



中国科学技术大学

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# Compressing the Tree of Canonical Huffman Coding

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The introduction of two conventional methods for compressing the tree of canonical Huffman coding

## two conventional methods

encode the number of leaf nodes in each layer

The first method need to encode the string of number  $(l_{max}, n_1, \dots, n_{l_{max}})$ .

Use an integer with  $\lceil \log_2(N-1) \rceil$  bits to encode  $l_{max}$ .

Then use  $L = \lceil \log_2(N+1) \rceil$  bits to encode each  $n_i$  for  $1 \leq i \leq l_{max}$ .

Thus, the total bits  $L_1$ :  $L_1 = \lceil \log_2(N-1) \rceil + l_{max} \lceil \log_2(N+1) \rceil$ .

Encodes  $n_i$  from the  $l_{min}$ -th layer instead of the first layer.

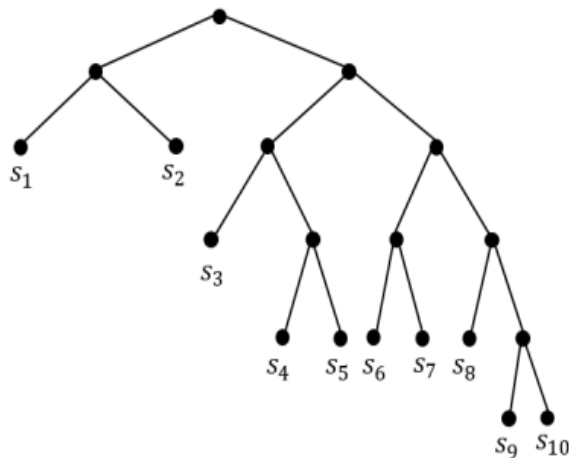
Then use  $M = \lceil \log_2(n_{max}+1) \rceil$  bits to encode each  $n_i$  for  $l_{min} \leq i \leq l_{max}$ .

The second method need to encode the string of number  $(l_{min}, l_{max}, M, n_{l_{min}}, \dots, n_{l_{max}})$ .

Thus, the total bits  $L_2$ :

$L_2 = 2\lceil \log_2(N-1) \rceil + \lceil \log_2(\lceil \log_2(N+1) \rceil - 1) \rceil + (l_{max} - l_{min})\lceil \log_2(N+1) \rceil$ .

- an example



Layer	Number	Symbols
1	0	
2	2	$s_1, s_2$
3	1	$s_3$
4	5	$s_4, s_5, s_6, s_7, s_8$
5	2	$s_9, s_{10}$

- The first method need total bits  $L_1$ :  
$$L_1 = \lceil \log_2(N-1) \rceil + l_{max} \lceil \log_2(N+1) \rceil = 24.$$

- The second method need total bits  $L_2$ :  
$$L_2 = 2 \lceil \log_2(N-1) \rceil + \lceil \log_2(\lceil \log_2(N+1) \rceil - 1) \rceil + (l_{max} - l_{min}) \lceil \log_2(N+1) \rceil = 22.$$

**Motivation:** how to compress the canonical Huffman coding tree **more effective?**

The new method for compressing the tree of canonical Huffman coding

**idea:** store the number of internal nodes  $\{b_i\}$  instead of the leaves  $\{n_i\}$

**encode method:** the mixed radix conversion

**the algorithm :**

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**Algorithm 1** ENCODING( $S_1$ ) : Encode a series of numbers  $\{b_i\}$  into an integer

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**Require:** A series of numbers  $\{b_i\}_{i=0}^{l_{max}-1}$

**Ensure:**  $S_1$

1:  $b_0 \leftarrow 1$

2:  $S_{l_{max}} \leftarrow -2$

3: **for**  $i = l_{max} - 1$  **to** 1 **do**

4:    $S_i = S_{i+1} \times 2b_{i-1} + (b_i - 1)$

5: **end for**

6: **return**  $S_1$

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# Proposed method



The proposed method need total bits  $L_3$ :

$$L_3 = \lceil \log_2 |S_1| \rceil + \lceil \log_2 (N-1) \rceil.$$

to encode the number of bits for encoding  $S_1$

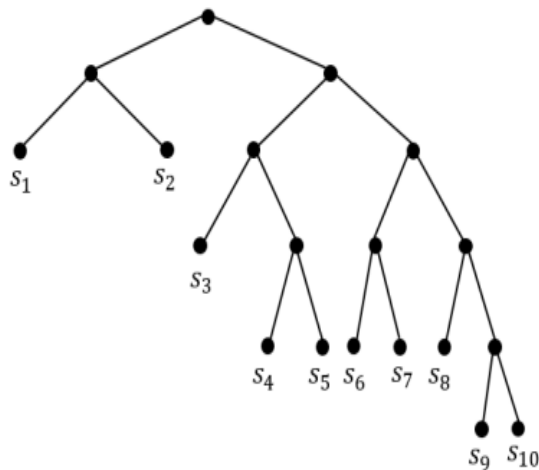


Table 1: The encoding of the string of numbers (2, 2, 3, 1).

$i$	1	2	3	4	5
Number of internal nodes $b_i$	2	2	3	1	
Upper bound $2b_{i-1}$	2	4	4	6	
$b_i - 1$	1	1	2	0	
$S_i$	-365	-183	-46	-12	-2

$$L_3 = \lceil \log_2 |-365| \rceil + \lceil \log_2 (10-1) \rceil = 13$$

# Experimental results



The comparison of the compression ratio between the proposed method and two conventional methods.

## Experimental data: Calgary corpus

We use a conventional way to construct the canonical Huffman tree for various data sets.

Table 2: Number of bits.

text	[4]	[13]	ours
bib	119	87	52
book1	126	77	55
book2	119	87	58
geo	96	79	52
news	105	77	52
obj1	54	55	28
obj2	84	62	47
paper1	112	82	56
paper2	119	73	57
pic	144	104	70
progc	105	77	49
progl	105	77	52
progp	112	73	54
trans	119	87	57

Table 3: Compression ratio.

text	$R_1$	$R_2$
bib	56.30%	40.23%
book1	56.35%	28.57%
book2	51.26%	33.33%
geo	45.83%	34.18%
news	50.48%	32.47%
obj1	48.15%	49.09%
obj2	44.05%	24.19%
paper1	50%	31.71%
paper2	52.10%	21.92%
pic	51.39%	32.69%
progc	53.33%	36.36%
progl	50.48%	32.47%
progp	51.79%	26.03%
trans	52.10%	34.48%
average	50.97%	32.69%

# This work

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graph LR; A((This work)) --- B[motivation]; A --- C[method]; A --- D[results]; B --- E[how to compress the canonical Huffman coding tree more effective?]; C --- F[Use the mixed radix conversion to encode the number of the internal nodes.]; D --- G[Simulation shows that the proposed method can reduce 51% and 33% space with the conventional methods on average.]; G --> B;
```

motivation

how to compress the canonical Huffman coding tree more effective?

method

Use the mixed radix conversion to encode the number of the internal nodes.

results

Simulation shows that the proposed method can reduce 51% and 33% space with the conventional methods on average.



# Data Compression Conference



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# Thank you !

For more details, please refer to our paper:  
**Compressing the Tree of Canonical Huffman Coding**

That is all my presentation

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