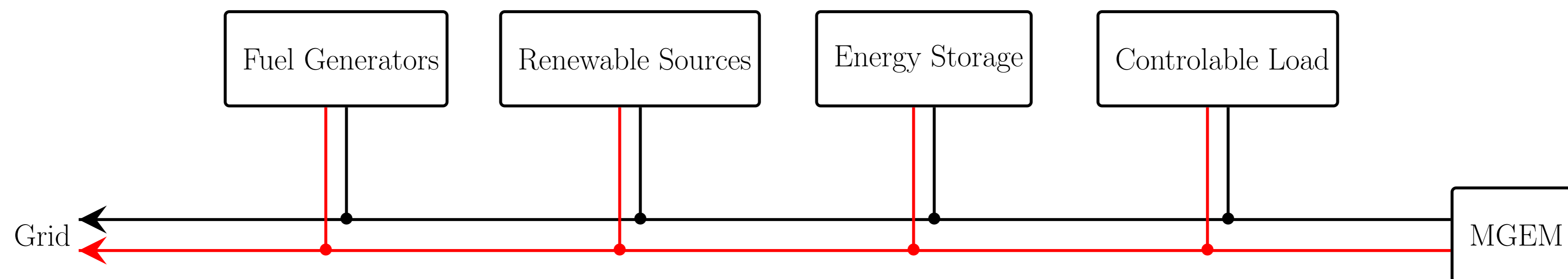




# Dynamic Power Allocation for Smart Grids via ADMM

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Decentralized infrastructure of a micro-grid with communications (black) and energy flow (red) networks. MGEM: Micro-grid energy manager

## Dynamic Economic Dispatch

$$\begin{aligned} \min_{\{\mathbf{p}_i \in \mathbb{R}^p\}_{i=1}^N} & \sum_{i=1}^M C_i(\mathbf{p}_i, k) - \sum_{i=M+1}^N U_i(\mathbf{p}_i, k) \\ \text{s.t.} & \underline{\mathbf{p}}_i^{[k]} \leq \mathbf{p}_i \leq \bar{\mathbf{p}}_i^{[k]}, i = 1, \dots, N \\ & \sum_{i=1}^N \mathbf{p}_i = \mathbf{P}^{[k]} \end{aligned}$$

**Assumption 1** (Regularity Conditions). Each of the *cost functions*  $C_i(\cdot, k)$  is  $\mu$ -strongly convex and has  $L$ -Lipschitz continuous gradient. Analogously, each *utility function*  $U_i(\cdot, k)$  is  $\mu$ -strongly concave and has  $L$ -Lipschitz continuous gradient.

**Assumption 2** (Bounded deviations). The primal dual optimal point from one iterate to the next deviates a bounded amount, i.e.  $\|\mathbf{p}^*(k) - \mathbf{p}^*(k+1)\| \leq \Delta p$  and  $\|\boldsymbol{\lambda}^{[k]} - \boldsymbol{\lambda}^{[k+1]}\| \leq \Delta \lambda \forall k$ .

What is desirable in a time varying set-up?

- If changes are fast, only one iteration per problem change
- Feasibility guarantees at each iteration
- Remain close to the optimal point

What do we require?

- For  $\lim_{k \rightarrow \infty} \|\mathbf{p}^{[k]} - \mathbf{p}^{*[k]}\| \leq C(\Delta p, \Delta \lambda)$  under the bounded deviations assumption we require  $Q$ -linear convergence of the method
- Strong coordination to guarantee per iteration constraint fulfillment

## ADMM friendly reformulation

$$\begin{aligned} \min_{\{\mathbf{p}_i\}_{i=1}^N, \{\mathbf{q}_i\}_{i=1}^N} & \sum_{i=1}^N f_i(\mathbf{p}_i) \triangleq \\ & \sum_{i=1}^M C_i(\mathbf{p}_i, k) - \sum_{i=M+1}^N U_i(\mathbf{p}_i, k) \\ \text{s.t.} & \underline{\mathbf{p}}_i^{[k]} \leq \mathbf{q}_i \leq \bar{\mathbf{p}}_i^{[k]}, i = 1, \dots, N \\ & \sum_{i=1}^N \mathbf{q}_i = \mathbf{P}^{[k]}, \\ & \mathbf{p}_i = \mathbf{q}_i, i = 1, \dots, N. \end{aligned}$$

Advantages of the problem formulation:

- For fixed  $\{\mathbf{q}_i\}_{i=1}^N$  the problem is smooth with strongly convex objectives in  $\{\mathbf{p}_i\}_{i=1}^N$
- For fixed  $\{\mathbf{p}_i\}_{i=1}^N$  the problem projects over the feasible set defined by the constraints. We will denote it by  $Q^{[k]}$ .

