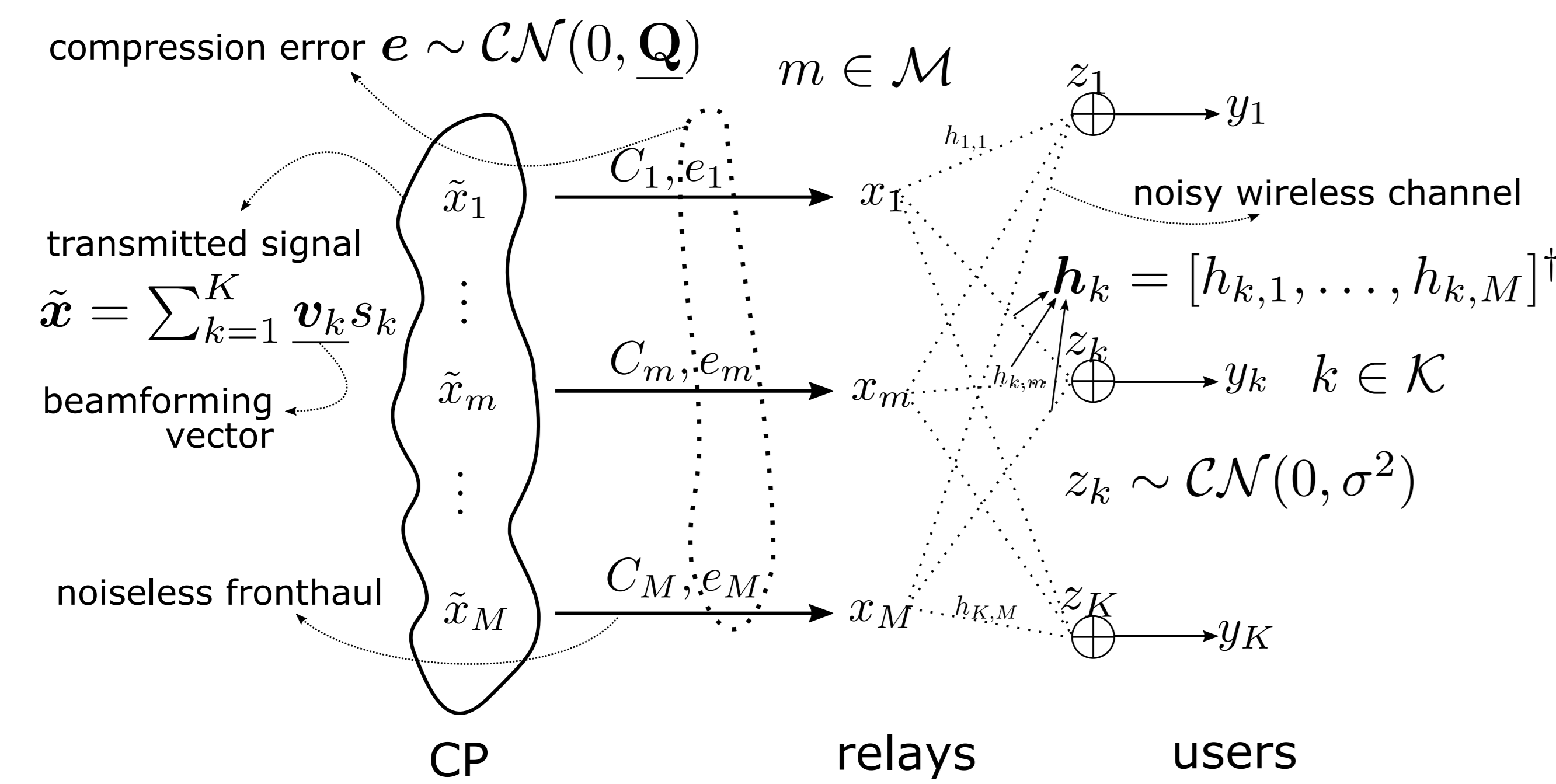


Main Contribution

- The **downlink joint beamforming and compression problem** is a challenging task in the **cooperative cellular network**.
- The **rate-limited fronthaul** between the central processor (CP) and base stations (BSs) poses constraints that are difficult to handle.
- The **joint optimization** of beamforming vectors and the covariance matrix of the compression error also makes the problem highly non-convex.
- We establish the **tightness** of the semidefinite relaxation (SDR) of the considered problem and thus the equivalence of the two problems.
- We propose a **global** and **efficient** algorithm for the considered problem via Lagrangian duality.

