

Determined Blind Source Separation via Proximal Splitting Algorithm

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Introduction

- Many **independence-based blind source separation (BSS)** methods reduces to the following minimization problem:

$$\text{Minimize}_{\{W[f]\}_{f=1}^F} \mathcal{P}(W[f]\mathbf{x}[t, f]) - \sum_{f=1}^F \log |\det(W[f])|$$

- Laplace-distribution-based independent component analysis (**FDICA**)

$$\mathcal{P}(\mathbf{y}[t, f]) = C \|\mathbf{y}[t, f]\|_1 = C \sum_{m=1}^M \sum_{t=1}^T \sum_{f=1}^F |y_m[t, f]|$$

- Spherical-Laplace-distribution-based Independent vector analysis (**IVA**)

$$\mathcal{P}(\mathbf{y}[t, f]) = C \|\mathbf{y}[t, f]\|_{2,1} = C \sum_{m=1}^M \sum_{t=1}^T \left(\sum_{f=1}^F |y_m[t, f]|^2 \right)^{\frac{1}{2}}$$

- Independent low-rank matrix analysis (**ILRMA**) [1]

$$\mathcal{P}(\mathbf{y}[t, f]) = C \sum_{m=1}^M \min_{\varphi_{f,r}^{[m]} \geq 0, \psi_{r,t}^{[m]} \geq 0} \sum_{t=1}^T \sum_{f=1}^F \left(\frac{|y_m[t, f]|^2}{\sum_{r=1}^R \varphi_{f,r}^{[m]} \psi_{r,t}^{[m]}} + \log \sum_{r=1}^R \varphi_{f,r}^{[m]} \psi_{r,t}^{[m]} \right)$$

- **State-of-the-art algorithms** based on the majorization-minimization (MM) principle [1–3] require specially designed upper bounds whose **derivation might be complicated and time-consuming** for some \mathcal{P} .

Primal-dual Splitting Algorithm

$$\begin{cases} \tilde{\mathbf{w}} = \text{prox}_{\mu_1 g} [\mathbf{w}^{[k]} - \mu_1 \mu_2 L^H \mathbf{y}^{[k]}] \\ \mathbf{z} = \mathbf{y}^{[k]} + L(2\tilde{\mathbf{w}} - \mathbf{w}^{[k]}) \\ \tilde{\mathbf{y}} = \mathbf{z} - \text{prox}_{h/\mu_2} [\mathbf{z}] \\ (\mathbf{w}^{[k+1]}, \mathbf{y}^{[k+1]}) = \alpha(\tilde{\mathbf{w}}, \tilde{\mathbf{y}}) + (1 - \alpha)(\mathbf{w}^{[k]}, \mathbf{y}^{[k]}) \end{cases}$$

- Primal-dual splitting (PDS) algorithm is one of the **proximal splitting algorithms** which can solve the convex optimization problem.

$$\text{Minimize}_{\mathbf{w}} g(\mathbf{w}) + h(L\mathbf{w})$$

- PDS iteratively utilizes the **proximity operator** of each function.

$$\text{prox}_{\mu g}[\mathbf{y}] = \arg \min_{\mathbf{z}} \left[g(\mathbf{z}) + \frac{1}{2\mu} \|\mathbf{y} - \mathbf{z}\|_2^2 \right]$$

- Since the **proximity operators are easier subproblems** which can be solved in many ways, the PDS algorithm applied to the independence-based BSS should be **able to handle complicated source model** \mathcal{P} more easily than the state-of-the-art MM-based algorithms.

Proposed Primal-dual Algorithm for Independence-based BSS problem

- By rewriting matrix determinant with singular values and vectorizing variables, the BSS problem can be seen as a **PDS applicable form**.

$$\text{Minimize}_{\mathbf{w}} \mathcal{I}(\mathbf{w}) + \mathcal{P}(X\mathbf{w})$$

- **Matrix determinant was rewritten by singular values** because the proximity operators of orthogonally invariant functions are known. This modification allows direct handling of an over-determined case.

$$\mathcal{I}(\mathbf{w}) = -\sum_{f=1}^F \sum_{m=1}^M \log \sigma_m(\mathcal{M}(\mathbf{w})[f]) \quad (\mathcal{M}(\mathbf{w})[f] = W[f])$$

- **Variables are vectorized** so that the role of observed data X in terms of the optimization variable \mathbf{w} becomes apparent.