## ADMM with Total Feasibility

- 1: Initialize  $\{\mathbf{p}_i\}_{i=1}^N$  and  $\{\boldsymbol{\lambda}\}_{i=1}^N$ . Set  $k = 0$
- 2: Each node  $i$  obtains  $f_i(\cdot, k + 1)$ ,  $\underline{\mathbf{p}}_i^{[k+1]}$ ,  $\bar{\mathbf{p}}_i^{[k+1]}$ , MGEM obtains  $\mathbf{P}^{[k+1]}$
- 3: MGEM and nodes jointly compute

$$\begin{aligned} \{\mathbf{q}_i^{[k+1]}\}_{i=1}^N &= \arg \min_{\{\mathbf{q}_i\}_{i=1}^N} g(\mathbf{q}, k) \\ \text{s.t.} & \underline{\mathbf{p}}_i^{[k+1]} \leq \mathbf{q}_i \leq \bar{\mathbf{p}}_i^{[k+1]} \\ & \sum_{i=1}^N \mathbf{q}_i = \mathbf{P}^{[k+1]}, \end{aligned}$$

$$g(\mathbf{q}, k) \triangleq \frac{1}{2} \sum_{i=1}^N \left\| \mathbf{q}_i - \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right) \right\|^2.$$

See "Cooperative Projection"

- 4: Each node computes

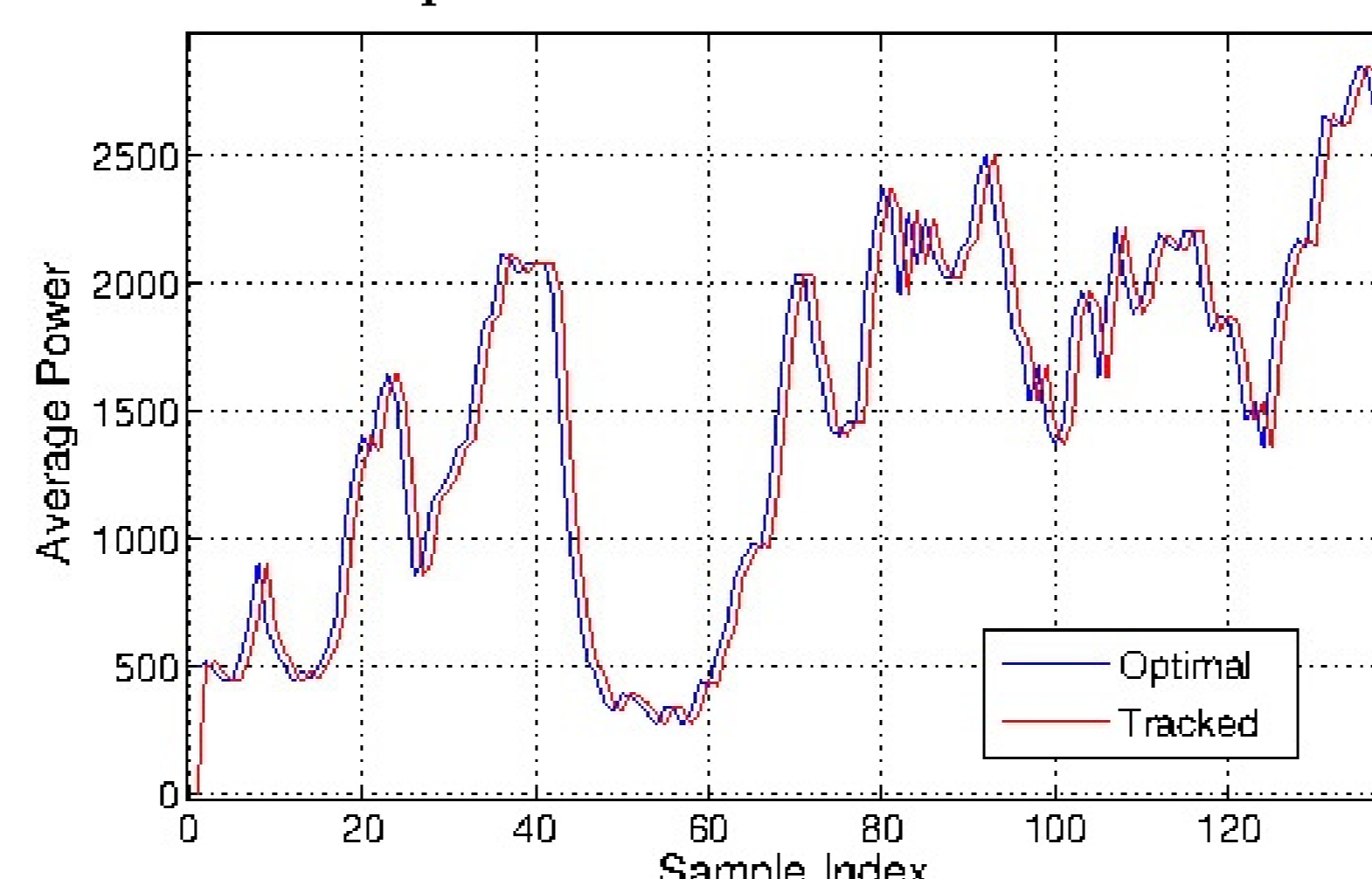
$$\begin{aligned} \mathbf{p}_i^{[k+1]} &= \arg \min_{\mathbf{p}_i} f_i(\mathbf{p}_i, k + 1) + \boldsymbol{\lambda}_i^{[k]T} \mathbf{p}_i \\ & \quad + \frac{\rho}{2} \left\| \mathbf{p}_i - \mathbf{q}_i^{[k+1]} \right\|^2 \\ \boldsymbol{\lambda}_i^{[k+1]} &= \boldsymbol{\lambda}_i^{[k]} + \rho \left( \mathbf{p}_i^{[k+1]} - \mathbf{q}_i^{[k+1]} \right) \end{aligned}$$