System Model



Problem Formulation

- The equivalent joint beamforming and compression problem [1, Proposition 4]:

$$\begin{aligned} \min_{\{\mathbf{v}_k\}, \mathbf{Q}} \quad & \sum_{k=1}^K \|\mathbf{v}_k\|^2 + \mathbf{Q} \bullet \mathbf{I} \\ \text{s.t.} \quad & \mathbf{v}_k^\dagger \mathbf{H}_k \mathbf{v}_k - \gamma_k \left(\sum_{j \neq k} \mathbf{v}_j^\dagger \mathbf{H}_k \mathbf{v}_j + \mathbf{Q} \bullet \mathbf{H}_k + \sigma^2 \right) \geq 0, \quad \forall k \in \mathcal{K}, \\ & \eta_m \tilde{\mathbf{Q}}_m - \mathbf{E}_m^\dagger \left(\sum_{k=1}^K \mathbf{v}_k \mathbf{v}_k^\dagger + \mathbf{Q} \right) \mathbf{E}_m \succeq \mathbf{0}, \quad \forall m \in \mathcal{M}, \\ & \mathbf{Q} \succeq \mathbf{0}, \end{aligned} \quad (\text{P})$$

where $\mathbf{H}_k = \mathbf{h}_k \mathbf{h}_k^\dagger$, γ_k is the rate target for user k , $\eta_m = 2^{C_m}$ and the right bottom of $\tilde{\mathbf{Q}}_m$ is filled with $\mathbf{Q}^{(m:M, m:M)}$ and the other parts are zero.

Problem Formulation (Cont.)

- The semidefinite relaxation problem (SDR) and its dual problem can be derived accordingly and have the following forms:

$$\begin{aligned} \min_{\mathbf{V}_k, \mathbf{Q}} \quad & P^{\text{dl}} \\ \text{s.t.} \quad & \beta_k : a_k \geq 0, \quad k \in \mathcal{K}, \\ & \mathbf{A}_m : \mathbf{B}_m \succeq \mathbf{0}, \quad m \in \mathcal{M}, \\ & \mathbf{V}_k, \mathbf{Q} \succeq \mathbf{0}, \quad k \in \mathcal{K}, \end{aligned} \quad \begin{aligned} \max_{\beta_k, \mathbf{A}_m} \quad & P^{\text{ul}} \\ \text{s.t.} \quad & \mathbf{V}_k : \mathbf{C}_k - \beta_k \mathbf{H}_k \succeq \mathbf{0}, \quad \forall k \in \mathcal{K}, \\ & \mathbf{Q} : \mathbf{D} \succeq \mathbf{0}, \\ & \beta_k \geq 0, \quad \forall k \in \mathcal{K}, \\ & \mathbf{A}_m \succeq \mathbf{0}, \quad \forall m \in \mathcal{M}, \end{aligned}$$

where a_k , \mathbf{B}_m and P^{dl} are functions of $\{\mathbf{V}_k\}$ and \mathbf{Q} ; P^{ul} is a function of $\{\beta_k\}$, $\{\mathbf{A}_m\}$ and

$$\mathbf{C}_k(\{\beta_k\}, \{\mathbf{A}_m\}) = \mathbf{I} + \sum_{m=1}^M \mathbf{E}_m^\dagger \mathbf{A}_m \mathbf{E}_m + \sum_{j \neq k} \beta_j \gamma_j \mathbf{H}_j,$$

$$\mathbf{D}(\{\beta_k\}, \{\mathbf{A}_m\}) = \mathbf{I} - \sum_{m=1}^M \eta_m \tilde{\mathbf{Q}}_m + \sum_{k=1}^K \beta_k \gamma_k \mathbf{H}_k + \sum_{m=1}^M \mathbf{E}_m^\dagger \mathbf{A}_m \mathbf{E}_m.$$

