

Distributed Online Learning with Adversarial Participants in An Adversarial Environment

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Background

Online Learning

- Online learning is a powerful tool to process **streaming data**.
- In response to an environment that provides **(adversarial) losses** sequentially, an online learning algorithm makes one-step-ahead decisions.

Distributed Online Learning

- **Multiple participants** separately collect streaming data, make local decisions.
- **Server aggregates** all local decisions to a global one.
- Applications: online web ranking and online advertisement recommendation.

Performance of an online learning algorithm is characterized by (adversarial) regret, and a **sublinear (adversarial) regret** is preferred.

Background

But not all participants are honest.

Byzantine Attack

Adversarial participants (called Byzantine participants) can collude and arbitrarily modify the messages sent to the server.

Is it possible to develop a Byzantine-robust distributed online learning algorithm with provable **sublinear adversarial regret**, in an **adversarial environment** and in the presence of **adversarial participants** ?

Answer Is Negative

- ✗ Distributed online gradient descent with mean: infinite adversarial regret.
- ✗ Even with robust aggregation rules: linear adversarial regret.

Problem Statement: Adversarial Regret

- Consider n participants in \mathcal{N} , h honest in \mathcal{H} , b Byzantine in \mathcal{B} , $n = h + b$.
- Suppose the ratio of Byzantine participants is less than half: $\alpha := \frac{b}{n} < \frac{1}{2}$.
- **Goal:** minimize **adversarial regret** over T steps

$$R_T := \sum_{t=1}^T f_t(w_t) - \min_{w \in \mathbb{R}^d} \sum_{t=1}^T f_t(w), \quad (1)$$

where

$$f_t(w) := \frac{1}{h} \sum_{j \in \mathcal{H}} f_t^j(w), \quad (2)$$

and f_t^j is the loss revealed to $j \in \mathcal{H}$ at the end of step t .

Byzantine-robust Distributed Online Gradient Descent

Each honest participant j makes its local decision by **online gradient descent**:

$$w_{t+1}^j = w_t - \eta_t \nabla f_t^j(w_t), \quad \text{step size } \eta_t > 0. \quad (3)$$

- **Baseline**: distributed **online gradient descent** (3) with mean aggregation
Server aggregates messages z_{t+1}^j (w_{t+1}^j from honest and arbitrary from Byzantine)

$$w_{t+1} = \frac{1}{n} \sum_{j=1}^n z_{t+1}^j. \quad (4)$$

- **Ours**: Byzantine-robust distributed **online gradient descent** (3) with **AGG**

$$w_{t+1} = \text{AGG}(z_{t+1}^1, z_{t+1}^2, \dots, z_{t+1}^n). \quad (5)$$

AGG is Robust Bounded Aggregation, if

$$\|w_{t+1} - \bar{z}_{t+1}\|^2 = \|\text{AGG}(z_{t+1}^1, z_{t+1}^2, \dots, z_{t+1}^n) - \bar{z}_{t+1}\|^2 \leq C_\alpha^2 \zeta^2, \quad \bar{z}_{t+1} := \frac{1}{h} \sum_{j \in \mathcal{H}} z_{t+1}^j,$$

where $\|\bar{z}_{t+1} - z_{t+1}^j\|^2 \leq \zeta^2$, C_α is a constant dependent on α and aggregation.

Assumptions & Theorem 1

Define $\nabla \bar{f}_t(w_t) := \frac{1}{h} \sum_{j \in \mathcal{H}} \nabla f_t^j(w_t)$ and $w^* := \arg \min_{w \in \mathbb{R}^d} \sum_{t=1}^T f_t(w)$. For any honest participant's loss f_t^j where $j \in \mathcal{H}$ and any $x, y \in \mathbb{R}^d$ we assume

Assumption 1 L -smoothness. $\|\nabla f_t^j(x) - \nabla f_t^j(y)\| \leq L\|x - y\|$.

Assumption 2 μ -strong convexity. $\langle \nabla f_t^j(x), x - y \rangle \geq f_t^j(x) - f_t^j(y) + \frac{\mu}{2}\|x - y\|^2$.

Assumption 3 Bounded deviation. $\|\nabla f_t^j(w_t) - \nabla \bar{f}_t(w_t)\|^2 \leq \sigma^2$.

Assumption 4 Bounded gradient at the overall best solution.

$$\left\| \frac{1}{h} \sum_{j \in \mathcal{H}} \nabla f_t^j(w^*) \right\|^2 \leq \xi^2.$$

Theorem 1

Under Assumptions 1, 2, 3 and 4, if $\eta = \mathcal{O}(\frac{1}{\sqrt{T}})$, Byzantine-robust distributed online gradient descent has a **linear** adversarial regret bound

$$R_T = \mathcal{O}((C_\alpha^2 \sigma^2 + \xi^2) \sqrt{T}) + \mathcal{O}(C_\alpha^2 \sigma^2 T). \quad (6)$$

We construct a counter-example to demonstrate $\mathcal{O}(\sigma^2 T)$ is tight.