**Tracking statement** The proposed algorithm generates a sequence of iterates  $\{\mathbf{q}^{[k]}, \mathbf{p}^{[k]}\}$  that fulfills

$$\begin{aligned} \limsup_{k \rightarrow \infty} \|\mathbf{p}^{[k]} - \mathbf{p}^{*[k]}\| &\leq c_1 \\ \limsup_{k \rightarrow \infty} \|\mathbf{q}^{[k]} - \mathbf{q}^{*[k]}\| &\leq c_2 \end{aligned}$$

where  $c_1 \triangleq \frac{g}{\sqrt{1+\delta_{\max}-1}}$ ,  $c_2 \triangleq 3c_1^2 + \frac{1}{\rho}g^2 + \frac{3}{\sqrt{\rho}}c_1g$ ,  $\delta_{\max} \triangleq \frac{1}{\sqrt{L/\mu}}$  and  $g \triangleq \sqrt{\rho(\Delta p)^2 + \frac{1}{\rho}(\Delta \lambda)^2}$ , and the sequence  $\{\mathbf{q}^{[k]}\}$  is always primal feasible.

## Numerical Experiments



## Cooperative Projection

- 1: Nodes compute

$$\mathbf{m}_i = \left( \mathcal{P}_{Q_i^{[k+1]}} - \mathcal{I} \right) \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right)$$

- 2: Transmit  $\mathcal{P}_{Q_i^{[k+1]}} \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right)$  MGEM receives

$$\mathbf{s}^{[k]} = \sum_{i=1}^N \mathcal{P}_{Q_i^{[k+1]}} \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right)$$

- 3: MGEM transmits  $\mathbf{b}^{[k]} \triangleq \text{sign}(-\mathbf{s}^{[k]} + \mathbf{P}^{[k+1]})$

- 4: **for**  $j = 1, \dots, p$  **do**

- 5: **if**  $\text{sign}(b^{[k]}(j)) = \text{sign}(m_i(j)) |m_i(j)| = 0$  **then**

- 6: **if**  $m_i(j) > 0$  **then**

- 7: Node transmits  $(m_i(j), \bar{x}_{ij})$ ,  $\bar{x}_{ij} = \bar{p}_i(j) - \mathcal{P}_{Q_i^{[k+1]}} \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right)(j)$

- 8: **end if**

- 9: **if**  $m_i(j) < 0$  **then**

- 10: Node transmits  $(m_i(j), \underline{x}_{ij})$ ,  $\underline{x}_{ij} = \underline{p}_i(j) - \mathcal{P}_{Q_i^{[k+1]}} \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right)(j)$

- 11: **end if**

- 12: **if**  $m_i(j) = 0$  **then** node tx  $(0, \bar{x}_{ij})$  if  $b_i(j) > 0$  and  $\underline{x}_{ij}$  otherwise.

- 13: **end if**

- 14: **end if**

- 15: **end for**

$\mathcal{T}(j)$  set of nodes transmits regarding component  $j$

- 16: MGEM solves

$$\begin{aligned} \min_{\{\Delta q_{ij} \in \mathcal{T}(j), j=1, \dots, p\}} & \sum_{ij} \|\Delta q_{ij} + m_i(j)\|^2 \\ \text{s.t.} & 0 \leq \Delta q_{ij} \leq \bar{x}_{ij} \\ & \quad \text{if } b_i^{[k]}(j) > 0 \\ & \underline{x}_{ij} \leq \Delta q_{ij} \leq 0 \\ & \quad \text{if } b_i^{[k]}(j) < 0 \end{aligned}$$

- 17: MGEM transmits  $\Delta q_{ij}$

- 18: Each node computes  $q_i^{[k+1]}(j) = \Delta q_{ij} + \mathcal{P}_{Q_i^{[k+1]}} \left( \mathbf{p}_i^{[k]} + \frac{\boldsymbol{\lambda}_i^{[k]}}{\rho} \right)(j)$

## Information Exchange

Each iteration requires

- In worse case scenario, exchange of  $3N$   $p$ -sized real vectors and broadcast 1  $p$  sized binary vector as compared to  $3N$   $p$ -sized real vectors if all information is sent to MGEM.
- this scenario will occur when no user lays in the boundary of the constraint set in any component.