The Proposed Algorithm

- Design an efficient algorithm for solving (P).
- We show that (SDR) is a **tight** relaxation of (P).
- Solve the KKT optimality conditions of (SDR) based on its special structure:

- Write out the equivalent **KKT conditions**:

$$\begin{aligned} \left\{ \begin{aligned} \mathbf{D}(\{\beta_k\}, \{\mathbf{A}_m\}) &= \mathbf{0}, & (1) \\ \text{rank}(\mathbf{A}_m) &= 1, \quad \mathbf{A}_m \succeq \mathbf{0}, \quad \forall m \in \mathcal{M}, \\ \mathbf{A}_m^{(1:m-1, 1:m-1)} &= \mathbf{0}, \quad \mathbf{A}_m^{(m:M, 1:m-1)} = \mathbf{0}, \quad \forall m \in \mathcal{M}, \end{aligned} \right\} & (2) \\ \left\{ \begin{aligned} \text{rank}(\mathbf{C}_k(\{\beta_k\}, \{\mathbf{A}_m\}) - \beta_k \mathbf{H}_k) &= M - 1, \quad \forall m \in \mathcal{M}, \\ \mathbf{C}_k(\{\beta_k\}, \{\mathbf{A}_m\}) - \beta_k \mathbf{H}_k &\succeq \mathbf{0}, \quad \forall m \in \mathcal{M}, \end{aligned} \right\} & (3) \\ \left\{ \begin{aligned} \beta_k &\geq 0, \quad \forall k \in \mathcal{K}, & (4) \\ \mathbf{V}_k \bullet (\mathbf{C}_k(\{\beta_k\}, \{\mathbf{A}_m\}) - \beta_k \mathbf{H}_k) &= 0, \quad \forall k \in \mathcal{K}, & (5) \\ \mathbf{V}_k &\succeq \mathbf{0}, \quad \text{rank}(\mathbf{V}_k) = 1, \quad \forall k \in \mathcal{K}, & (6) \\ a_k(\{\mathbf{V}_k\}, \mathbf{Q}) &= 0, \quad \forall k \in \mathcal{K}, & (7) \\ \mathbf{B}_m(\{\mathbf{V}_k\}, \mathbf{Q}) &\succeq \mathbf{0}, \quad \forall m \in \mathcal{M}, & (8) \\ \mathbf{A}_m \bullet \mathbf{B}_m(\{\mathbf{V}_k\}, \mathbf{Q}) &= 0, \quad \forall m \in \mathcal{M}, & (9) \\ \mathbf{Q} &\succeq \mathbf{0}. & (10) \end{aligned} \right. \end{aligned}$$

- Separate the equations into **two sets** (Eqs. (1)–(4) and Eqs. (5)–(10)) and solve the equations involving the dual variables first and then the equations involving the primal variables.
- Show that each set of equations can be solved elegantly via **fixed-point iteration**.

The Proposed Algorithm (Cont.)

- First, consider solving Eqs. (1)–(4).
- Given** $\{\beta_k\}$, one equivalent form of Eq. (1) is

$$\sum_{m=1}^M \eta_m \tilde{\mathbf{A}}_m - \sum_{m=1}^M \mathbf{E}_m^\dagger \mathbf{A}_m \mathbf{E}_m = \mathbf{I} + \sum_{k=1}^K \beta_k \gamma_k \mathbf{H}_k \triangleq \mathbf{\Gamma}.$$