Algorithm 1 PDS-BSS

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1: Input:  $X, \mathbf{w}^{[1]}, \mathbf{y}_1^{[1]}, \mu_1, \mu_2, \alpha$ 
2: Output:  $\mathbf{w}^{[K+1]}$ 
3: for  $k = 1, \dots, K$  do
4:    $\tilde{\mathbf{w}} = \text{prox}_{\mu_1 \mathcal{I}} [\mathbf{w}^{[k]} - \mu_1 \mu_2 X^H \mathbf{y}^{[k]}]$ 
5:    $\mathbf{z} = \mathbf{y}^{[k]} + X(2\tilde{\mathbf{w}} - \mathbf{w}^{[k]})$ 
6:    $\tilde{\mathbf{y}} = \mathbf{z} - \text{prox}_{\frac{1}{\mu_2} \mathcal{P}} [\mathbf{z}]$ 
7:    $\mathbf{y}^{[k+1]} = \alpha \tilde{\mathbf{y}} + (1 - \alpha) \mathbf{y}^{[k]}$ 
8:    $\mathbf{w}^{[k+1]} = \alpha \tilde{\mathbf{w}} + (1 - \alpha) \mathbf{w}^{[k]}$ 
9: end for

```

Modification of singular values

Proximity operator of source term

- The proximity operator of \mathcal{I} is to apply the proximity operator of $-\log$ function to each singular value of the un-vectorized variables.

$$\text{prox}_{-\mu \log}[\sigma_m] = (\sigma_m + \sqrt{\sigma_m^2 + 4\mu})/2$$

- **Source models defined by adding multiple penalty terms** can also be handled by the proposed algorithm with slight modifications. Such multiple penalty terms can be difficult to handle for MM algorithms.

$$\text{Minimize}_{\mathbf{w}} \mathcal{I}(\mathbf{w}) + \sum_{q=1}^Q \mathcal{P}_q(X\mathbf{w})$$

Algorithm 2 PDS-BSS-multiPenalty

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1: Input:  $X, \mathbf{w}^{[1]}, \mathbf{y}_1^{[1]}, \dots, \mathbf{y}_Q^{[1]}, \mu_1, \mu_2, \alpha$ 
2: Output:  $\mathbf{w}^{[K+1]}$ 
3: for  $k = 1, \dots, K$  do
4:    $\tilde{\mathbf{w}} = \text{prox}_{\mu_1 \mathcal{I}} [\mathbf{w}^{[k]} - \mu_1 \mu_2 X^H (\sum_{q=1}^Q \mathbf{y}_q^{[k]})]$ 
5:   for  $q = 1, \dots, Q$  do
6:      $\mathbf{z}_q = \mathbf{y}_q^{[k]} + X(2\tilde{\mathbf{w}} - \mathbf{w}^{[k]})$ 
7:      $\tilde{\mathbf{y}}_q = \mathbf{z}_q - \text{prox}_{\frac{1}{\mu_2} \mathcal{P}_q} [\mathbf{z}_q]$ 
8:      $\mathbf{y}_q^{[k+1]} = \alpha \tilde{\mathbf{y}}_q + (1 - \alpha) \mathbf{y}_q^{[k]}$ 
9:   end for
10:   $\mathbf{w}^{[k+1]} = \alpha \tilde{\mathbf{w}} + (1 - \alpha) \mathbf{w}^{[k]}$ 
11: end for

```

Modification of singular values

Proximity operator of each source term

New BSS Models and Experimental Results

- Four BSS models were investigated as example problems that can be easily handled by the proposed algorithms. ($\lambda = 0.002$)

$$\text{IVA} \quad \text{Minimize}_{\mathbf{w}} \mathcal{I}(\mathbf{w}) + \|X\mathbf{w}\|_{2,1}$$

$$\text{sparse IVA} \quad \text{Minimize}_{\mathbf{w}} \mathcal{I}(\mathbf{w}) + \|X\mathbf{w}\|_{2,1} + \lambda \|X\mathbf{w}\|_1$$