How to derive sublinear regret under Byzantine Attacks?
→ Not fully adversarial environment.

Byzantine-Robust Distributed Online Momentum

- **Not fully adversarial environment:** losses are independent and identically distributed (i.i.d.), meaning $f_t^j \sim \mathcal{D}$ for all $j \in \mathcal{H}$ and all t .
- Define the expected loss $F(w) := \mathbb{E}_{\mathcal{D}} f_t^j(w)$ for all $j \in \mathcal{H}$ and all t .
- **New Goal:** minimize **stochastic regret** over T steps

$$S_T := \mathbb{E} \sum_{t=1}^T F(w_t) - T \cdot \min_{w \in \mathbb{R}^d} F(w). \quad (7)$$

- Each honest participant j maintains a **momentum** vector to reduce variance

$$m_t^j = \nu_t \nabla f_t^j(w_t) + (1 - \nu_t) m_{t-1}^j, \quad (8)$$

where $0 < \nu_t < 1$ is momentum parameter. Then, it makes a local decision

$$w_{t+1}^j = w_t - \eta_t m_t^j. \quad (9)$$

- **Ours:** Byzantine-Robust distributed **online momentum** (9) with **AGG** (5).

Assumptions & Theorem 2

For expected loss $F(w)$ and any $x, y \in \mathbb{R}^d$, we assume

Assumption 5 L -smoothness. $\|\nabla F(x) - \nabla F(y)\| \leq L\|x - y\|$.

Assumption 6 μ -strong convexity. $\langle \nabla F(x), x - y \rangle \geq F(x) - F(y) + \frac{\mu}{2}\|x - y\|^2$.

Assumption 7 Bounded variance. $\mathbb{E}_{\mathcal{D}} \|\nabla f_t^j(w_t) - \nabla F(w_t)\|^2 \leq \sigma^2$.

Theorem 2

Supposed losses are i.i.d., under Assumptions 5, 6 and 7, if $\eta = \mathcal{O}(\frac{1}{\sqrt{T}})$ and $\nu = \mathcal{O}(\frac{1}{\sqrt{T}})$, Byzantine-robust distributed online momentum has a **sublinear** stochastic regret bound

$$S_T = \mathcal{O} \left(\left(1 + \frac{\sigma^2}{h} \left(1 + (h+1)C_\alpha^2 \right) \frac{L^4}{\mu^4} \right) \sqrt{T} \right). \quad (10)$$

Numerical Experiments¹

- Softmax regression on the i.i.d. MNIST dataset.
- Measurement: **adversarial regret and accuracy**.

Byzantine-robust distributed online gradient descent show robustness.

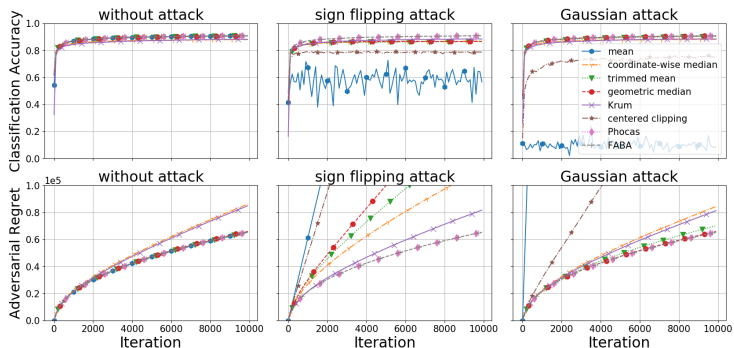


Figure 1: Performance of Byzantine-robust distributed online gradient descent.

¹More results and codes are available at <https://github.com/wanger521/OGD>.

Numerical Experiments

Momentum show improvement!

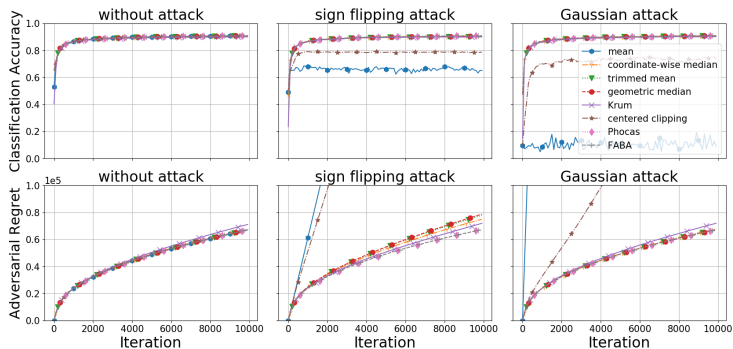


Figure 2: Performance of Byzantine-robust distributed online momentum.

More experiment results on non-i.i.d. data are shown in the paper.

- Investigate **Byzantine-robustness of distributed online learning** for first time.
- Show **tight linear adversarial regret** bound for Byzantine-robust distributed online gradient descent.
- Establish **sublinear stochastic regret** bound for Byzantine-robust distributed online momentum with i.i.d. distribution.

Thank You!