

## 1. Average Consensus Problem (ACP)

**SETUP:**  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  is a connected network with  $|\mathcal{V}| = n$  nodes (e.g., sensors) and  $|\mathcal{E}| = m$  edges (e.g., communication links). Node  $i \in \mathcal{V}$  stores a private value  $c_i \in \mathbb{R}$  (e.g., temperature).

**GOAL:** Compute the average of the private values (i.e., the quantity  $\bar{c} := \frac{1}{n} \sum_i c_i$ ) in a **distributed** fashion. That is, exchange of information can only occur along the edges.

## 2. Optimization Formulation of ACP

The optimal solution of the optimization problem

$$\min_{x \in \mathbb{R}^n} \frac{1}{2} \|x - c\|^2 = \frac{1}{2} \sum_i (x_i - c_i)^2 \quad \text{subject to} \quad x_i = x_j \quad \text{for all} \quad (i, j) \in \mathcal{E} \quad (1)$$

is  $x_i^* = \bar{c}$  for all  $i$ . The constraints can be written compactly as  $\mathbf{A}x = 0$ , where  $\mathbf{A} \in \mathbb{R}^{m \times n}$ , and the rows of the system enforce the constraints  $x_i = x_j$  for  $(i, j) \in \mathcal{E}$ .

**QUESTIONS:** Can we interpret Randomized Gossip (RG) algorithms for ACP as instances of specific randomized optimization methods for (1)? **Can new accelerated RG methods be developed this way?**

## 3. Accelerated Randomized Kaczmarz Method

**Best Approximation Problem:**  $\min_{x \in \mathbb{R}^n} \frac{1}{2} \|x - x^0\|^2 \quad \text{subject to} \quad \mathbf{A}x = b$

**Algorithm 1** Accelerated Randomized Kaczmarz Method (AccRK)

- 1: Data: Matrix  $\mathbf{A} \in \mathbb{R}^{m \times n}$ ; vector  $b \in \mathbb{R}^m$
- 2: Choose  $x^0 \in \mathbb{R}^n$  and set  $v^0 = x^0$
- 3: Parameters: Evaluate the sequences of the scalars  $\alpha_k, \beta_k, \gamma_k$  following one of two possible options.
- 4: **for**  $k = 0, 1, 2, \dots$ , **do**
- 5:  $y^k = \alpha_k v^k + (1 - \alpha_k) x^k$
- 6: Draw a fresh sample  $i_k \in [m]$  with probability  $1/m$
- 7:  $x^{k+1} = y^k - \frac{\mathbf{A}_{i_k} y^k - b_{i_k}}{\|\mathbf{A}_{i_k}\|^2} \mathbf{A}_{i_k}^\top$
- 8:  $v^{k+1} = \beta_k v^k + (1 - \beta_k) y^k - \gamma_k \frac{\mathbf{A}_{i_k} y^k - b_{i_k}}{\|\mathbf{A}_{i_k}\|^2} \mathbf{A}_{i_k}^\top$
- 9: **end for**

**Option 1, [4]:**  $\lambda \in [0, \lambda_{\min}^+(\mathbf{A}^\top \mathbf{A})]$  and  $\gamma_{-1} = 0$ .

Generate  $\{\gamma_k\}$ :  $\gamma_k$  largest root of

$$\gamma_k^2 - \frac{\gamma_k}{m} = (1 - \frac{\gamma_k}{\lambda} m) \gamma_{k-1}^2.$$

Generate  $\{\alpha_k\}$  and  $\{\beta_k\}$  by setting,

$$\alpha_k = (m - \gamma_k \lambda) / (\gamma_k (m^2 - \lambda)), \quad \beta_k = 1 - (\gamma_k \lambda) / m.$$

**Convergence Rate:** For  $k \rightarrow \infty$ ,

$$\mathbb{E}[\|x^k - x^*\|^2] \leq \left(1 - \sqrt{\lambda/m}\right)^k 4\lambda \|x^0 - x^*\|_{(\mathbf{A}^\top \mathbf{A})}^2$$