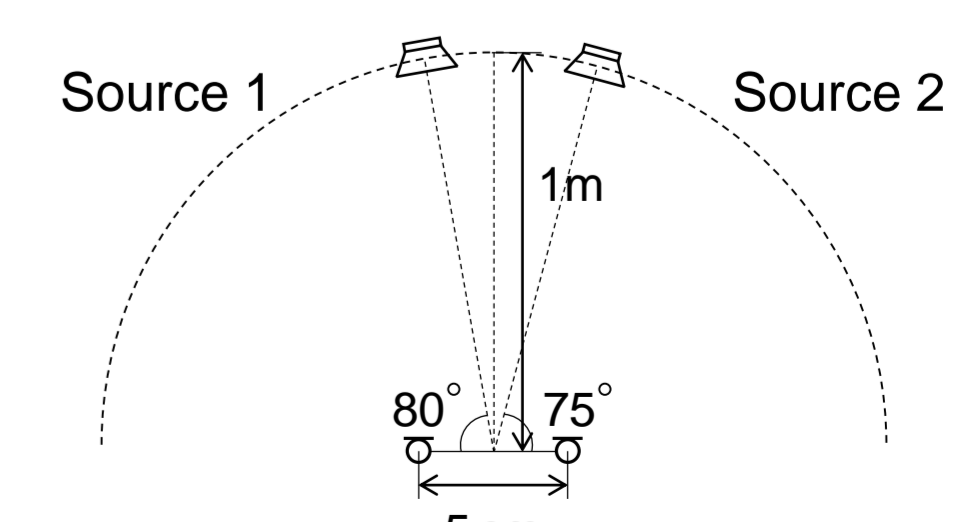
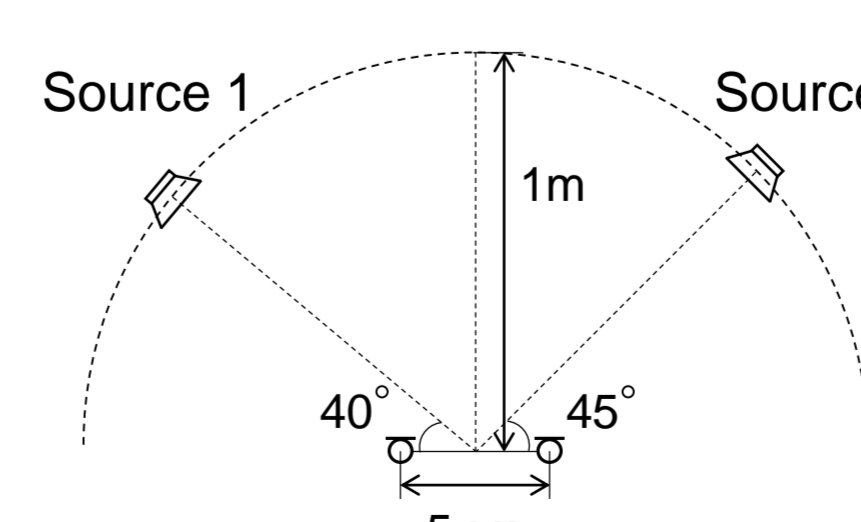
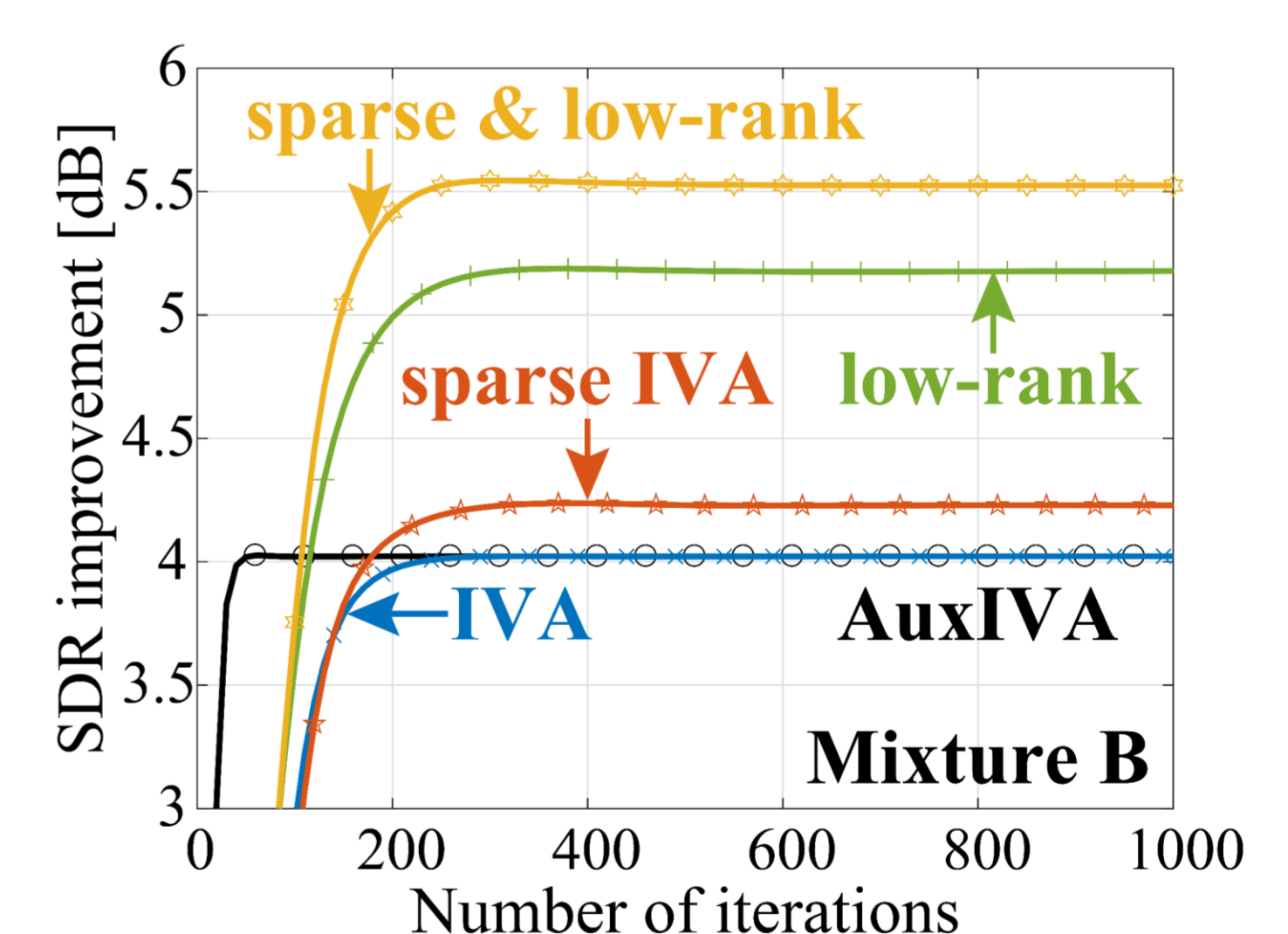
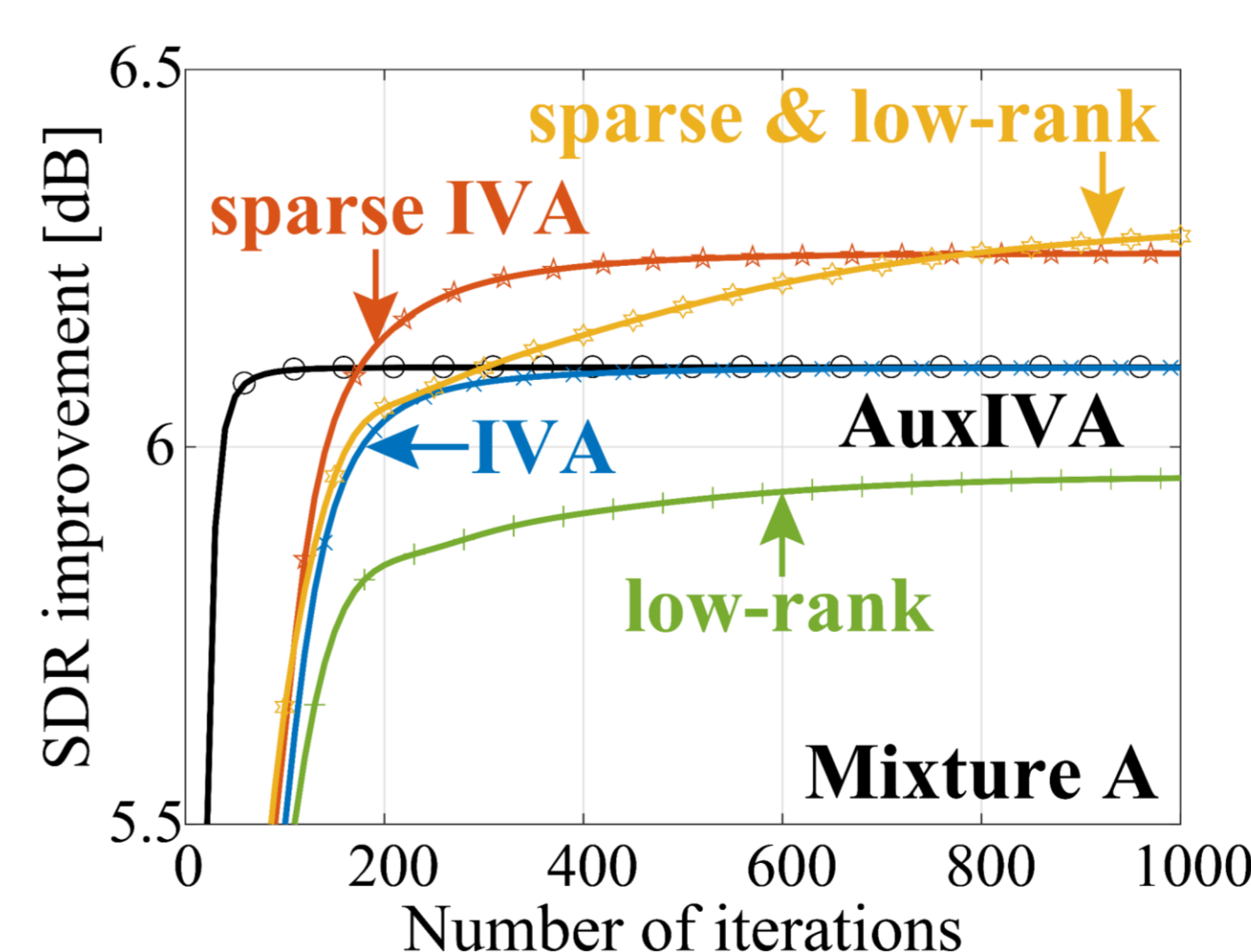
$$\text{low-rank (nuclear)} \quad \text{Minimize}_{\mathbf{w}} \mathcal{I}(\mathbf{w}) + \|X\mathbf{w}\|_*$$

$$\text{sparse \& low-rank} \quad \text{Minimize}_{\mathbf{w}} \mathcal{I}(\mathbf{w}) + \|X\mathbf{w}\|_* + \lambda \|X\mathbf{w}\|_1$$

- **Live recording** (liverec) of four female speech contained in UND task of the **SiSEC 2011** database was utilized as an test data.

- In the case of IVA, computational time per iteration of the proposed algorithm was **1.7x faster than AuxIVA** (MM algorithm [2]).

- Modification of the codes was able to be done within a few minutes (inside train) that indicate easiness of the proposed algorithms.



Reverberation time: 130ms, Window: 128ms Hann (half-overlap)