- Only \mathbf{A}_1 affects the first row and column of matrix $\mathbf{\Gamma}$. \Rightarrow The entries in the first row of \mathbf{A}_1 should be $\left[\frac{1}{\eta_1-1} \mathbf{\Gamma}^{(1,1)}, \frac{1}{\eta_1} \mathbf{\Gamma}^{(1,2:M)} \right]$.
- \mathbf{A}_1 is of rank one. (Eq. (2)) \Rightarrow Further obtain all entries of \mathbf{A}_1 .
- Subtract all terms related to \mathbf{A}_1 from both sides:

$$\sum_{m=2}^M \eta_m \tilde{\mathbf{A}}_m - \sum_{m=2}^M \mathbf{E}_m^\dagger \mathbf{A}_m \mathbf{E}_m = \mathbf{\Gamma} - \eta_1 \mathbf{A}_1 + \mathbf{E}_1 \mathbf{A}_1 \mathbf{E}_1 \triangleq \mathbf{\Gamma}'.$$

- Repeat the above procedure to find all \mathbf{A}_m . Denote the solution to Eqs. (1)–(2) as $\{\mathbf{A}_m(\{\beta_k\})\}$.
- Given** $\{\mathbf{A}_m\}$, by Eq. (3) and since $\mathbf{H}_k \succeq \mathbf{0}$ is of rank one, we have

$$\beta_k(\{\mathbf{A}_m\}, \{\beta_j\}_{j \neq k}) = (\mathbf{h}_k^\dagger \mathbf{C}_k^{-1} \mathbf{h}_k)^{-1} > 0.$$

- If one find $(\{\beta_k\}, \{\mathbf{A}_m(\{\beta_k\})\})$ that satisfy

$$\beta_k = I_k(\{\beta_k\}) \triangleq \beta_k(\{\mathbf{A}_m(\{\beta_k\})\}, \{\beta_j\}_{j \neq k}), \quad \forall k \in \mathcal{K}, \quad (11)$$

then all Eqs. (1)–(4) holds.

- Define $\boldsymbol{\beta} = [\beta_1, \dots, \beta_K]^\top$ and $I(\boldsymbol{\beta}) = [I_1(\{\beta_k\}), \dots, I_K(\{\beta_k\})]^\top$, then (11) is to find the fixed-point of function $I(\cdot)$.
- We can show that $I(\cdot)$ is a standard interference function.
- The fixed-point iteration $\boldsymbol{\beta}^{(i+1)} = I(\boldsymbol{\beta}^{(i)})$ will converge to the unique fixed point of $I(\cdot)$. (Lemma and [3, Theorem 2])
- Solving Eqs. (5)–(10) can also be reduced to find a fixed-point of some standard interference function $J(\cdot)$. Hence, it can also be solved by performing a fixed-point iteration.

