\mathcal{E}_{col}^k = \{S \mid \|S_{\bullet,i}\|_0 \leq k\}$: at most k nonzero entries per column.
- $\mathcal{E}_{tot}^k = \{S \mid \|S\|_0 \leq k\}$: at most k nonzero entries in total.

Problem Approach - State of The Art

- **MAIN TOOLS**: proximal operator and proximal algorithm¹.

¹Neal Parikh and Stephen Boyd. "Proximal Algorithms". In: *Found. Trends Optim.* 1.3 (Jan. 2014), 127–239. ISSN: 2167-3888. DOI: [10.1561/24000000003](https://doi.org/10.1561/24000000003).

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- **PALM IN A NUTSHELL**:

1: **for** $j \in \{1, \dots, J\}$ **do**

$$2: \quad S_j^{i+1} \leftarrow P_{\mathcal{E}_j} \left(\underbrace{S_j^i - \frac{1}{c_j^i} \nabla_{S_j} \|A - \lambda \left(\prod_{l=1}^{j-1} S_l^{i+1} \right) \left(\prod_{l=j}^J S_l^i \right)\|_F^2}_{\text{gradient step}} \right)$$

3: **end for**

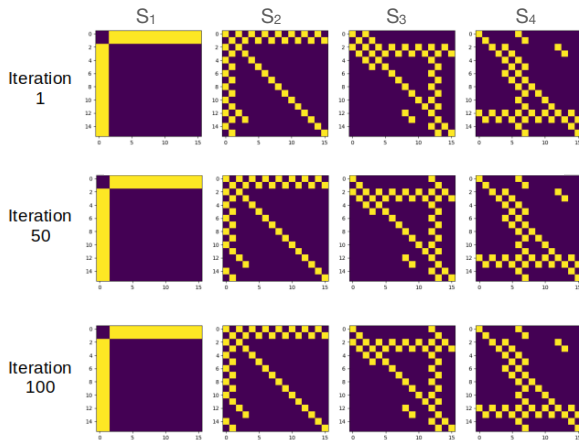
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- **PROBLEM 1:** PALM is **lazy** at exploring support.
CONTRIBUTION 1: Enforce the search of support with a new algorithm.
- **PROBLEM 2:** Proximal operators of existing sparse matrix sets often produce **rank deficient** factors.
CONTRIBUTION 2: Propose a new family of sparse matrices to avoid rank deficiency.

Support Exploration Laziness and Rank Deficiency



The evolution of support of factors of PALM during the Hadamard Transform factorization. Yellow indicates the support.

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Bilinear Hard Thresholding Pursuit (BHTP)

- Consider 2-factor factorization: $A = S_1 S_2 = XY, X \in \mathcal{E}_X, Y \in \mathcal{E}_Y$.

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1: $Y_0 = 0$.

2: **for** $n \in \{1, \dots, N\}$ **do**

3: $T_n = \text{supp} \left(\underbrace{P_{\mathcal{E}_Y} (Y_{n-1} - \lambda \nabla_Y \|A - XY_{n-1}\|^2)}_{\text{gradient step}} \right)$.

4: $Y_n = \arg \min_Y \|A - XY\|$, $\text{supp}(Y) \subseteq T_n$ (**orthogonal projection**).⁵

5: **end for**

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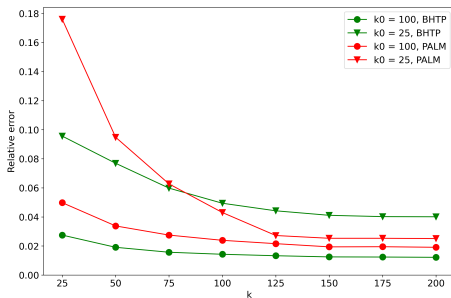
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Experimental Results

- Factorize a leadfield matrix MEG, (in functional brain imaging), $A \in \mathbb{R}^{8193 \times 204}$.
- Constraints: $\mathcal{E}_X = \mathcal{E}_{col}^{k_0}$, $\mathcal{E}_Y = \mathcal{E}_{row}^k$.
- Measure: Relative error = $\|A - XY\|_F / \|A\|_F$.



The error of BHTP and PALM factorizing the matrix MEG.

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Motivation

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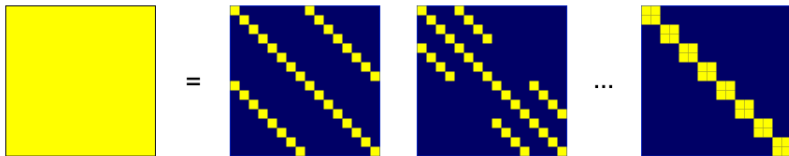
Definition

A k -regular sparse matrix $U \in \mathbb{C}^{n \times n}$ is a matrix whose columns and rows contain at most k non-zero entries each.

Let \mathcal{R}_k be the set of all k -regular matrices.

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Problem: How to project a matrix onto \mathcal{R}_k .

Generalized Hungarian Method (GHM)

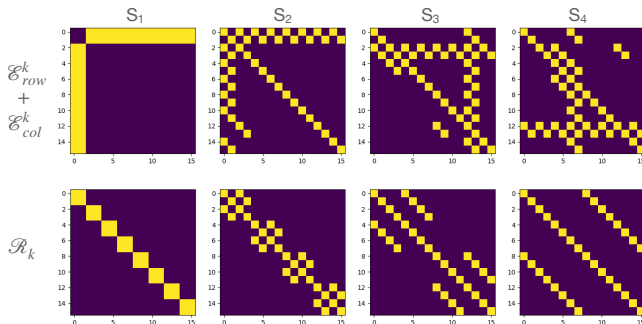
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The result of PALM with different proximal operators.

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Highlight of the paper:

- Analysis of the laziness of PALM in exploring the support. This holds with many other proximal algorithms as well
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What is left for future work:

- Other operators (such as the Discrete Fourier Transform (DFT)) resist the current algorithms.
→ Possible causes: Complex value optimization, the DFT admits fewer exact factorizations.
- Extension for nonlinear setting (Neural Networks, etc).

Practical implementation - Pyfaust

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Primal for Lorenz component L_k

$$\begin{aligned}
 \min \quad & kr_k - \sum_{i=1}^n d_{ki} \\
 \text{s.t.} \quad & r_k - y_i \geq d_{ki} \quad \forall i \in [n], \forall k \in [n] \\
 & d_{ki} \geq Mz_{ki} \quad \forall i \in [n], \forall k \in [n] \\
 & \sum_{i=1}^n z_{ki} \leq k - 1 \quad \forall k \in [n] \\
 & d_{ki} \leq 0 \quad \forall i \in [n], \forall k \in [n].
 \end{aligned} \tag{3}$$