**Option 2, [3]:** Let  $\mathbf{W} = \frac{\mathbf{A}^\top \mathbf{A}}{2m}$  and

$$\nu = \max_{u \in \text{Range}(\mathbf{A}^\top)} \frac{u^\top [\sum_{i=1}^m \mathbf{A}_{i \cdot}^\top \mathbf{A}_{i \cdot} (\mathbf{A}^\top \mathbf{A})^\dagger \mathbf{A}_{i \cdot}^\top \mathbf{A}_{i \cdot}] u}{u^\top \frac{\mathbf{A}^\top \mathbf{A}}{m} u},$$

where  $1 \leq \nu \leq m$ .  $\{\beta_k\} \equiv 1 - \sqrt{\lambda_{\min}^+(\mathbf{W})/\nu}$ ,

$$\{\gamma_k\} \equiv \sqrt{\frac{1}{\lambda_{\min}^+(\mathbf{W})\nu}}, \quad \{\alpha_k\} \equiv \frac{1}{1 + \gamma_k \nu} \in (0, 1).$$

Let  $\Psi^k = \mathbb{E}[\|v^k - x^*\|_{\mathbf{W}^\dagger}^2 + \frac{1}{\mu} \|x^k - x^*\|^2]$ .

**Convergence Rate:**

$$\Psi^k \leq \left(1 - \sqrt{\lambda_{\min}^+(\mathbf{W})/\nu}\right)^k \Psi^0$$

## 4. Iteration Complexity

Let  $\mathbf{L}$  be the Laplacian matrix of the network  $\mathcal{G}$ . Then, for solving  $\mathbf{A}x = 0$ , where  $\mathbf{A}$  the incidence matrix,

**RK** [5, 2]:  $O\left(\left(\frac{2m}{\lambda_{\min}^+(\mathbf{L})}\right) \log(1/\epsilon)\right)$ ,

**AccRK (Option 1):**  $O\left(m\sqrt{2/\lambda_{\min}^+(\mathbf{L})} \log(1/\epsilon)\right)$ , **AccRK (Option 2):**  $O\left(\sqrt{2m\nu/\lambda_{\min}^+(\mathbf{L})} \log(1/\epsilon)\right)$

## 5. Accelerated Randomized Gossip Algorithms

**Algorithm 2** Accelerated Randomized Gossip Algorithm (AccGossip)

- 1: **Data:** Matrix  $\mathbf{A} \in \mathbb{R}^{m \times n}$  (normalized incidence matrix); vector  $b = 0 \in \mathbb{R}^m$
- 2: Choose  $x^0 \in \mathbb{R}^n$  and set  $v^0 = x^0$
- 3: **Parameters:** Evaluate the sequences  $\alpha_k, \beta_k, \gamma_k$  following option 1 or 2.
- 4: **for**  $k = 0, 1, 2, \dots$ , **do**
- 5: Each node  $\ell \in \mathcal{V}$  evaluates

$$y_\ell^k = \alpha_k v_\ell^k + (1 - \alpha_k) x_\ell^k.$$

- 6: Pick an edge  $e = (i, j) \in \mathcal{E}$  uniformly at random.

- Nodes  $i$  and  $j$  update their values via:

$$x_i^{k+1} = x_j^{k+1} = (y_i^k + y_j^k)/2$$

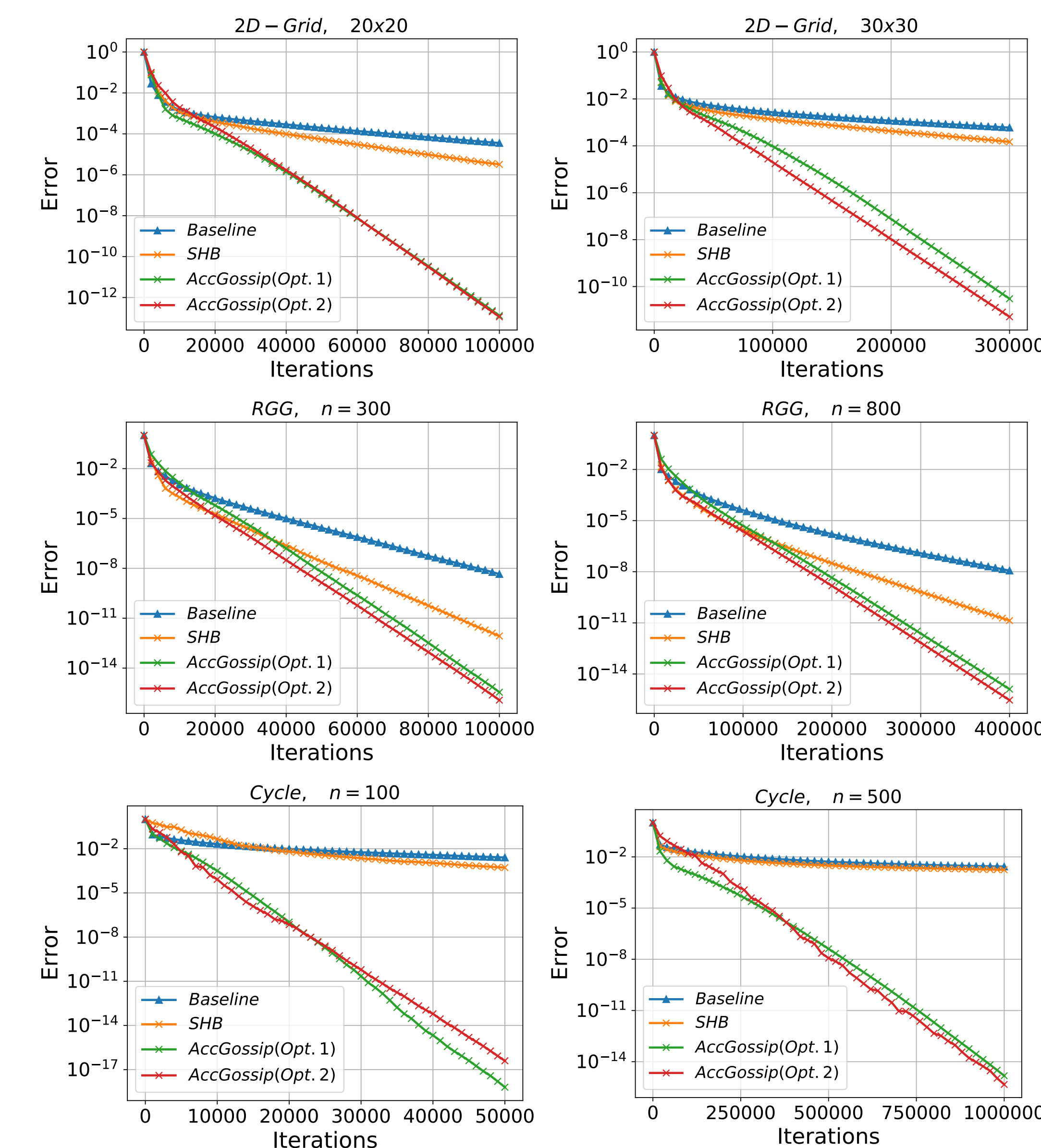
$$v_i^{k+1} = \beta_k v_i^k + (1 - \beta_k) y_i^k - \gamma_k (y_i^k - y_j^k)/2$$

$$v_j^{k+1} = \beta_k v_j^k + (1 - \beta_k) y_j^k - \gamma_k (y_j^k - y_i^k)/2$$

- Any other node  $\ell \in \mathcal{V} \setminus \{i, j\}$ :

$$x_\ell^{k+1} = y_\ell^k, \quad v_\ell^{k+1} = \beta_k v_\ell^k + (1 - \beta_k) y_\ell^k$$

- 7: **end for**



**Figure 1:** Performance of AccGossip in a 2-dim grid, random geometric graph (RGG) and a cycle graph. The baseline method corresponds to the randomized pairwise gossip algorithm proposed in [1] and SHB to the fast gossip algorithm proposed in [6]. The  $n$  in the title of each plot indicates the # of nodes of the network. For the grid graph, this is  $n \times n$ .

## 6. References

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