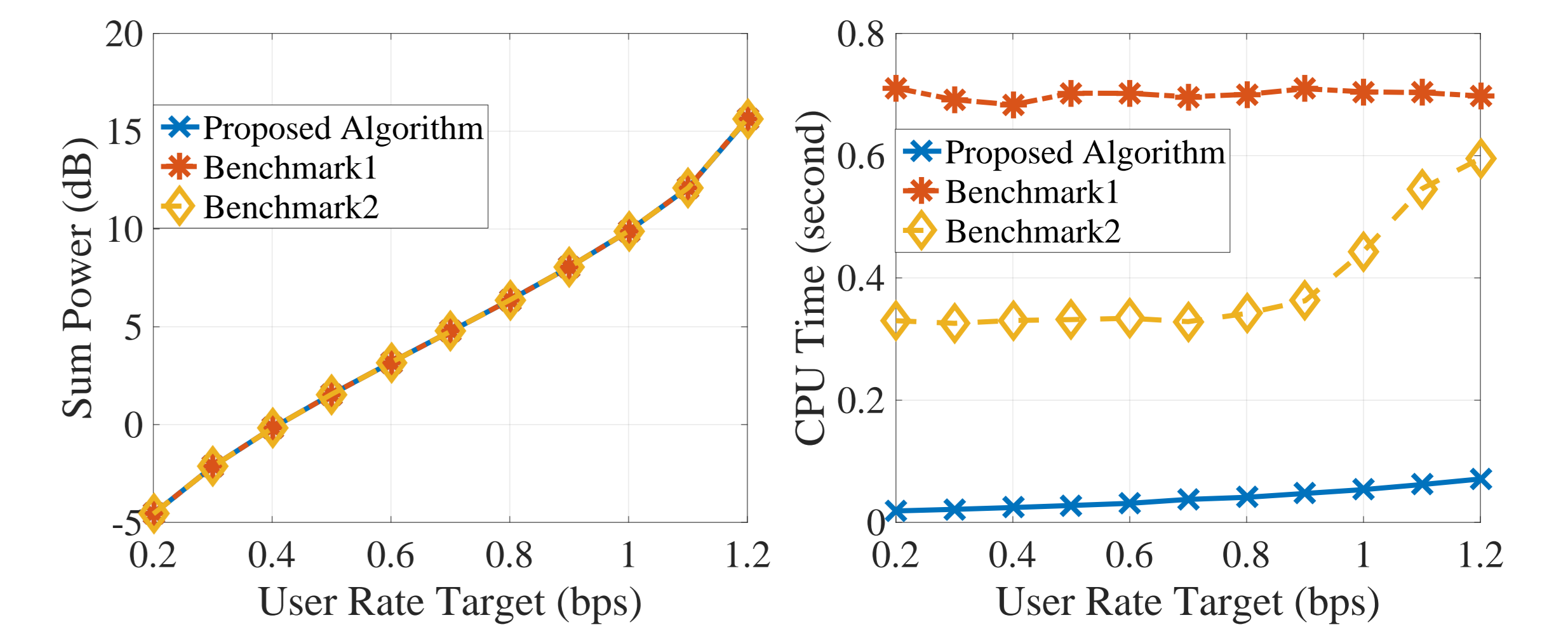
Algorithm 1 Proposed Algorithm for Solving Problem (P)

- Find $\{\beta_k\}$ and $\{\mathbf{A}_m\}$ that satisfy Eqs. (1)–(4) by performing a fixed-point iteration via $I(\cdot)$ on $\{\beta_k\}$ until the desired error bound is met.
- Find $\{\mathbf{V}_k\}$ and \mathbf{Q} that satisfy Eqs. (5)–(10) by performing a fixed-point iteration via $J(\cdot)$ on $\{\mathbf{V}_k\}$ until the desired error bound is met.
- Find \mathbf{v}_k such that $\mathbf{V}_k = \mathbf{v}_k \mathbf{v}_k^\dagger$, $\forall k \in \mathcal{K}$.
- Output**: $\{\mathbf{v}_k\}$ and \mathbf{Q} .

- Global optimality guarantee**: If (SDR) is feasible, then proposed Algorithm 1 returns the optimal solution of problem (P).
- Computationally efficient**: The evaluation of fixed-point iteration functions, namely $I(\cdot)$ and $J(\cdot)$, are cheap.

Simulation Results

- Consider a network with
 - $M = 8$ relays and $K = 10$ users,
 - the wireless channels between relays and users are generated based on the i.i.d. **Rayleigh fading** model following $\mathcal{CN}(0, 1)$,
 - and the **fronthaul capacities** between all relays and the CP are set to be **3 bits per symbol** (bps).
- Moreover, the **noise powers** at the users are set to be $\sigma^2 = 1$.
- The **rate targets** for all the users are assumed to be **identical**.
- All results are obtained by **averaging over 200 Monte-Carlo runs**.



Left: Average sum power versus the user rate target; Right: Average CPU time versus the user rate target.

- Benchmark1: directly call CVX to solve (SDR) \Rightarrow **verify the tightness**
- Benchmark2: the proposed algorithm in [1] \Rightarrow **compare the efficiency**
- The left figure verifies the **tightness** of the SDR and the global optimality of the solution returned by the proposed algorithm.
- The right figure shows the **high efficiency** of our proposed algorithm.

Conclusion

- Propose an **efficient** and **global** algorithm for solving the downlink beamforming and compression problem.
- Solve the KKT conditions by judiciously exploiting the **problem structure**.
- Achieve the **global optimality** as the state-of-the-art algorithm proposed in [1] but with a **significant less CPU time**.

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